

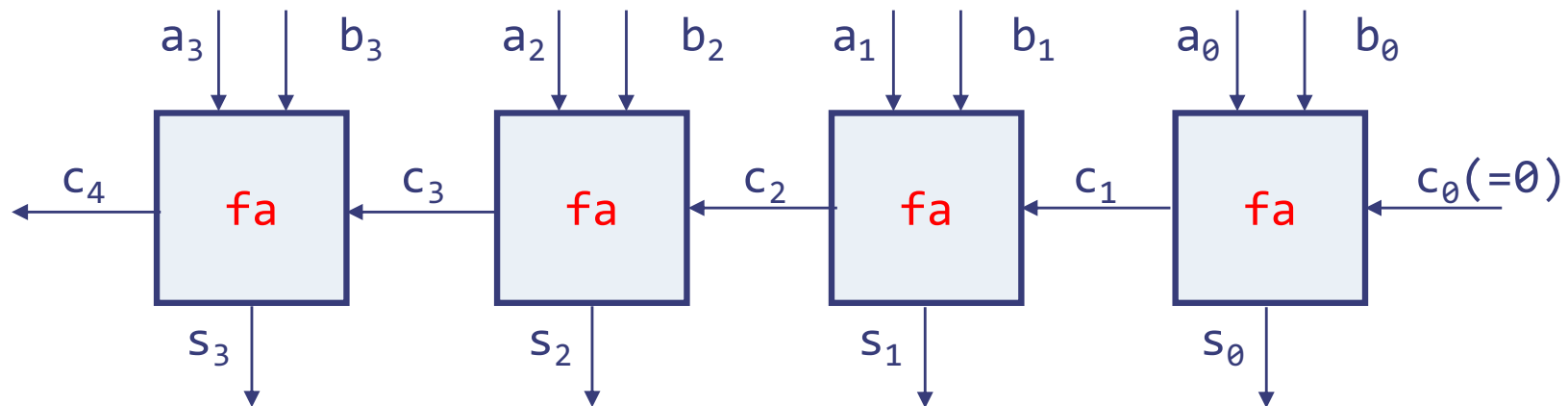
Combinational Circuits in Bluespec

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Ripple-carry adder

- Cascade FAs to perform binary addition



Full Adder: A one-bit adder

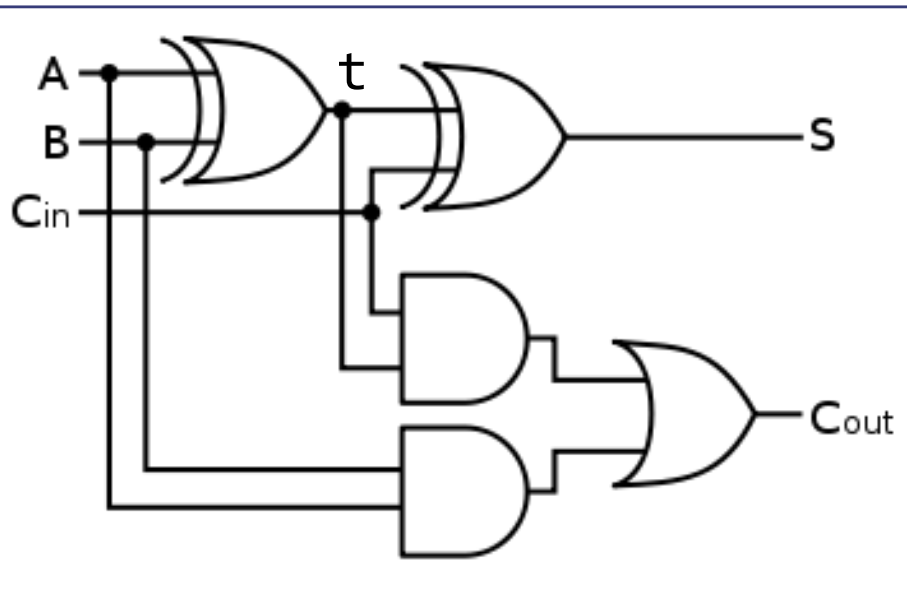
Boolean Equations

$$t = a \oplus b$$

$$s = t \oplus c_{in}$$

$$c_{out} = a \cdot b + c_{in} \cdot t$$

```
function fa(a, b, c_in);  
    t = (a ^ b);  
    s = t ^ c_in;  
    c_out = (a & b) | (c_in & t);  
    result[0] = s;  
    result[1] = c_out;  
    return result;  
endfunction
```



