Accelerators (I)

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"Compute has been the oxygen of deep learning" – Ilya Sutskever (Open AI)

GPU Usage for ImageNet Challenge



Challenges

From EE Times – September 27, 2016

"Today the job of training machine learning models is limited by compute, if we had faster processors we'd run bigger models...in practice we train on a reasonable subset of data that can finish in a matter of months. We could use improvements of several orders of magnitude – 100x or greater."

– Greg Diamos, Senior Researcher, SVAIL, Baidu

Compute Demands Growing Exponentially

AlexNet to AlphaGo Zero: A 300,000x Increase in Compute

Deep and steep

Computing power used in training AI systems Days spent calculating at one petaflop per second*, log scale



What is Moore's Law

- every two years*
- CPU performance will double every two years*
- Chip performance will double every two years*
- The speed of transistors will double every two years*
- Transistors will shrink to half size every two years*
- Transistors per die will double every two years*
- The economic sweet spot for the number of devices on a chip will double every two years*

* Or 18 months...

Compute Demands for Deep Neural Networks

Common carbon footprint benchmarks

in lbs of CO2 equivalent



Chart: MIT Technology Review

[Strubell, ACL 2019]

Energy and Power Consumption







[Dennard et al., "Design of ion-implanted MOSFET's with very small physical dimensions", JSSC 1974]



[Dennard et al., "Design of ion-implanted MOSFET's with very small physical dimensions", JSSC 1974]



[Dennard et al., "Design of ion-implanted MOSFET's with very small physical dimensions", JSSC 1974]









Gen X+1















	Gen X	Gen X+1
Gate Width	1.0	0.7
Device Area/ Capacitance	1.0	0.5 0.5
•	1.0	0.7 0.7
Voltage	1.0	0.7
Energy ~ <u>1.0 x 1.0²</u> = 1.0		~ <u>2 x 0.7 x 0.7</u> ² = 0.65
Delay 1.0		0.7
Frequency $1/1.0 = 1.0$		1/ <mark>0.7</mark> = 1.4
Power ~ <u>1.0 x</u>	<u>x 1.0²</u> x 1.0 = 1.0	

[Dennard et al., "Design of ion-implanted MOSFET's with very small physical dimensions", JSSC 1974]

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[Dennard et al., "Design of ion-implanted MOSFET's with very small physical dimensions", JSSC 1974]

Technology Trends



Figure 2. Sources of computing performance have been challenged by the end of Dennard scaling in 2004. All additional approaches to further performance improvements end in approximately 2025 due to the end of the roadmap for improvements to semiconductor lithography. Figure from Kunle Olukotun, Lance Hammond, Herb Sutter, Mark Horowitz and extended by John Shalf. (Online version in colour.)

During the Moore + Dennard's Law Era

- Instruction-level parallelism (ILP) was largely mined out by early 2000s
- Voltage (Dennard) scaling ended in 2005
- Hit the power limit wall in 2005
- Performance is coming from parallelism using more transistors since ~2007
- But....

Technology Trends



Architecture Metrics

- Speed The rate at which the hardware finishes tasks. Limited by the number of computation units and their utilization.
- Energy The total energy, e.g., in Joules, consumed to perform a task. Often constrained by battery capacity or desire to reduce carbon footprint.
- Power The rate at which energy is consumed, e.g., in Watts. Often limited by delivery or packaging constraints
- Accuracy The precision of the results produced. Can be dictated by bit width of compute units.
- Flexibility The range of problems that can be solved, which is constrained by the limitations of the architecture.

Deep Learning Platforms

- CPU
 - Intel, ARM, AMD...
- GPU
 - NVIDIA, AMD...
- Fine Grained Reconfigurable (FPGA)
 - Xilinx, Altera (Microsoft BrainWave)
- Coarse Grained Programmable/Reconfigurable
 - Wave Computing, Graphcore, Samba Nova...
- Application Specific
 - Neuflow, *DianNao, Eyeriss, TPU, Cnvlutin, SCNN, ...

There Are Myriad Tensor Accelerators

Selected tensor accelerator designs





Eyeriss V2 [JETCAS2019]



SCNN [ISCA2017]



ExTensor [MICRO2019]



Gamma [ASPLOS2021]



RAELLA [ISCA2023]











What Can Einsums Do For You?

