### 6.823 Spring 2014

## Handout \#14 - Multi-producer/Single-consumer Messaging Queues <br> http://csg.csail.mit.edu/6.823/

You are writing a queue to be used in a multi-producer/single-consumer application. (Producer threads write messages that are read by one consumer.) We assume here a queue with infinite space.

TST rs, $\operatorname{Imm}(r t)$ is the test-and-set instruction, which atomically loads the value at $\operatorname{Imm}(r t)$ into $r s$, and if the value is zero, updates the memory location at $\operatorname{Imm}(r t)$ to 1. This atomic instruction is useful for implementing locks: a value of 1 at the memory location indicates that someone holds the lock, and a value of 0 means the lock is free.

Producer pushes a message onto queue: (memory operations in bold)

```
void push(int** tail_ptr, int* tail_write_lock, int message) {
    while (lock_try(tail_write_lock) == false);
    **tail_ptr = message;
    *tail_ptr++;
    lock_\overline{release(tail_write_lock);}
}
# R1 - contains address of data to enqueue
# R2 - contains the address of the tail pointer of queue
# R3 - address of tail pointer write lock
P1 SpinLock:TST R4, 0(R3) # try to acquire tail write lock
P2 BNEZ R4, R4, SpinLock
P3 LD R4, 0(R2) # get tail pointer
P4 ST R1, 0(R4) # write message to tail
P5 ADD R4, R4, 4 # update tail pointer
P6 ST R4, 0(R2)
P7 ST R0, 0(R3) # release lock
```

Consumer pops a message off queue: (memory operations in bold)

```
int pop(int** head_ptr, int** tail_ptr) {
    while (*head_ptr == *tail_pt\overline{r});
    int message = **head_ptr;
    *head_ptr++;
    return message;
}
# R1 - will receive address contained in message
# R2 - contains the address of the head pointer of queue
# R3 - contains the address of the tail pointer of the queue
C1 Retry: LD R4, 0(R2) # get head pointer
C2 LD R5, 0(R3) # get tail pointer
C3 SUB R5, R4, R5 # is there a message?
C4 BNEZ R5, Pop
C5 JMP Retry
C6 Pop: LD R1, 0(R4) # read message from queue
C7 ADD R4, R4, 4 # update head pointer
C8 ST R4, 0(R2)
```