Structural code – only specifies interconnection between boxes

Not quite correct - needs type annotations

Full Adder: A one-bit adder

corrected

```
function Bit#(2) fa(Bit#(1) a, Bit#(1) b, Bit#(1) c_in);
  Bit#(1) t;
  Bit#(1) s;
  Bit#(1) c_out;
  Bit#(2) result;
  t = a ^ b;
  s = t ^ c_in;
  c_out = (a & b) | (c_in & t);
  result[0] = s;
  result[1] = c_out;
  return result;
endfunction
```

Type declaration
“`Bit#(1) a`” says
that `a` is one bit wide

Full Adder *more convenient syntax*

```
function Bit#(2) fa(Bit#(1) a, Bit#(1) b, Bit#(1) c_in);  
    Bit#(1) t = a ^ b;  
    Bit#(1) s = t ^ c_in;  
    Bit#(1) c_out = (a & b) | (c_in & t);  
  
    return {c_out, s};  
endfunction
```

A variable's type declaration and definition can be combined

{c_out, s} represents bit concatenation; can be used to avoid naming intermediate results

How big is {c_out, s}?

Type checking

- The Bluespec compiler checks if all the declared types are used consistently

```
function Bit#(2)3 fa(Bit#(1) a, Bit#(1) b, Bit#(1) c_in);  
    Bit#(1) t = a ^ b;  
    Bit#(1) s = t ^ c_in;  
    Bit#(1) c_out = (a & b) | (c_in & t);  
    return {c_out, s};  
endfunction
```

The compiler will flag this as an error because {c_out, abc[0]} is Bit#(2)

- In fact, the compiler can reduce the programmer's burden by deducing some types and not asking for explicit type declarations

⇒ The "let" syntax

“let” syntax

```
function Bit#(2) fa(Bit#(1) a, Bit#(1) b, Bit#(1) c_in);  
  Bit#(1) t = a ^ b;  
  Bit#(1) s = t ^ c_in;  
  Bit#(1) c_out = (a & b) | (c_in & t);  
  return {c_out, s};  
endfunction
```

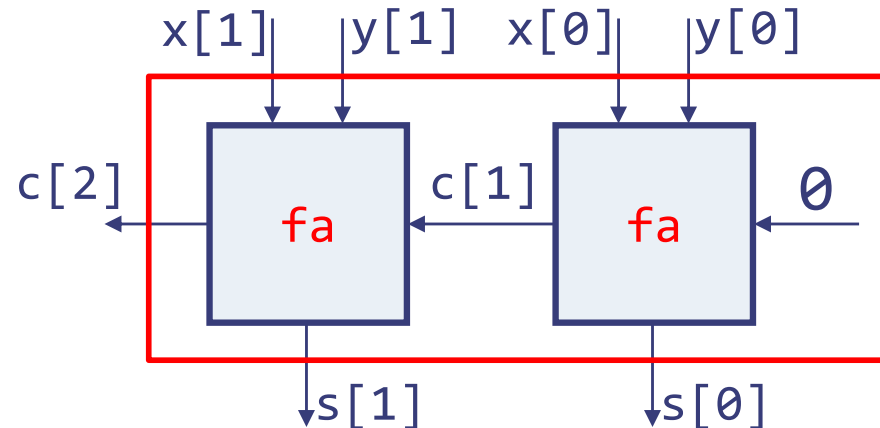
```
function Bit#(2) fa(Bit#(1) a, Bit#(1) b, Bit#(1) c_in);  
  let t = a ^ b;  
  let s = t ^ c_in;  
  let c_out = (a & b) | (c_in & t);  
  return {c_out, s};  
endfunction
```

Types of t, s and c_out can be deduced from the types of the corresponding right-hand-side expressions