- Simultaneously more precise and concise representation
- Extends tensor algebra far beyond matrix multiplication
- Now includes algorithms for: AI, graphs, fft, crypto, point cloud
- Intermediate point between GEMM and full programmability
- Allows for bounds analysis (SoL) on compute and data movement
- Admits of optimization via algebraic transformations
- Tip of the pyramid of separation of concerns
Rank-0: Scalar







Rank-2: Matrix





Rank-2: Matrix

Rank-3: Cube















Rank-2: Matrix



Rank-3: Cube





Rank-2: Matrix



Rank-3: Cube













• The elements of each "rank" (dimension) are identified by their "coordinates", e.g., rank H has coordinates 0, 1, 2



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- Each element of the tensor is identified by the tuple of coordinates from each of its ranks, i.e., a "point".
 So (1,2) -> "f"



- The elements of each "rank" (dimension) are identified by their "coordinates", e.g., rank H has coordinates 0, 1, 2
- Each element of the tensor is identified by the tuple of coordinates from each of its ranks, i.e., a "point".
 So (1,2) -> "f"

Matrix Multiply



Matrix Multiply









A = Tensor(shape=[M, K]) B = Tensor(shape=[N, K]) Z = Tensor(shape=[M, N]) for n in [0..N): for m in [0..M):

```
for m in [0...M):
for m in [0...M):
    for k in [0...K):
        Z[m][n] += A[m][k] × B[n][k]
```





$$Z_{m,n} = A_{m,k} \times B_{n,k}$$

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Operational Definition for Einsums (ODE):

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- Traverse all points in space of all legal index values (iteration space)

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Operational Definition for Einsums (ODE):

- Traverse all points in space of all legal index values (iteration space)
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 - Calculate value on right hand at specified indices for each operand

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 - Assign value to operand at specified indices on left hand side
 - Unless that operand is non-zero, then reduce value into it

$$Z_{m,n} = \sum_{k} A_{m,k} \times B_{n,k}$$

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Einsum – More notation

$$Z_{m,n}^{\boldsymbol{M},\boldsymbol{N}} = A_{m,k}^{\boldsymbol{M},\boldsymbol{K}} \times B_{n,k}^{\boldsymbol{N},\boldsymbol{K}}$$

Superscript represents the name of rank of the tensor

Einsum – More notation

$$Z_{m,n}^{M,N} = A_{m,k}^{M,K} \times B_{n,k}^{N,K}$$

Superscript represents the name of rank of the tensor

$$Z_{m,n}^{M=3,N=3} = A_{m,k}^{M=3,K=6} \times B_{n,k}^{N=3,K=6}$$

Equals in superscript represents the shape of the rank of the tensor, by default rank shape is assumed to be same as rank name

Einsum-level analysis - MM

$$Z_{m,n}^{M,N} = A_{m,k}^{M,K} \times B_{n,k}^{N,K}$$

Number of multiplies: $M \times N \times K$

Minimum amount of data to read: $M \times K + N \times K$

Minimum amount of data to write: $M \times N$

Einsum-level analysis – M dot products

$$Z_m^M = A_{m,k}^{M,K} \times B_{m,k}^{M,K}$$

Number of multiplies: $M \times K$

Minimum amount of data to read: $M \times K + M \times K$

Minimum amount of data to write: M

$$Z_{m,n} = A_{m,k} \times B_{n,k}$$

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Convolution (CONV) Layer



CONV Computation


Cycle through input fmap and weights (hold psum of output fmap) input fmap output fmap 3 filters 3 8 3 2 2 ← 3 2 _____ Filter overlay L______ 8 Incomplete partial sum \leftarrow 2

Cycle through input fmap and weights (hold psum of output fmap) input fmap output fmap 3 filters 37 8 3 2 2 2 ← 3 2 _____ Filter overlay _____ 2 8 Incomplete partial sum ← 2

Cycle through input fmap and weights (hold psum of output fmap) input fmap filters



2

2

Cycle through input fmap and weights (hold psum of output fmap) input fmap output fmap 3 filters 3 18 3 2 2 ← 3 2 _____ Filter overlay L______ 2 8 Incomplete partial sum 2 \leftarrow



