Cache Coherence

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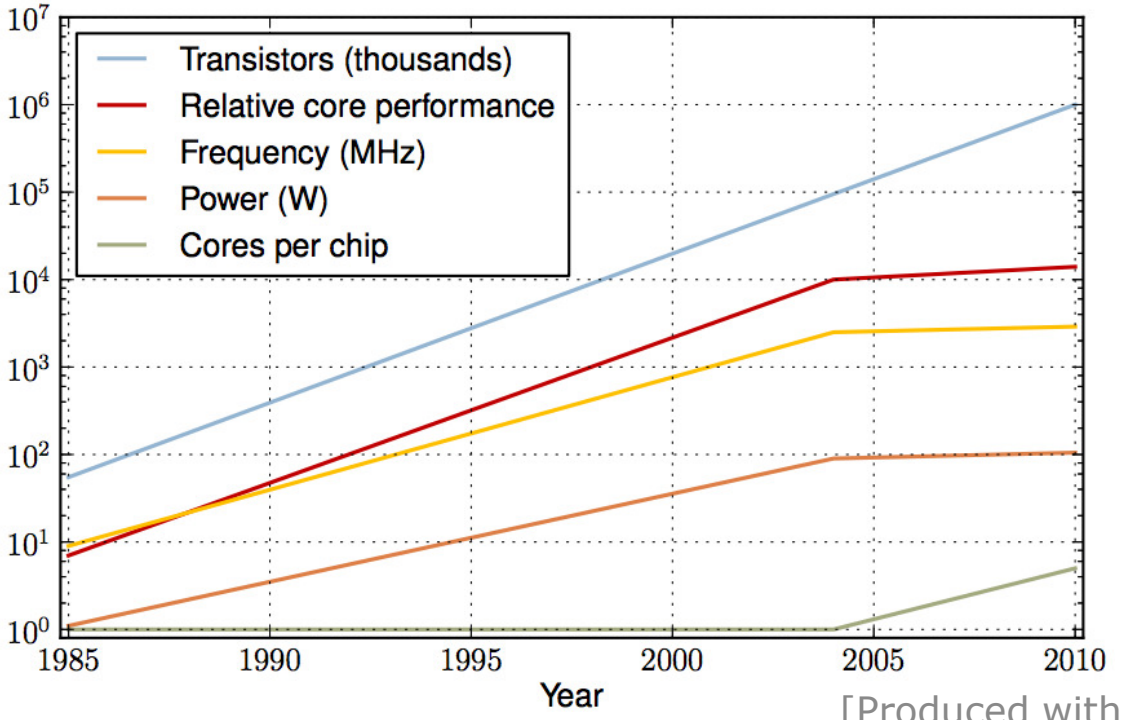
http://www.csg.csail.mit.edu/6.823

The Shift to Multicore



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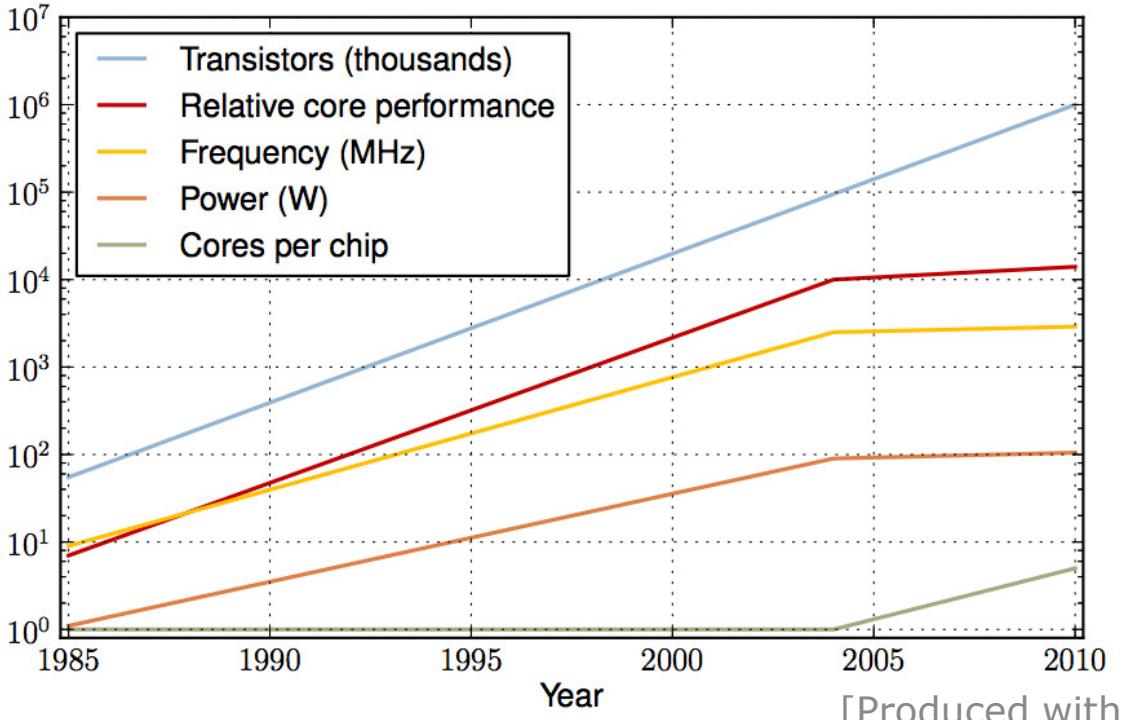

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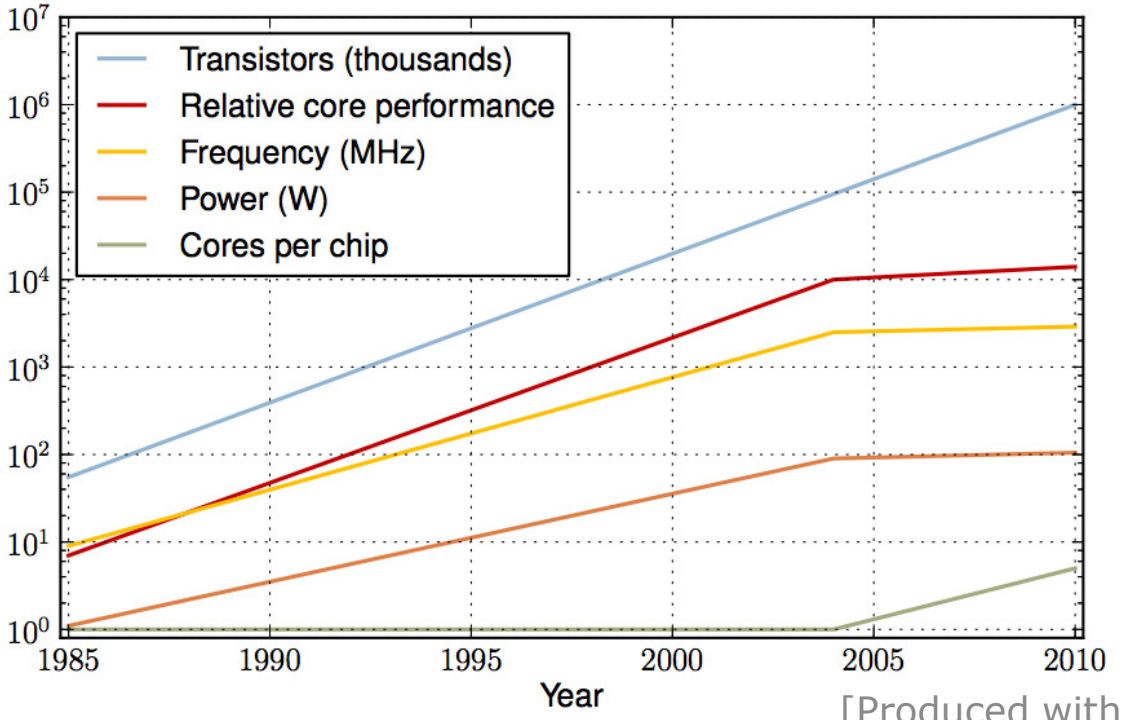