“let” syntax is very convenient, we will use it extensively

2-bit Ripple-Carry Adder

cascading full adders



Use **fa** as a black-box

```
function Bit#(3) add2(Bit#(2) x, Bit#(2) y);
```

```
  let s = 2b'00; Bit#(3) c = 3b'000;
```

```
  c[0] = 0;
```

```
  let cs0 = fa(x[0], y[0], c[0]);
```

```
  s[0] = cs0[0]; c[1] = cs0[1];
```

```
  let cs1 = fa(x[1], y[1], c[1]);
```

```
  s[1] = cs1[0]; c[2] = cs1[1];
```

```
  return {c[2], s};
```

```
endfunction
```

s has two wires,
initially each s
wire is zero

wire s[0] is
updated

wire s[1] is
updated

The same as writing
{c[2], s[1], s[0]};

32-bit Ripple-Carry Adder (RCA)

- We could have written the chain of RCA explicitly, but we can also use loops!

```
function Bit#(33) add32(Bit#(32) x, Bit#(32) y, Bit#(1) c0);
    Bit#(32) s = 0;
    Bit#(33) c = 0;
    c[0] = c0;
    for (Integer i=0; i<32; i=i+1) begin
        Bit#(2) cs = fa(x[i],y[i],c[i]);
        c[i+1] = cs[1];
        s[i] = cs[0];
    end
    return {c[32],s};
endfunction
```

Now we discuss how the gates are generated (synthesized) from a loop

Loop is unfolded by the compiler

```
for(Integer i=0; i<32; i=i+1) begin
    Bit#(2) cs = fa(x[i], y[i], c[i]);
    c[i+1] = cs[1];
    s[i] = cs[0];
end
```

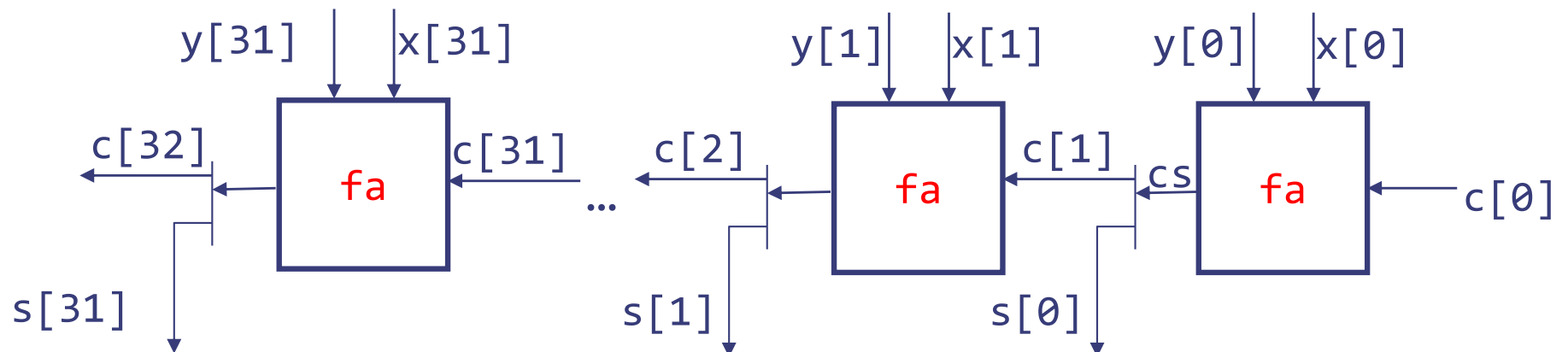
```
cs = fa(x[0], y[0], c[0]);
c[1] = cs[1];
s[0] = cs[0];
cs = fa(x[1], y[1], c[1]);
c[2] = cs[1];
s[1] = cs[0];
...
cs = fa(x[31], y[31], c[31]);
c[32] = cs[1];
s[31] = cs[0];
```

cs in the loop body is a local variable. Hence each of these cs refers to a different value. We could have named them cs0, ... cs31.