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Cycle through input fmap and weights (hold psum of output fmap)



Cycle through input fmap and weights (hold psum of output fmap)



CONV Layer Implementation

Naïve 7-layer for-loop implementation:



Einsum - Convolution

$$O_{p,q,m} = I_{c,p+r,q+s} \times F_{m,c,r,s}$$

- N Number of input fmaps/output fmaps (batch size)
- C Number of channels in input fmaps (activations) & filters (weights)
- H Height of input fmap (activations)
- W Width of input fmap (activations)
- R Height of filter (weights)
- S Width of filter (weights)
- M Number of channels in output fmaps (activations)
- P Height of output fmap (activations)
- Q Width of output fmap (activations)
- U Stride of convolution

CONV Variants

- -Depthwise layer M == C and $\forall_{c \mid = m} F_{c,m,r,s} = 0$
- -Pointwise layer R == S == 1
- -Matrix multiply R == S == H == 1
- -Compress (pointwise) M < C and R == S == 1
- -Expand (pointwise) M > C and R == S == 1

Compress...Expand sequences are called a "bottleneck"

Toeplitz (IM2COL) Cascade

$$O_{p,q,m} = I_{c,p+r,q+s} \times F_{m,c,r,s}$$

$$T_{c.p,q,r,s} = I_{c,p+r,q+s}$$
$$O_{p,q,m} = T_{c,p,q,r,s} \times F_{m,c,r,s}$$

Conventional Transformer Diagram

This time, composition of many kernels

Not a precise description of functionality.



Figure 1: The Transformer - model architecture.

[Attention, Vaswani et al. 2016]

Multi-head Attention (without initial embedding step)

$$K_{b,h,m,e} = I_{b,m,d} \times WK_{d,h,e}$$

$$Q_{b,h,m,e} = I_{b,m,d} \times WQ_{d,h,e}$$

$$QK_{b,h,m,p}^{B,H,M,P=M} = Q_{b,h,p,e}^{B,H,M,E} \times K_{b,h,m,e}$$

$$SN_{b,h,m,p} = exp(QK_{b,h,m,p})$$

$$SD_{b,h,p} = SN_{b,h,m,p}$$

$$A_{b,h,m,p} = SN_{b,h,m,p}/SD_{b,h,p}$$

$$V_{b,h,m,f} = I_{b,m,d} \times WV_{d,h,f}$$

$$AV_{b,h,p,f}^{B,H,P=M,F} = A_{b,h,m,p} \times V_{b,h,m,f}$$

$$C_{b,p,h \times F+f}^{B,P=M,G=H \times F} = AV_{b,h,p,f}$$

$$Z_{b,p,d} = C_{b,p,f} \times WZ_{g,d}$$

Separation of Concerns

Base transformer [Vaswani et al. 2016] SpAtten [Wang et al. 2020]

$$\begin{split} Q_{b,e,h,m} &= I_{b,d,m} \times WQ_{d,e,h} : \bigvee_{d} + (\cup) \\ K_{b,e,h,m} &= I_{b,d,m} \times WK_{d,e,h} : \bigvee_{d} + (\cup) \\ V_{b,f,h,m} &= I_{b,d,m} \times WV_{d,f,h} : \bigvee_{d} + (\cup) \\ QK_{b,h,m,p}^{B,H,M,P=M} &= \frac{1}{\sqrt{E}} \times Q_{b,e,h,p}^{B,E,H,M} \times K_{b,e,h,m} : \bigvee_{e} + (\cup) \\ SN_{b,h,m,p}^{B,H,M,P=M} &= e^{QK_{b,h,m,p}} \\ SD_{b,h,p}^{B,H,P=M} &= SN_{b,h,m,p} : \bigvee_{m} + (\cup) \\ A_{b,h,m,p} &= SN_{b,h,m,p} / SD_{b,h,p} \\ AV_{b,f,h,p} &= A_{b,h,m,p} \times V_{b,f,h,m} : \bigvee_{m} + (\cup) \\ C_{b,h \times F+f,p}^{B,G=H \times F,P=M} &= AV_{b,f,h,p} \\ Z_{b,d,p} &= C_{b,g,p} \times WZ_{d,g} : \bigvee_{g} + (\cup) \end{split}$$