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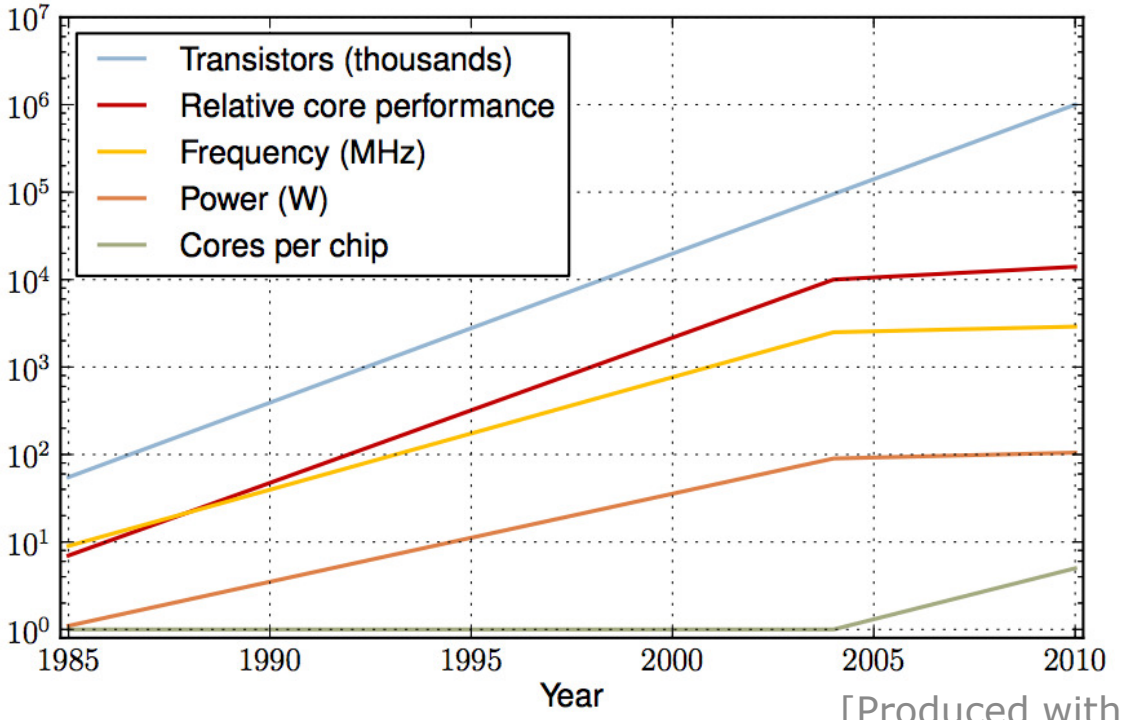

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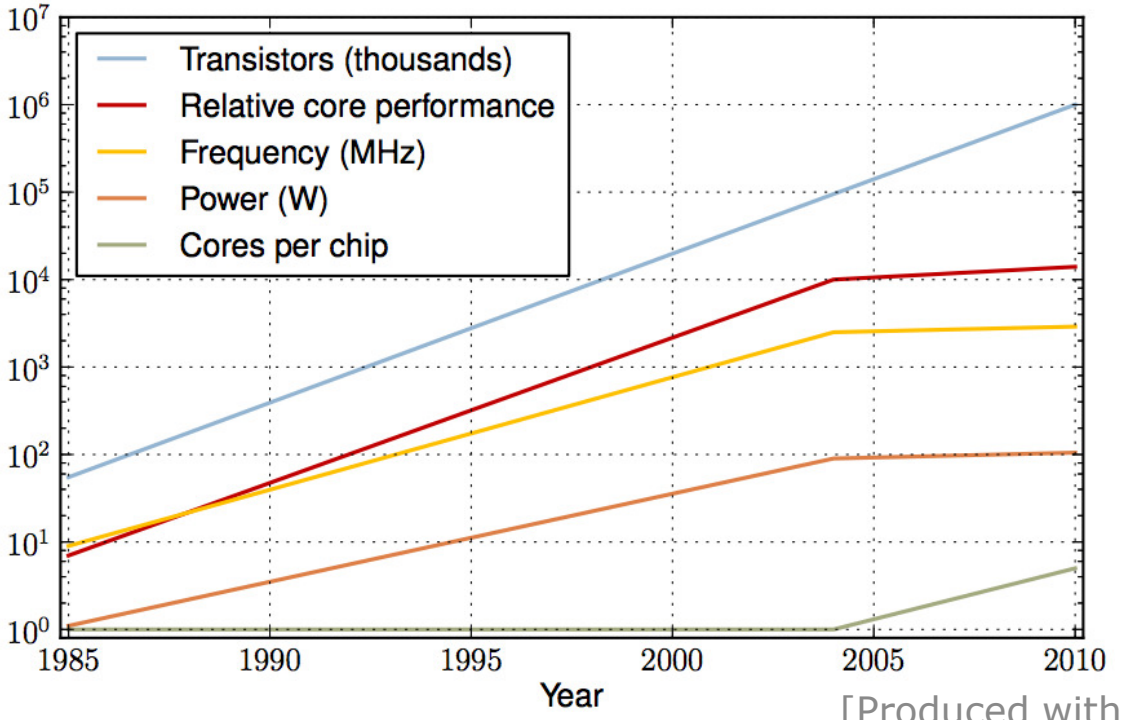

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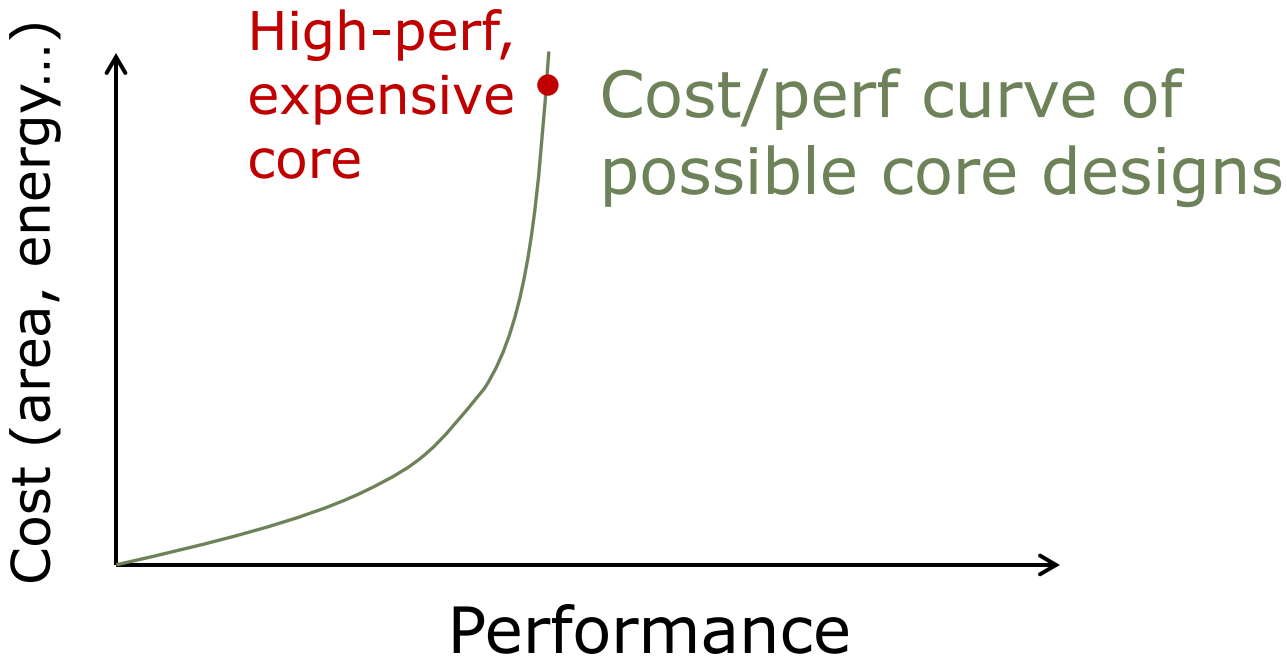
  Technology scaling