Loops to gates

```
cs0 = fa(x[0], y[0], c[0]); c[1]=cs0[1]; s[0]=cs0[0];  
cs1 = fa(x[1], y[1], c[1]); c[2]=cs1[1]; s[1]=cs1[0];  
...  
cs31 = fa(x[31], y[31], c[31]);  
c[32] = cs31[1]; s[31] = cs31[0];
```

Unfolded loop defines an acyclic wiring diagram



Each instance of function **fa** is replaced by its body

Types in Bluespec

Types

- Every expression in a Bluespec program has a type
- A type is a *grouping* of values, examples
 - `Bit#(16)` // 16-bit wide bit-vector (16 is a numeric type)
 - `Bool` // 1-bit value representing True or False
 - `Vector#(16, Bit#(8))` // Vector of size 16 containing `Bit#(8)`'s
- A type declaration can be parameterized by other types using the syntax ``#'`, for example
 - `Bit#(n)` represents *n* bits, e.g., `Bit#(8)`, `Bit#(32)`, ...
 - `Tuple2#(Bit#(8), Integer)` represents a pair of 8-bitvector and an integer.
 - `function Bit#(8) fname (Bit#(8) arg)` represents a function from `Bit#(8)` to `Bit#(8)` values
- A *type name* always begins with a capital letter, while a *variable identifier* begins with a small letter

Type synonyms

```
typedef Bit#(8) Byte;
```

```
typedef Bit#(32) Word;
```

```
typedef Tuple2#(a,a) Pair#(type a);
```

type variable



```
typedef 32 DataSize;
```

numeric type



```
typedef Bit#(DataSize) Data;
```

Enumerated types

A very useful typing concept

- Suppose we have a variable `c` whose values can represent three different colors
 - Declare the type of `c` to be `Bit#(2)` and adopt the convention that `00` represents Red, `01` Blue and `10` Green
- A better way is to create a new type called `Color`:

```
typedef enum {Red, Blue, Green}
Color deriving(Bits, Eq);
```

Why is this way better?

- Bluespec compiler automatically assigns a bit representation to the three colors and provides a function to test whether two colors are equal
- If you do not use “deriving” then you will have to specify your own encoding and equality function

Types prevent us from
mixing colors with raw bits

Parameterized Circuits

n-bit Ripple-Carry Adder

```
function Bit#(n+1) addN(Bit#(n) x, Bit#(n) y, Bit#(1) c0);
  Bit#(n) s = 0;
  Bit#(n+1) c = 0;
  c[0] = c0;
  for (Integer i=0; i<n; i=i+1) begin
    let cs = fa(x[i],y[i],c[i]);
    c[i+1] = cs[1];
    s[i] = cs[0];
  end
  return {c[n],s};
endfunction
```

Now can instantiate different sized adders by specifying n

Unfortunately, there are several subtle type errors in this program - we will fix them one by one

Fixing the type errors

Parameterized Ripple-Carry Adder

```
function Bit#(n+1) addN(Bit#(n) x, Bit#(n) y, Bit#(1) c0);  
  Bit#(n) s = 0;  
  Bit#(n+1) c = 0;  
  c[0] = c0;  
  for (Integer i=0; i<n; i=i+1) begin  
    let cs = fa(x[i],y[i],c[i]);  
    c[i+1] = cs[1];  
    s[i] = cs[0];  
  end  
  return {c[n],s};  
endfunction
```

- n is numeric type and Bluespec does not allow arithmetic on types, e.g., $n+1$, $i<n$, $c[n]$ are illegal!

Fixing the type errors

valueOf(n) versus n

- Each expression has a *type* and a *value*, and these two come from entirely disjoint worlds
- n in $\text{Bit}\#(n)$ is a *numeric type* variable and resides in the types world
- Sometimes we need to use values from the types world in actual computation. The function `valueOf` extracts the integer from a numeric type
 - Thus,
 - $i < n$ is not type correct
 - $i < \text{valueOf}(n)$ is type correct

Fixing the type errors

TAdd#(n, 1) versus n+1

- Sometimes we need to perform operations in the types world that are very similar to the operations in the value world
 - Examples: Addition, Multiplication, Logarithm base 2, ...
- Bluespec defines a few special operators in the types world for such operations
 - TAdd#(m,n), TSub#(m,n), TMul#(m,n), TDiv#(m,n), TLog#(n), TExp#(n), TMax#(m,n), TMin#(m,n)
 - Thus,
 - Bit#(n+1) is not type correct
 - Bit#(TAdd#(n,1)) is type correct