$IP_{i,b,d,m} = MT_{i,b,m} \times IKV_{b,d,m}$ $Q_{i,b,e,h,r} = MH_{i,b,h} imes IQ_{b,d,r} imes WQ_{d,e,h}: igvee_d + (\cup)$ $K_{i,b,e,h,m} = MH_{i,b,h} imes IP_{b,d,m} imes WK_{d,e,h} : igvee_d + (\cup)$ $V_{i,b,f,h,m} = MH_{i,b,h} imes IP_{b,d,m} imes WV_{d,f,h} : igvee_{I} + (\cup)$ $QK_{i,b,h,m,r} = rac{1}{\sqrt{E}} imes Q_{i,b,e,h,r} imes K_{i,b,e,h,m} : igvee + (\cup)$ $SN_{i,b,h,m,r}=e^{QK_{i,b,h,m,r}}$ $SD_{i,b,h,r}=SN_{i,b,h,m,r}:\bigvee_m+(\cup)$ $A_{i,b,h,m,r} = SN_{i,b,h,m,r}/SD_{i,b,h,r}$ $AV_{i,b,f,h,r} = A_{i,b,h,m,r} imes V_{i,b,f,h,m}: igvee_m + (\cup)$ $C^{B,G=H imes F,R}_{i,b,hst F+f,r}=AV_{i,b,f,h,r}$ $Z_{i,b,d,r} = C_{i,b,g,r} imes WZ_{d,g} : \bigvee_{a} + (\cup)$ $ZA_{i,b,f,h,r} = Z_{i,b,h*F+f,r}$ $MT_{i+1,b,m} = prune(A_{i,b,h,m,r}) imes MT_{i,b,m}$ $MH_{i+1,b,h} = prune(ZA_{i,b,f,h,r}) imes MH_{i,b,h}$ $ST_{i+1,b,m} = ST_{i,b,m} + A_{i,b,h,m,r}$

 $SH_{i+1,b,h} = SH_{i,b,h} + ZA_{i,b,f,h,r}$

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Base transformer [Vaswani et al. 2016]

$$\begin{split} Q_{b,e,h,m} &= I_{b,d,m} \times WQ_{d,e,h} : \bigvee_d + (\cup) \\ K_{b,e,h,m} &= I_{b,d,m} \times WK_{d,e,h} : \bigvee_d + (\cup) \\ V_{b,f,h,m} &= I_{b,d,m} \times WV_{d,f,h} : \bigvee_d + (\cup) \\ QK_{b,h,m,p}^{B,H,M,P=M} &= \frac{1}{\sqrt{E}} \times Q_{b,e,h,p}^{B,E,H,M} \times K_{b,e,h,m} : \bigvee_e + (\cup) \\ SN_{b,h,m,p}^{B,H,M,P=M} &= e^{QK_{b,h,m,p}} \\ SD_{b,h,p}^{B,H,P=M} &= SN_{b,h,m,p} : \bigvee_m + (\cup) \\ A_{b,h,m,p} &= SN_{b,h,m,p} / SD_{b,h,p} \\ AV_{b,f,h,p} &= A_{b,h,m,p} \times V_{b,f,h,m} : \bigvee_m + (\cup) \\ C_{b,h \times F+f,p}^{B,G=H \times F,P=M} &= AV_{b,f,h,p} \\ Z_{b,d,p} &= C_{b,g,p} \times WZ_{d,g} : \bigvee_g + (\cup) \end{split}$$

Sanger [Lu et al. 2021]