Multicore Performance



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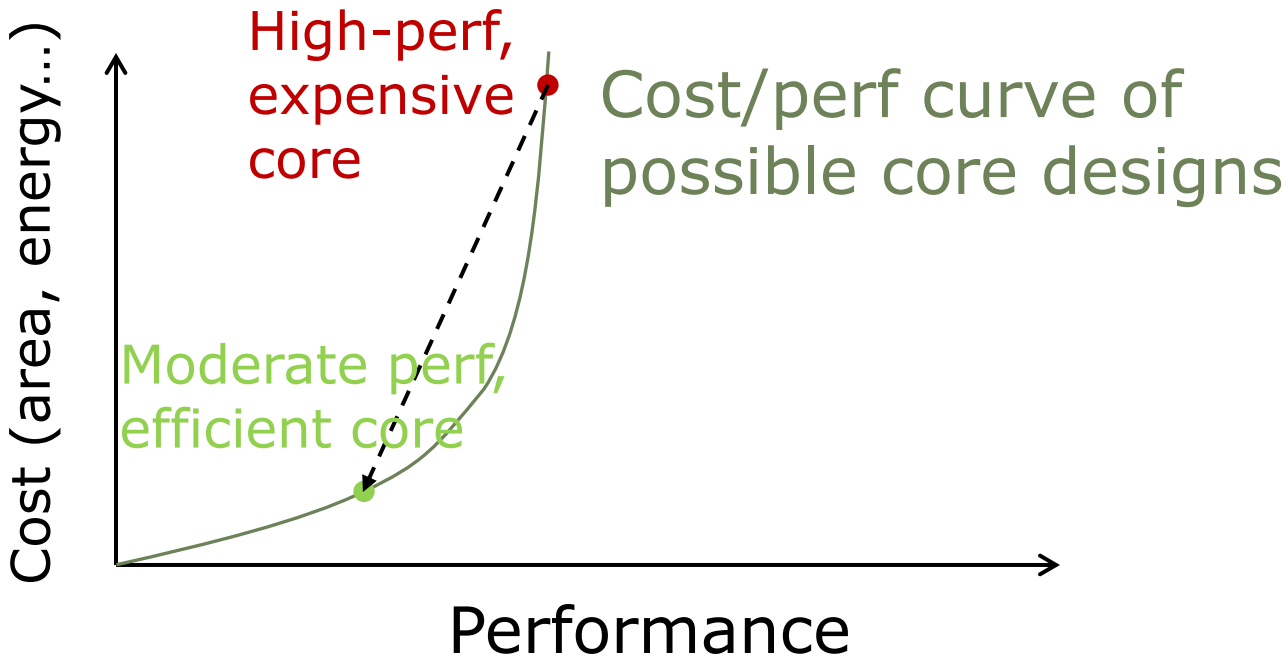

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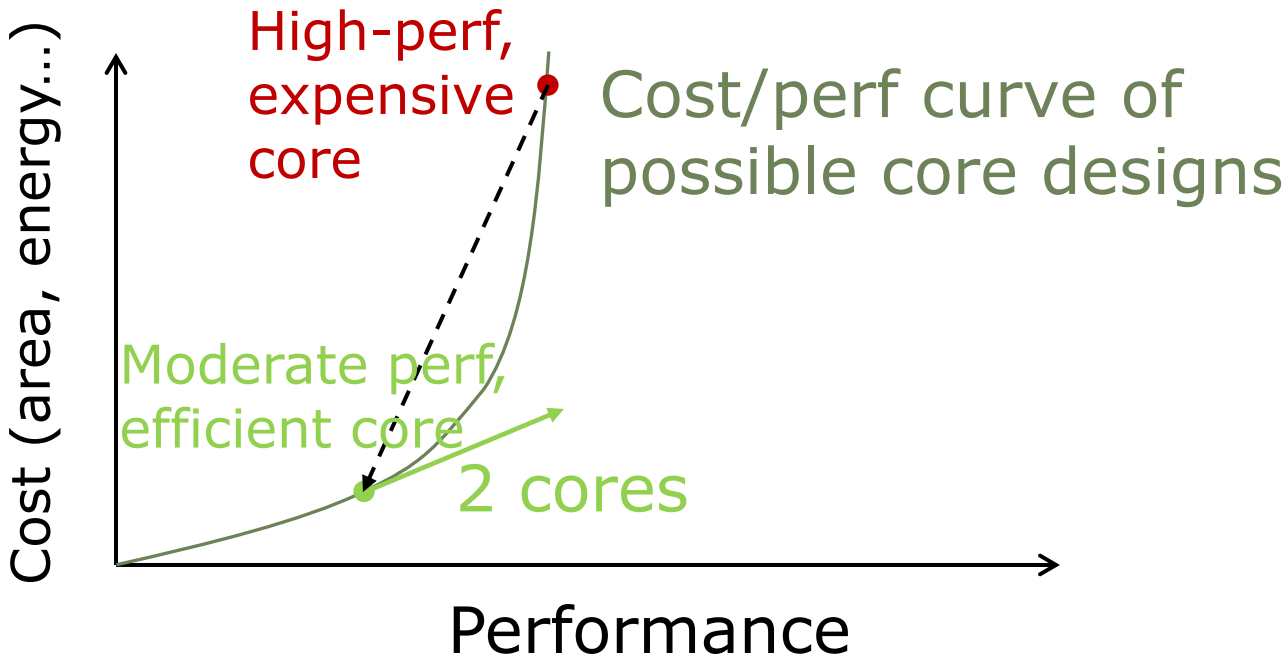

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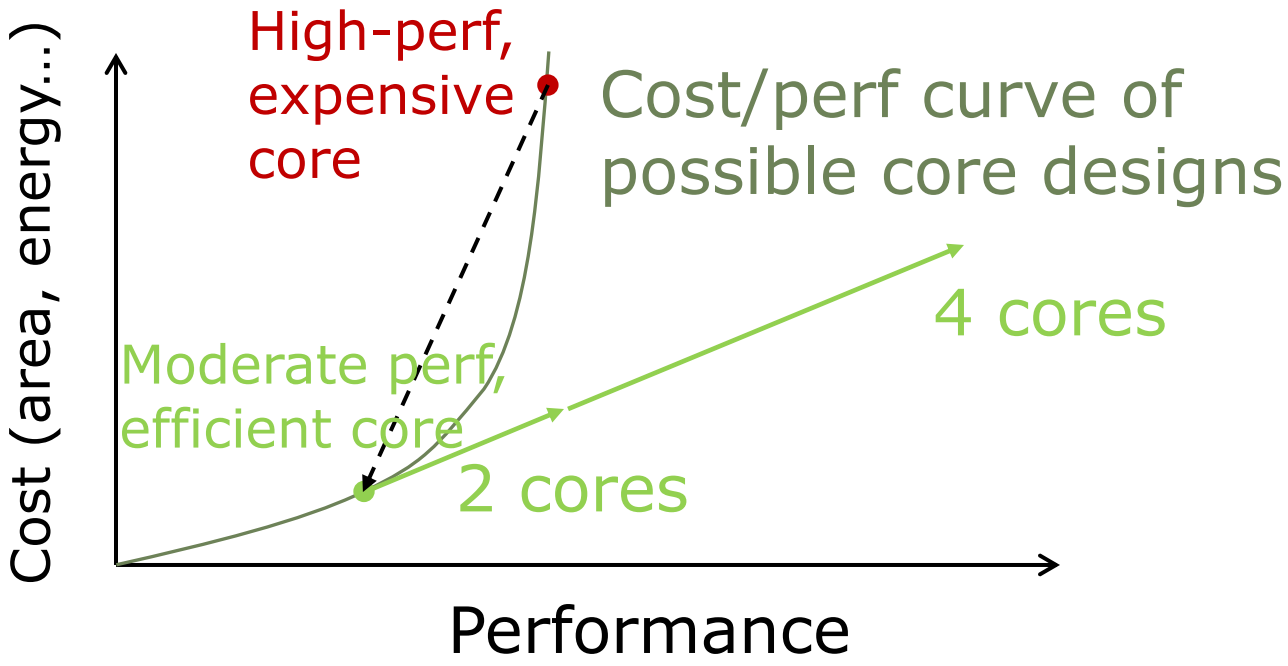