Parameterized Ripple-Carry Adder

corrected

```
function Bit#(TAdd#(n,1)) addN(Bit#(n) x, Bit#(n) y,  
                                Bit#(1) c0);
```

```
    Bit#(n) s = 0;
```

```
    Bit#(TAdd#(n,1)) c;
```

```
    c[0] = c0;
```

```
    let valn = valueOf(n);
```

```
    for (Integer i=0; i<valn; i=i+1) begin
```

```
        let cs = fa(x[i], y[i], c[i]);
```

```
        c[i+1] = cs[1];
```

```
        s[i] = cs[0];
```

```
    end
```

```
    return {c[valn],s};
```

```
endfunction
```

types world
equivalent of n+1

Lifting a type into
the value world

Instantiating the parametric Adder

```
function Bit#(Tadd#(n,1)) addN(Bit#(n) x, Bit#(n) y, Bit#(1) c0);
```

How do we define a add32, add3 ... using addN ?

```
// concrete instances of addN!
```

```
function Bit#(33) add32(Bit#(32) x, Bit#(32) y, Bit#(1) c0)  
    = addN(x,y,c0);
```

```
function Bit#(4) add3(Bit#(3) x, Bit#(3) y, Bit#(1) c0)  
    = addN(x,y,c0);
```

The numeric type n on the RHS implicitly gets instantiated to 32 and 3, respectively, because of the LHS declaration

Bluespec is for describing circuits

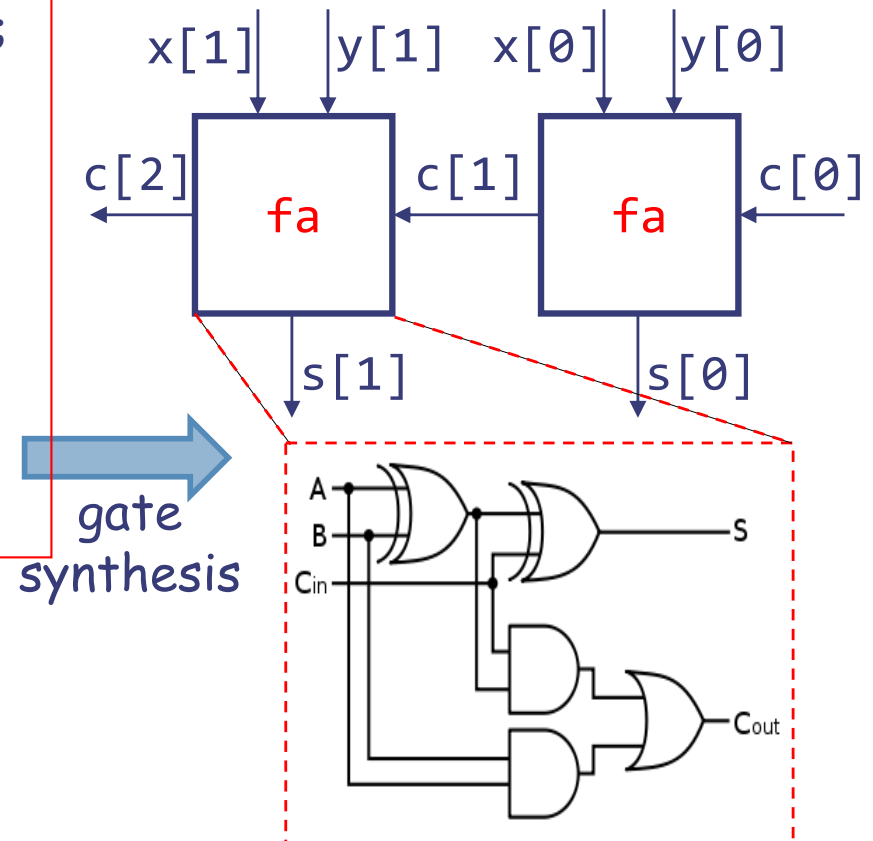
- Bluespec is like a language for drawing pictures of interconnected boxes
- Boxes happen to be Boolean gates with inputs and outputs
- However, unlike ordinary pictures, our boxes, i.e., gates, have computational meaning, and therefore, we can ask what values a circuit would produce on its output lines, given a specific set of values on its input lines
- Even though the primary purpose of the Bluespec compiler is to synthesize a network of gates, the ability to simulate the functionality of the resulting circuit is extremely important