 $Q_{b,e,h,m} = I_{b,d,m} imes WQ_{d,e,h} : igvee_{J} + (\cup)$ $K_{b,e,h,m} = I_{b,d,m} imes WK_{d,e,h} : igvee_d + (\cup)$ $V_{b,f,h,m} = I_{b,d,m} imes WV_{d,f,h} : igvee_d + (\cup)$ $\hat{Q}_{b,e,h,m} = Qt(Q_{b,e,h,m})$ $\hat{K}_{b,e,h,m} = Qt(K_{b,e,h,m})$ $\hat{QK}^{B,H,M,P=M}_{b,h,m,p} = rac{1}{\sqrt{E}} imes \hat{Q}^{B,E,H,M}_{b,e,h,p} imes \hat{K}_{b,e,h,m}: igvee_{a} + (\cup)$ $\hat{SN}_{b,h,m,p}=e^{\hat{QK}_{b,h,m,p}}$ $\hat{SD}_{b,h,p} = \hat{SN}_{b,h,m,p}: \bigvee_m + (\cup)$ $\hat{A}_{b,h,m,p} = \hat{SN}_{b,h,m,p}/\hat{SD}_{b,h,p}$ $M_{b,h,m,p} = prune(\hat{A}_{b,h,m,p})$ $QK^{B,H,M,P=M}_{b,h,m,p} = rac{1}{\sqrt{E}} imes M_{b,h,m,p} imes Q^{B,E,H,M}_{b,e,h,p} imes K_{b,e,h,m} : igvee + (\cup)$ $SN^{B,H,M,P=M}_{b\ b\ m\ n}=e^{QK_{b,h,m,p}}$ $SD^{B,H,P=M}_{b,h,p}=SN_{b,h,m,p}:\bigvee_m+(\cup)$ $A^{B,H,M,P=M}_{b,b,m,p} = SN_{b,h,m,p}/SD_{b,h,p}$ $AV^{B,F,H,P=M}_{b,f,h,p} = A_{b,h,m,p} imes V_{b,f,h,m}: \bigvee_m + (\cup)$ $C^{B,G=H imes F,P=M}_{b,h*F+f,p}=AV_{b,f,h,p}$ $Z^{B,D,P=M}_{b,d,p} = C_{b,g,p} imes WZ_{d,g}: igvee_{a} + (\cup)$ The second second

Base transformer [Vaswani et al. 2016]

$$\begin{split} Q_{b,e,h,m} &= I_{b,d,m} \times WQ_{d,e,h} : \bigvee_d + (\cup) \\ K_{b,e,h,m} &= I_{b,d,m} \times WK_{d,e,h} : \bigvee_d + (\cup) \\ V_{b,f,h,m} &= I_{b,d,m} \times WV_{d,f,h} : \bigvee_d + (\cup) \\ QK_{b,h,m,p}^{B,H,M,P=M} &= \frac{1}{\sqrt{E}} \times Q_{b,e,h,p}^{B,E,H,M} \times K_{b,e,h,m} : \bigvee_e + (\cup) \\ SN_{b,h,m,p}^{B,H,M,P=M} &= e^{QK_{b,h,m,p}} \\ SD_{b,h,p}^{B,H,P=M} &= SN_{b,h,m,p} : \bigvee_m + (\cup) \\ A_{b,h,m,p} &= SN_{b,h,m,p} / SD_{b,h,p} \\ AV_{b,f,h,p} &= A_{b,h,m,p} \times V_{b,f,h,m} : \bigvee_m + (\cup) \\ C_{b,h \times F+f,p}^{B,G=H \times F,P=M} &= AV_{b,f,h,p} \\ Z_{b,d,p} &= C_{b,g,p} \times WZ_{d,g} : \bigvee_g + (\cup) \end{split}$$

EdgeBERT [Tambe et al. 2021]

$$\begin{split} MH_{h} &= take(M_{h,m,p}, zeroes(m,p), 0) \\ Q_{b,e,h,m} &= MH_{h} \times I_{b,d,m} \times WQ_{d,e,h} : \bigvee_{d} + (\cup) \\ K_{b,e,h,m} &= MH_{h} \times I_{b,d,m} \times WK_{d,e,h} : \bigvee_{d} + (\cup) \\ V_{b,f,h,m} &= MH_{h} \times I_{b,d,m} \times WV_{d,f,h} : \bigvee_{d} + (\cup) \\ QK_{b,h,m,p}^{B,H,M,P=M} &= \frac{1}{\sqrt{E}} \times Q_{b,e,h,p}^{B,E,H,M} \times K_{b,e,h,m} : \bigvee_{e} + (\cup) \\ SN_{b,h,m,p}^{B,H,M,P=M} &= e^{QK_{b,h,m,p}} \\ SD_{b,h,p}^{B,H,P=M} &= SN_{b,h,m,p} : \bigvee_{m} + (\cup) \\ A_{b,h,m,p} &= SN_{b,h,m,p} \times V_{b,f,h,m} : \bigvee_{m} + (\cup) \\ C_{b,h*F+f,p}^{B,G=H \times F,P=M} &= AV_{b,f,h,p} \\ Z_{b,d,p}^{B,D,P=M} &= C_{b,g,p} \times WZ_{d,g} : \bigvee_{g} + (\cup) \end{split}$$