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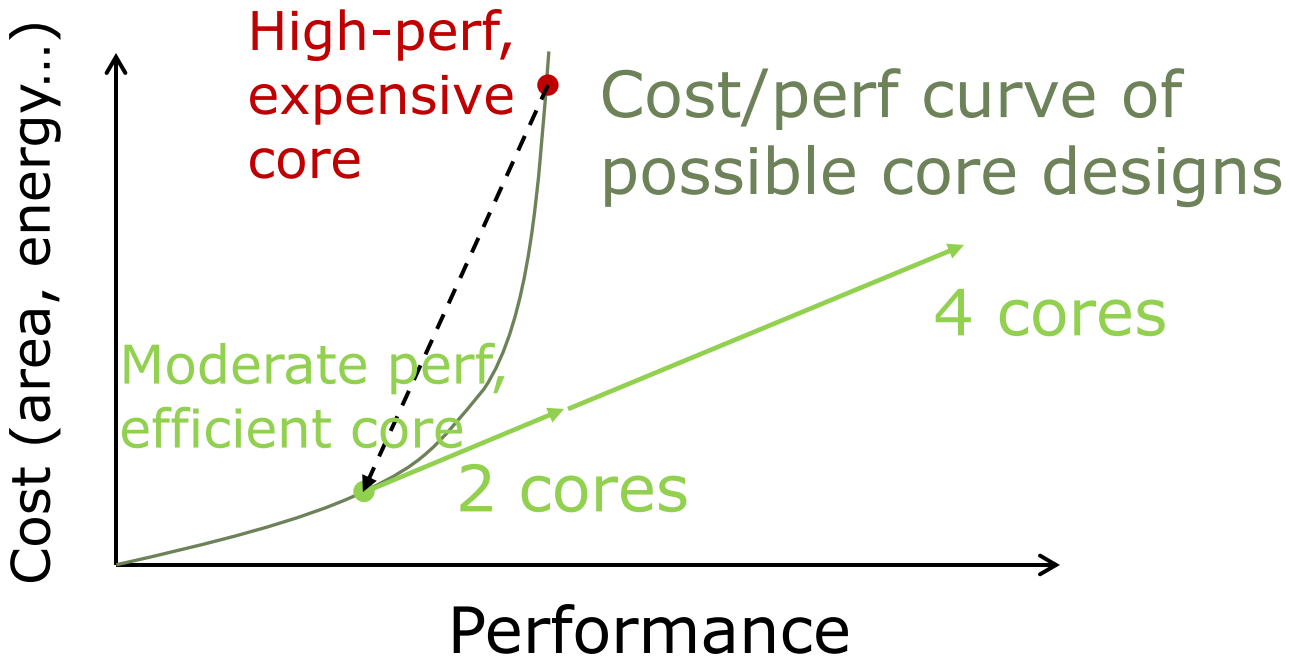
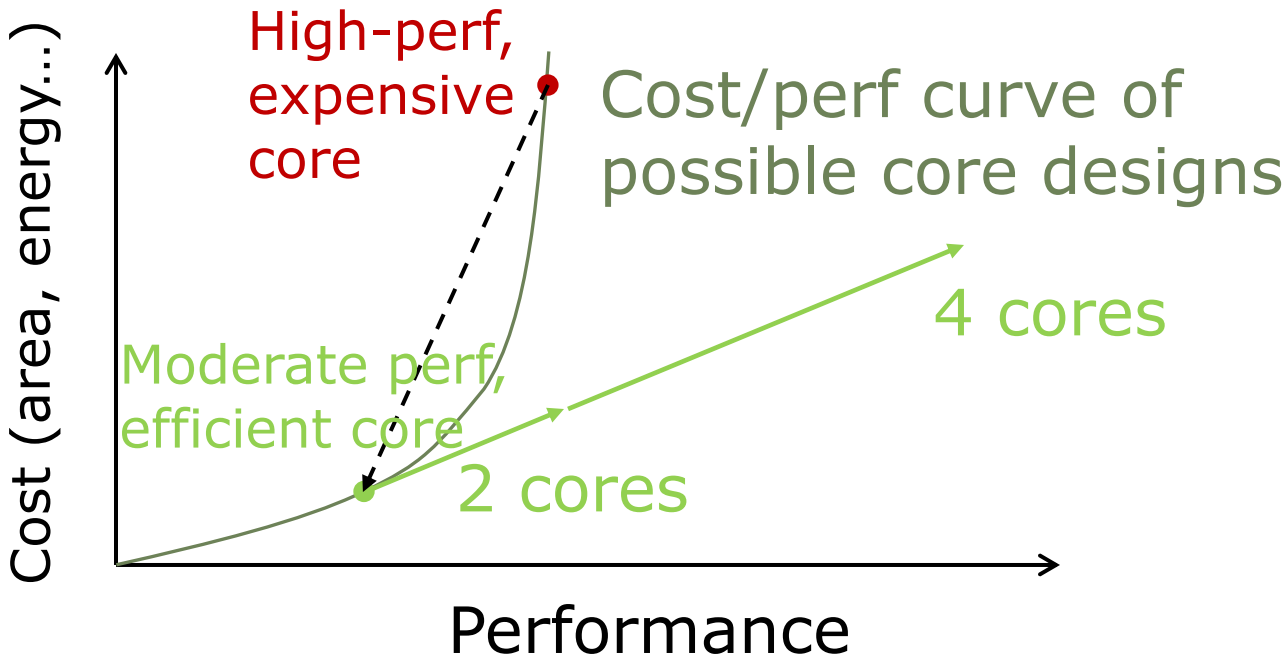


What factors may limit multicore performance?

Multicore Performance



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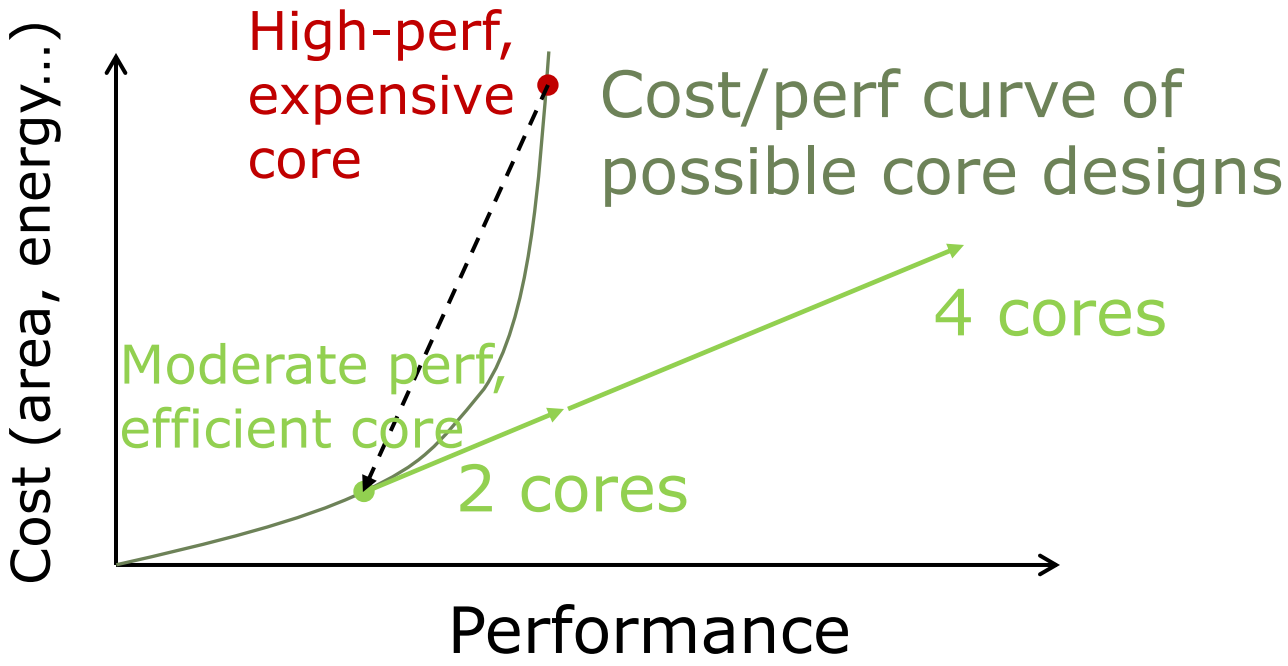