Bluespec: Gate synthesis versus simulation 2-bit adder

```
function Bit#(3) add2(Bit#(2) x, Bit#(2) y);
  let s = 2b'00;      Bit#(3) c = 3b'000;
  c[0] = 0;
  let cs0 = fa(x[0], y[0], c[0]);
  s[0] = cs0[0];    c[1] = cs0[1];
  let cs1 = fa(x[1], y[1], c[1]);
  s[1] = cs1[0];    c[2] = cs1[1];
return {c[2],s};
endfunction
```

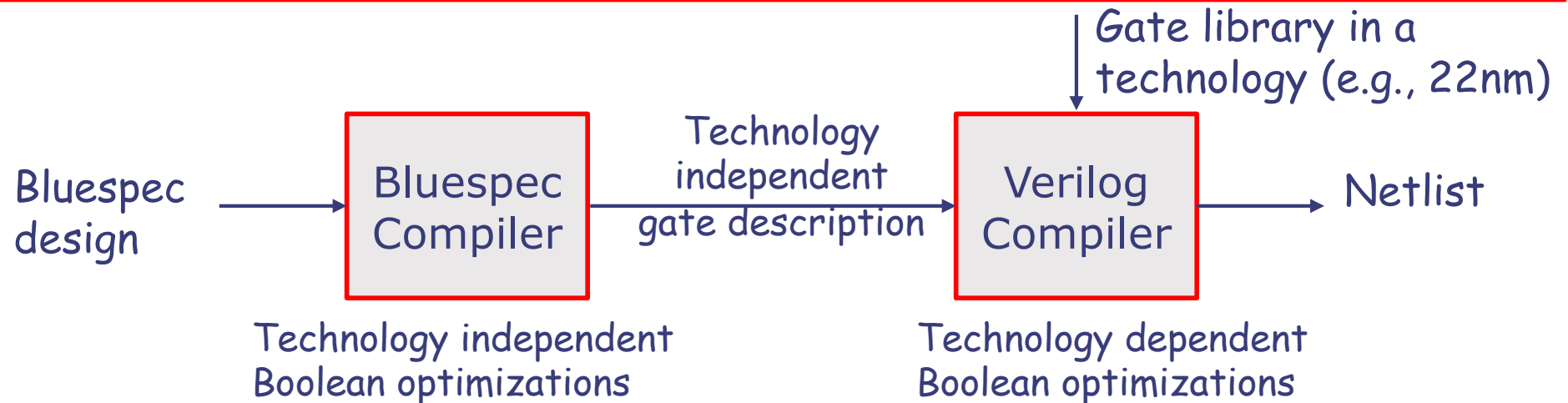
simulate

- $\text{add2}(2'b11, 2'b01) \Rightarrow 3'b100$
- $\text{add2}(2'b01, 2'b01) \Rightarrow 3'b010$



Caution: In spite of the fact that Bluespec programs, like programs in other software languages, produce outputs given inputs, the purpose of Bluespec programs is to describe circuits

Compiling Bluespec into circuits



- *Static elaboration:* Bluespec compiler eliminates all constructs which have no direct hardware meaning
 - All data structures are converted into bit vectors
 - Loops are unfolded
 - Functions are in-lined
 - What remains is an acyclic graph of Boolean gates
 - The compiler complains if it detects a cycle in your circuit

Takeaway

- Once we define a combinational circuit, we can use it repeatedly to build larger circuits
- Bluespec compiler, because of the type signatures of functions, prevents us from connecting functions and gates in obviously illegal ways
- We can use loop constructs and functions to express combinational circuits, but all loops are unfolded and functions are in-lined during the compilation phase
- We can also write parameterized circuits in Bluespec, for example an n-bit adder. Once n is specified, the correct circuit is automatically generated

The best way to learn about types is to try writing a few expressions and feeding them to the compiler