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Base transformer [Vaswani et al. 2016]

$$\begin{split} Q_{b,e,h,m} &= I_{b,d,m} \times WQ_{d,e,h} : \bigvee_{d} + (\cup) \\ K_{b,e,h,m} &= I_{b,d,m} \times WK_{d,e,h} : \bigvee_{d} + (\cup) \\ V_{b,f,h,m} &= I_{b,d,m} \times WV_{d,f,h} : \bigvee_{d} + (\cup) \\ QK_{b,h,m,p}^{B,H,M,P=M} &= \frac{1}{\sqrt{E}} \times Q_{b,e,h,p}^{B,E,H,M} \times K_{b,e,h,m} : \bigvee_{e} + (\cup) \\ SN_{b,h,m,p}^{B,H,M,P=M} &= e^{QK_{b,h,m,p}} \\ SD_{b,h,p}^{B,H,P=M} &= SN_{b,h,m,p} : \bigvee_{m} + (\cup) \\ A_{b,h,m,p} &= SN_{b,h,m,p} / SD_{b,h,p} \\ AV_{b,f,h,p} &= A_{b,h,m,p} \times V_{b,f,h,m} : \bigvee_{m} + (\cup) \\ C_{b,h \times F+f,p}^{B,G=H \times F,P=M} &= AV_{b,f,h,p} \\ Z_{b,d,p} &= C_{b,g,p} \times WZ_{d,g} : \bigvee_{q} + (\cup) \end{split}$$

DOTA [Qu et al. 2022]

$$\begin{aligned} Q_{b,e,h,m} &= I_{b,d,m} \times WQ_{d,e,h} : \bigvee_{d} + (\cup) \\ K_{b,e,h,m} &= I_{b,d,m} \times WK_{d,e,h} : \bigvee_{d} + (\cup) \\ V_{b,f,h,m} &= I_{b,d,m} \times WV_{d,f,h} : \bigvee_{d} + (\cup) \\ \tilde{Q}_{b,k,h,m} &= I_{b,d,m} \times P_{d,j}^{D,K} \times \tilde{WQ}_{j,k,h} : \bigvee_{d,j} + (\cup) \\ \tilde{K}_{b,k,h,m} &= I_{b,d,m} \times P_{d,j}^{D,K} \times \tilde{WK}_{j,k,h} : \bigvee_{d,j} + (\cup) \\ \tilde{K}_{b,k,h,m} &= I_{b,d,m} \times P_{d,j}^{D,K} \times \tilde{WK}_{j,k,h} : \bigvee_{d,j} + (\cup) \\ Q\tilde{K}_{b,h,m,p}^{B,H,M,P=M} &= \tilde{Q}_{b,k,h,p}^{B,K,H,M} \times \tilde{K}_{b,k,h,m} : \bigvee_{k} + (\cup) \\ M_{b,h,m,p} &= prune(Q\tilde{K}_{b,h,m,p}) \\ QK_{b,h,m,p}^{B,H,M,P=M} &= \frac{1}{\sqrt{E}} \times M_{b,h,m,p} \times Q_{b,e,h,p}^{B,E,H,M} \times K_{b,e,h,m} : \bigvee_{e} + (\cup) \\ SN_{b,h,m,p}^{B,H,M,P=M} &= e^{QK_{b,h,m,p}} \\ SD_{b,h,p}^{B,H,M,P=M} &= SN_{b,h,m,p} : \bigvee_{m} + (\cup) \\ A_{b,h,m,p}^{B,H,M,P=M} &= SN_{b,h,m,p} / SD_{b,h,p} \\ AV_{b,f,h,p}^{B,F,H,P=M} &= A_{b,h,m,p} \times V_{b,f,h,m} : \bigvee_{m} + (\cup) \\ C_{b,h*F+f,p}^{B,G=H \times F,P=M} &= AV_{b,f,h,p} \\ Z_{b,d,p}^{B,D,P=M} &= C_{b,g,p} \times WZ_{d,g} : \bigvee_{g} + (\cup) \end{aligned}$$