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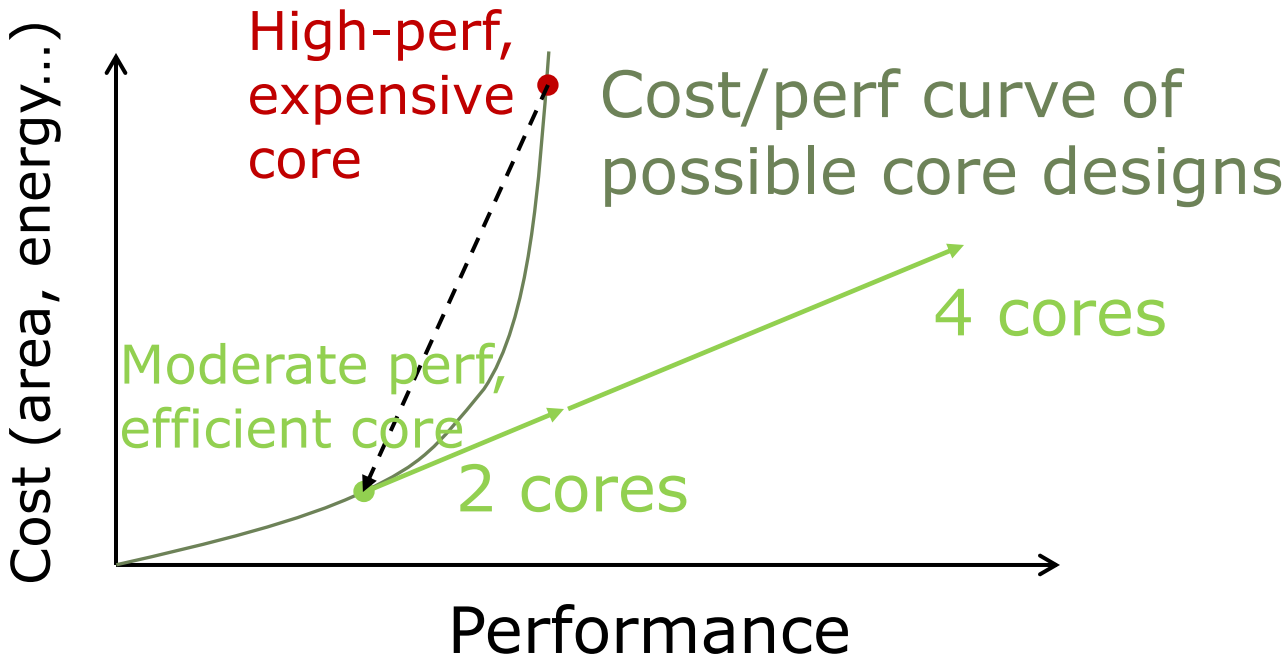

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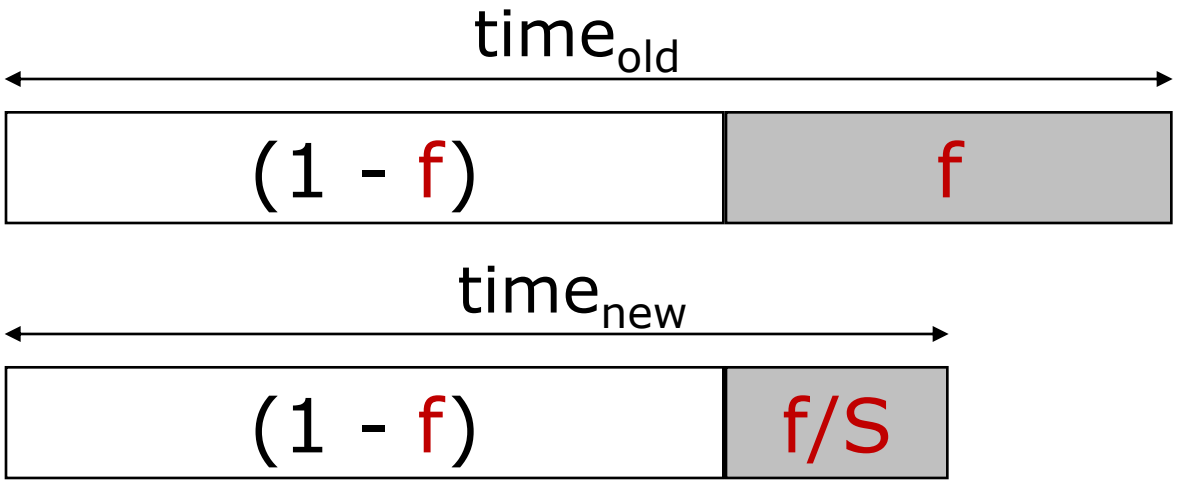
Memory accesses and inter-core communication
Programming complexity

Amdahl's Law

- Speedup= $time_{without\ enhancement}$ / $time_{with\ enhancement}$
- Suppose an enhancement speeds up a fraction f of a task by a factor of S

$$time_{new} = time_{old} \cdot ((1-f) + f/S)$$

$$S_{overall} = 1 / ((1-f) + f/S)$$



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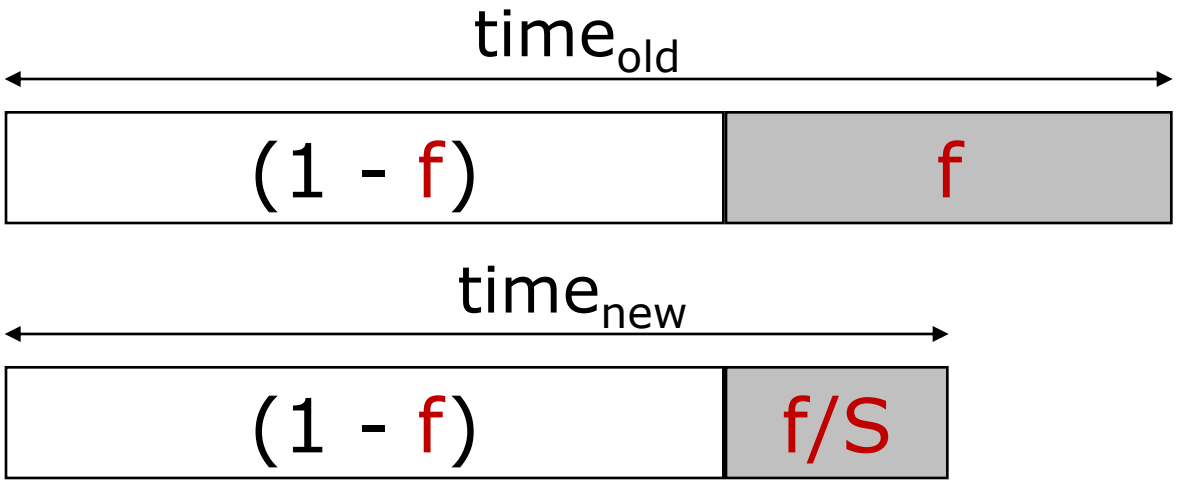

Corollary: Make the common case fast

Amdahl's Law and Parallelism

- Say you write a program that can do 90% of the work in parallel, but the other 10% is sequential
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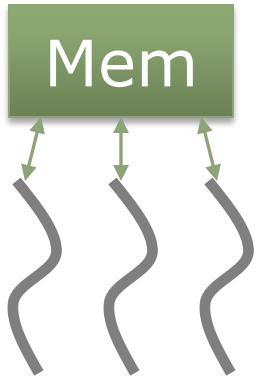
What f do you need to use a 1000-core machine well?

Communication Models

- Shared memory:
  - Single address space
  - Implicit communication by reading/writing memory
    - Data
    - Control (semaphores, locks, barriers, …)
  - Low-level programming model: threads



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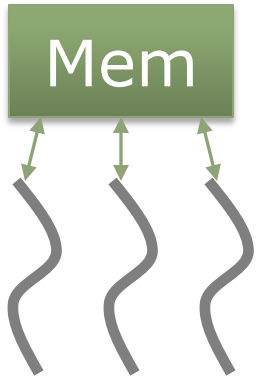

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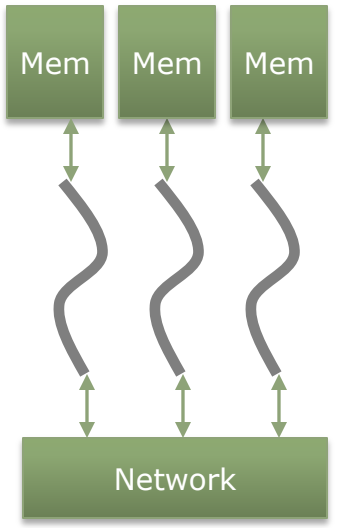

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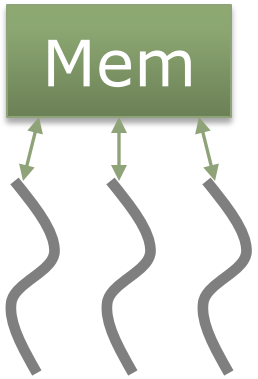

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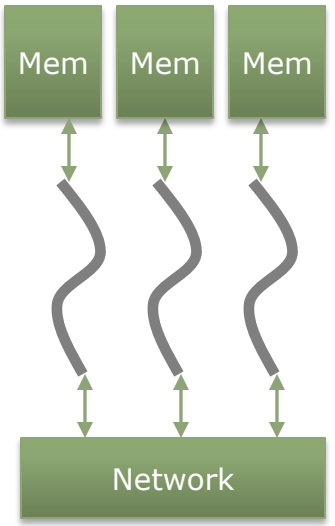

- Pros/cons of each model?

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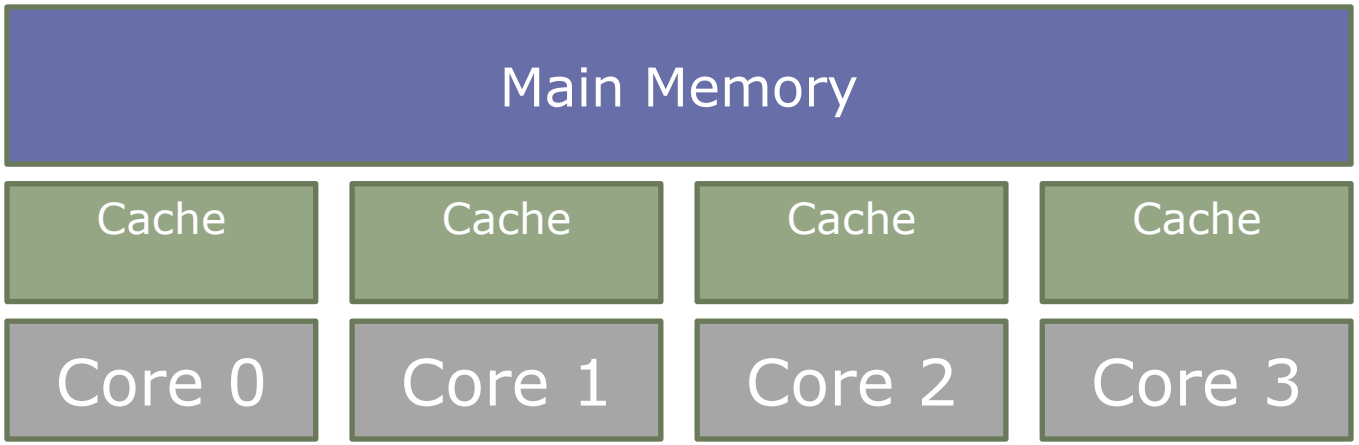

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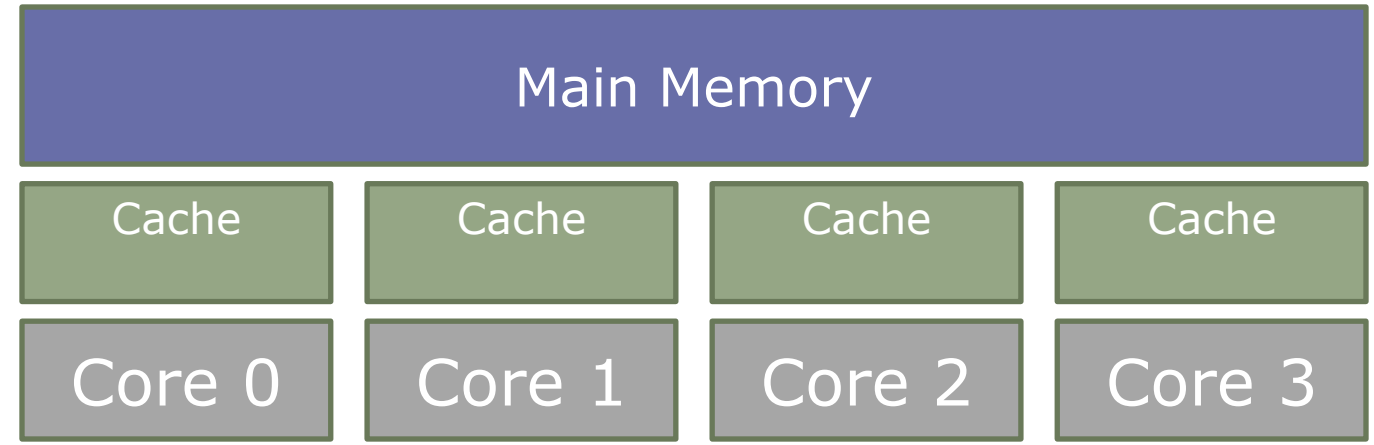

1. LD 0xA → 2

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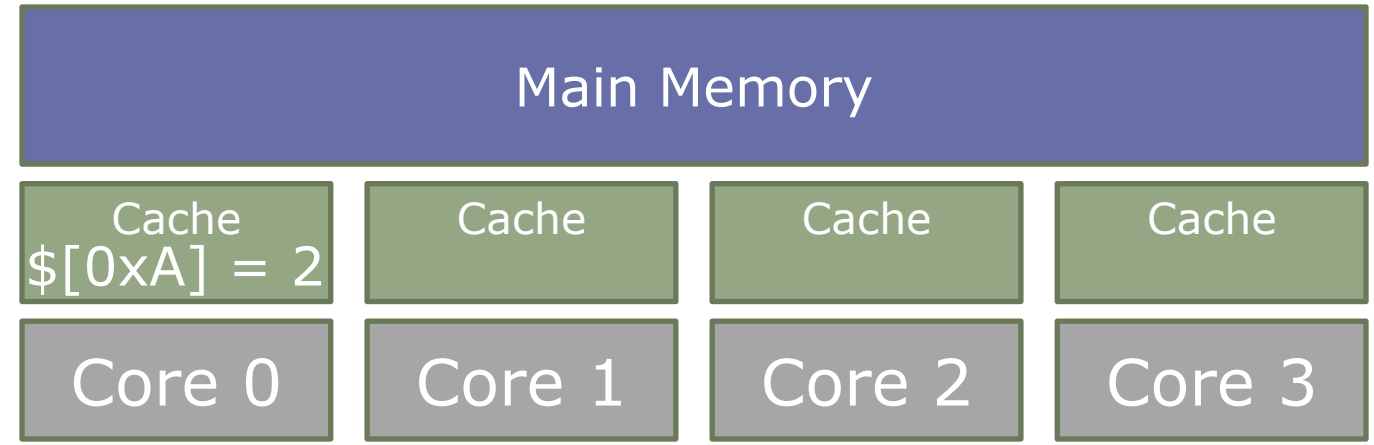

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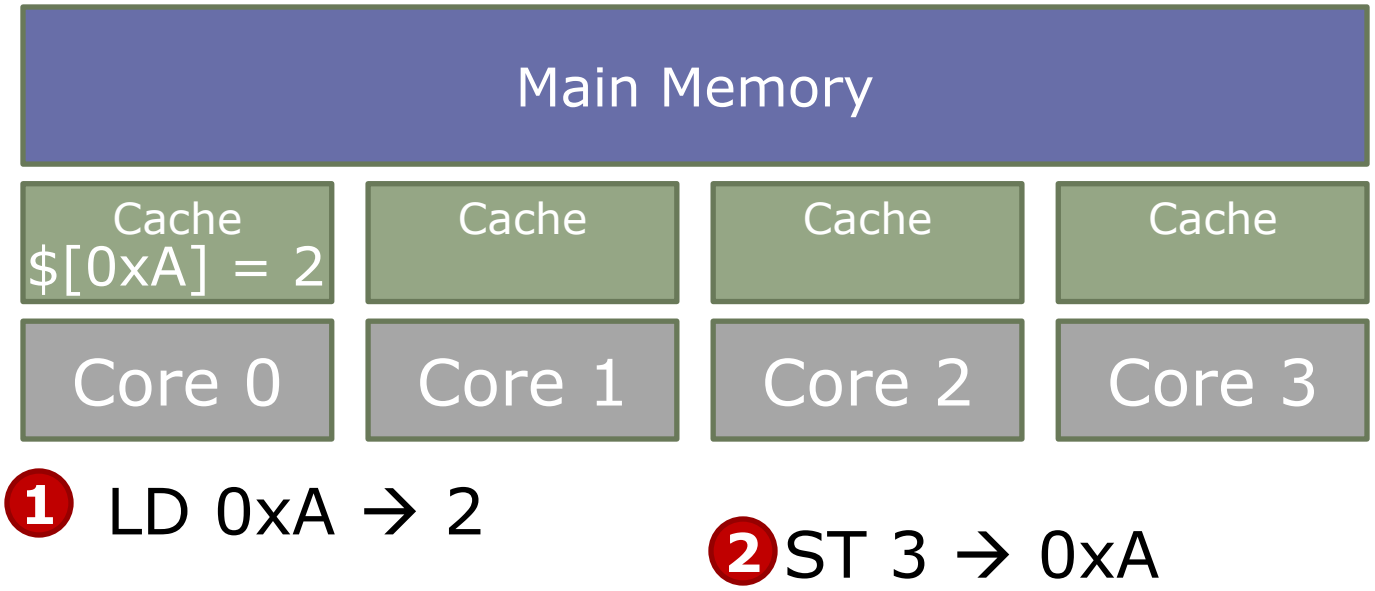