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Einsums: Precise and Concise

	Paper Length	Code length	# Einsums
Attention Is All You Need	15 pages	14 python files	14
FlashAttention	34 pages	17 pythop filos	24 (3 changed)
FlashAttention2	14 pages	47 python files	25 (11 changed)
Spatten [Wang]	15 pages	8 python files + CPP	19 (3 changed)
Sanger [Lu]	15 pages	16 python files + Scala	21 (1 changed)
EdgeBert [Tambe]	16 pages	80+ python files	15 (4 changed)
DOTA [Qu]	13 pages	N/A	17 (1 changed)

Separation of Concerns

The High Cost of Data Movement

Fetching operands more expensive than computing on them



Spatial Architecture for DNN



Types of Data Reuse in DNN

Convolutional Reuse

CONV layers only (sliding window)



Reuse: Activations Filter weights

Types of Data Reuse in DNN



Types of Data Reuse in DNN



















1-D Convolution - Movie

		Т	enso	or:	F[S]	
		R	ank:	: S			
			0	1	2		
		8	5	5 2			
Ter	nsor	: 1	[W]				
Rar	nk: I	W					
0	1	2	3	4	5	6	7
1	1	2	3	3	2	7	6
	Ten	sor	: 0	[Q]			
	Ran	k: (2				
	0	1	2	3	4	5	
	0	0	0	0	0	0	

Output Stationary – Spacetime View



	P	ank	: т															
		0	1	2	3 4	5	6	7	8	9	10	11	12	13	14	15	16	17
Rank: W	0-1	-	1 1	1	1							1						1
	1 1	-	1 1	1	- 1	1	1	1	1	1		1		1	1	1	1	1
	2 2	2	2 -2	2	2	2	-2	2	2									
	3 3		3 3	3	3	-3	3	3		3	3	3	3	3	3	3	3	3
	4 3		3 3	3		3	3		3		3		3	3	3	3	3	
	5 2		2 2	2	2	2	2	2				2		2	2	-2	2	2
	6 7		7 7	7											7	7	7	7
	7 6		6 6	6	6	6	6	6	6	6	6	6	6	6	6	6		6



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1-D Output Stationary



Output Stationary (OS)



- Minimize partial sum R/W energy consumption
 - maximize local accumulation
- Broadcast/Multicast filter weights and reuse activations spatially across the PE array

OS Example: ShiDianNao



• Partial sums accumulated in PE and streamed out

[Du et al., ISCA 2015]

OS Example: KU Leuven

16x16b / 1x16b



[Moons et al., VLSI 2016, ISSCC 2017]

Many Dataflows

• Output Stationary (OS)

[Peemen, *ICCD* 2013] [ShiDianNao, *ISCA* 2015] [Gupta, *ICML* 2015] [Moons, *VLSI* 2016] [Thinker, *VLSI* 2017]

• Weight Stationary (WS)

[Chakradhar, /SCA 2010] [nn-X (NeuFlow), CVPRW 2014] [Park, /SSCC 2015] [ISAAC, /SCA 2016] [PRIME, /SCA 2016] [TPU, /SCA 2017]

• Input Stationary (IS)

[Parashar (SCNN), ISCA 2017]

• Row Stationary (IS)

[Eyeriss, ISCA 2016] [Tetris ASPLOS 2017] [Eyeriss2, JETCAT 2019]

Many Mapping Options

Per storage level cross product of:

- Dataflow
- Tiling in time
- Tiling in space
- Bypassing
- Split/Shared storage

Variation in Traffic with Mapping





Williams, Samuel, Andrew Waterman, and David Patterson. "Roofline: an insightful visual performance model for multicore architectures." Communications of the ACM 52.4 (2009): 65-76.















Thank you!

Next Lecture: Accelerators (II)