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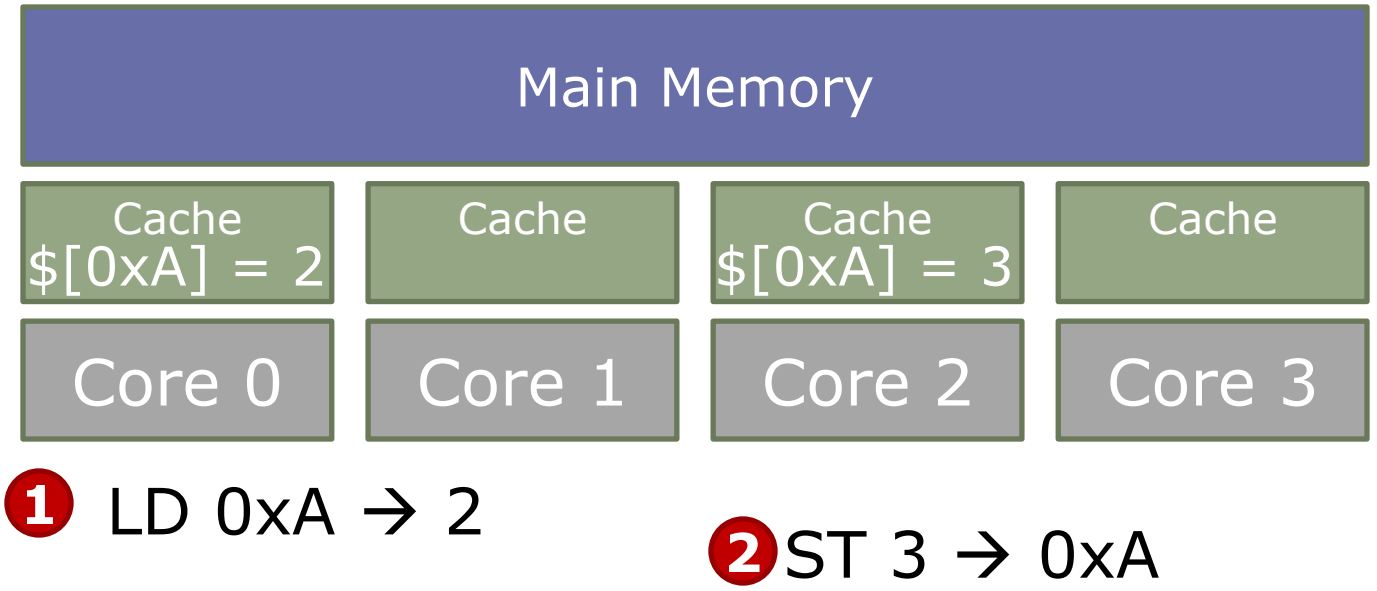

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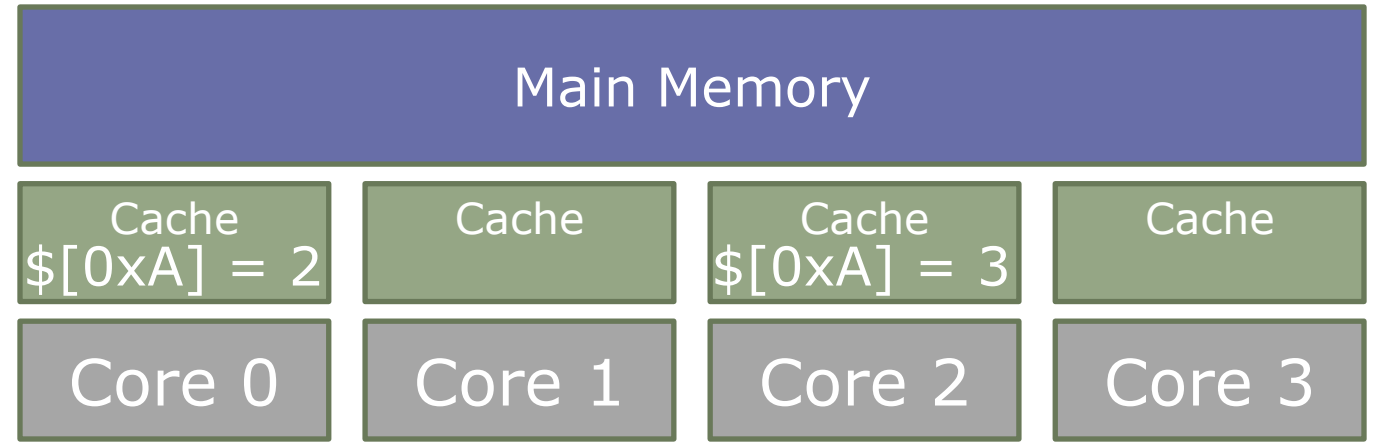
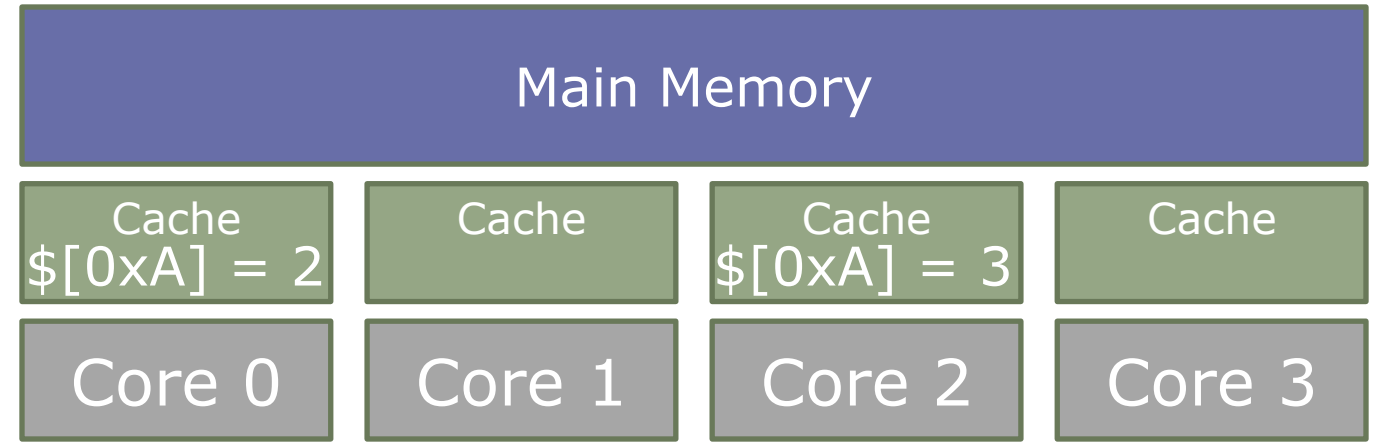


1. LD 0xA → 2
2. ST 3 → 0xA
3. LD 0xA → 2 (stale!)

Cache Coherence Avoids Stale Data



- A cache coherence protocol controls cache contents to avoid stale cache lines

Implementing Cache Coherence

- Coherence protocols must enforce two rules:
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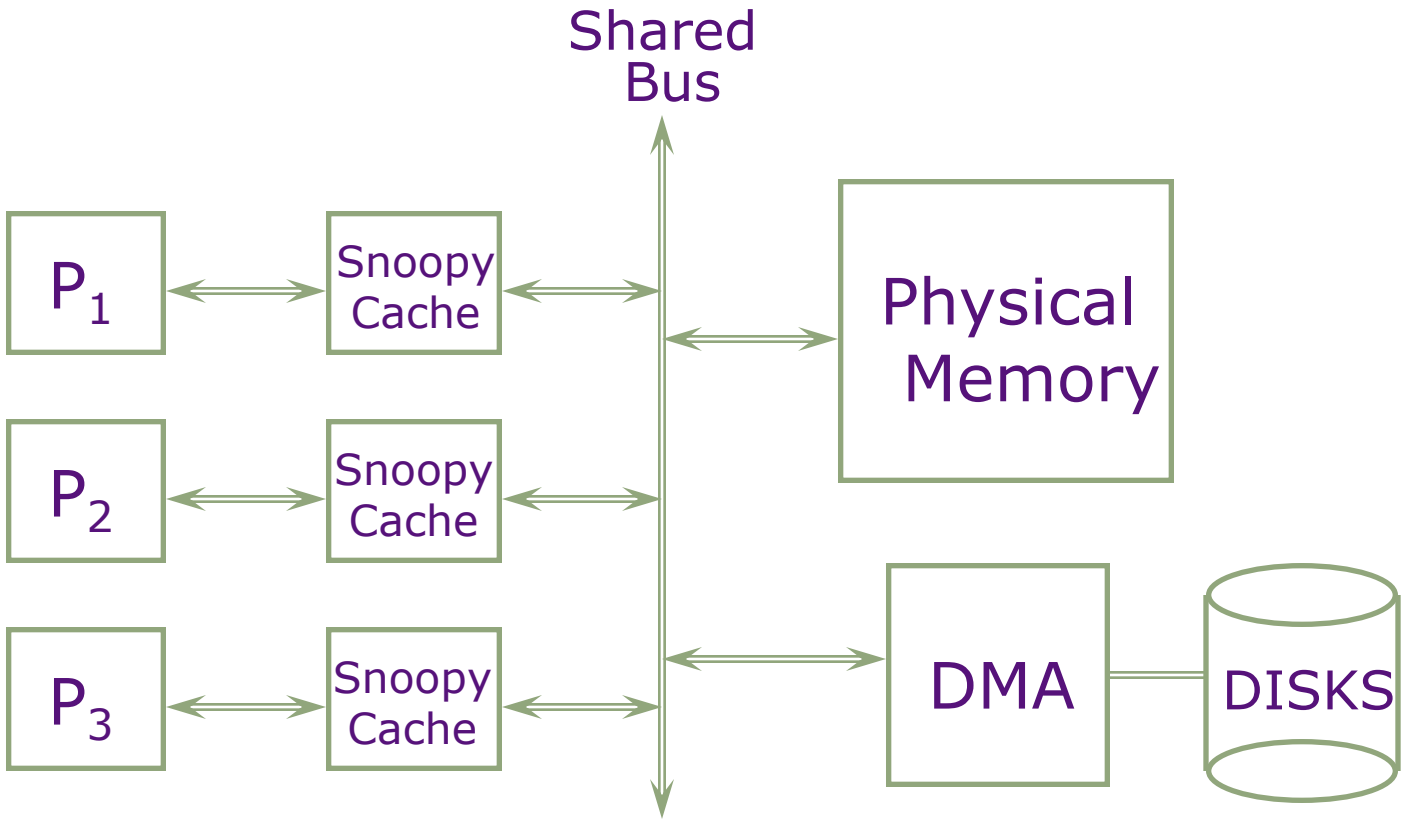
- How to track sharing state of cached data and serialize requests to the same address?
  - Snooping-based protocols: All caches observe each other's actions through a shared bus
  - Directory-based protocols: A coherence directory tracks contents of private caches and serializes requests

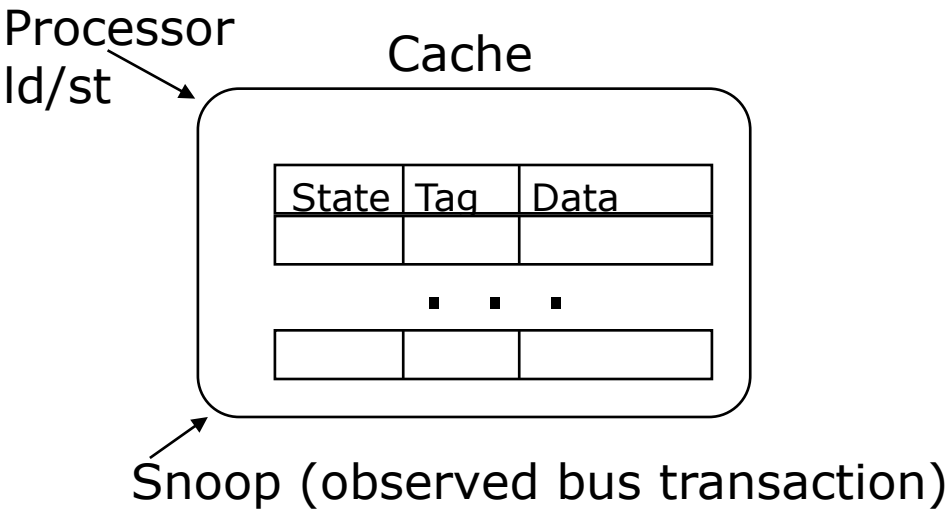
Snooping-Based Coherence [Goodman 1983]



Caches watch (snoop on) bus to keep all processors' view of memory coherent

Snooping-Based Coherence

- Bus provides serialization point
  - Broadcast, totally ordered
  - Each cache controller "snoops" all bus transactions
  - Controller updates state of cache in response to processor and snoop events and generates bus transactions



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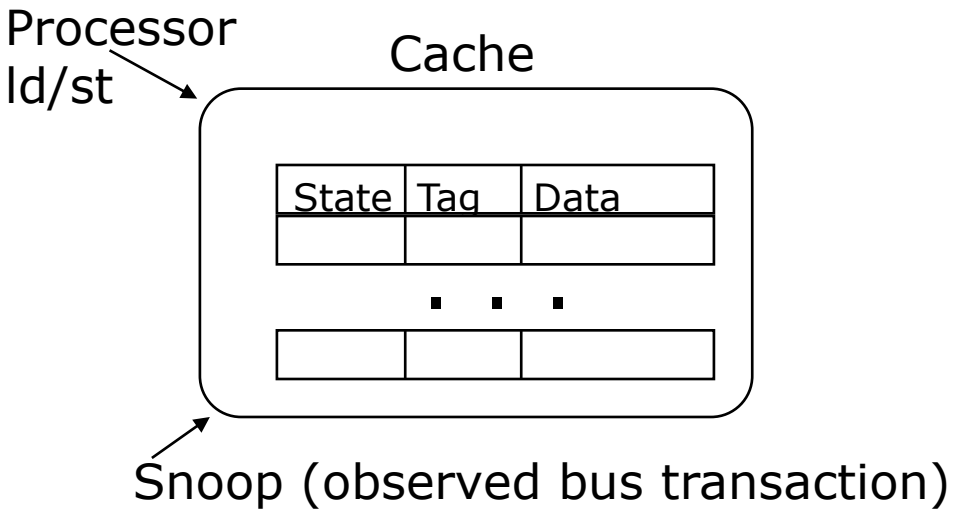
  - Actions



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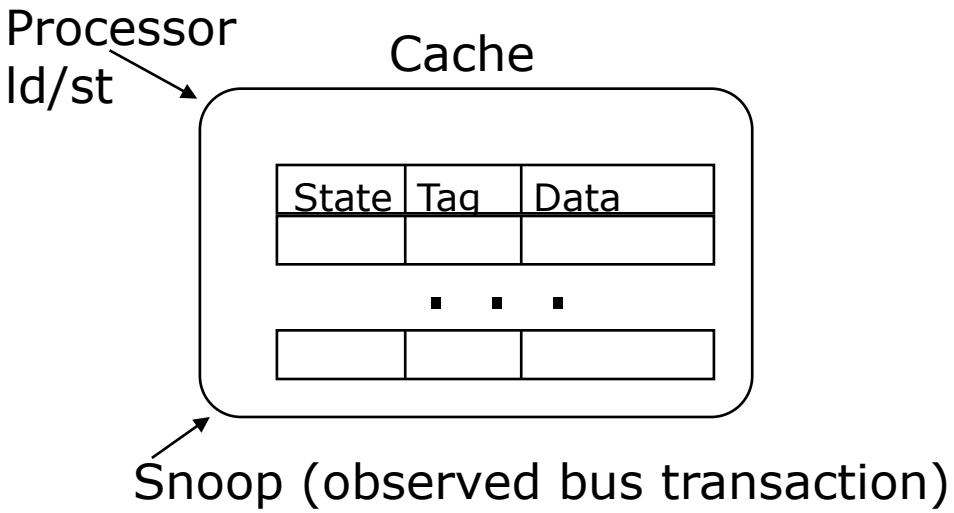
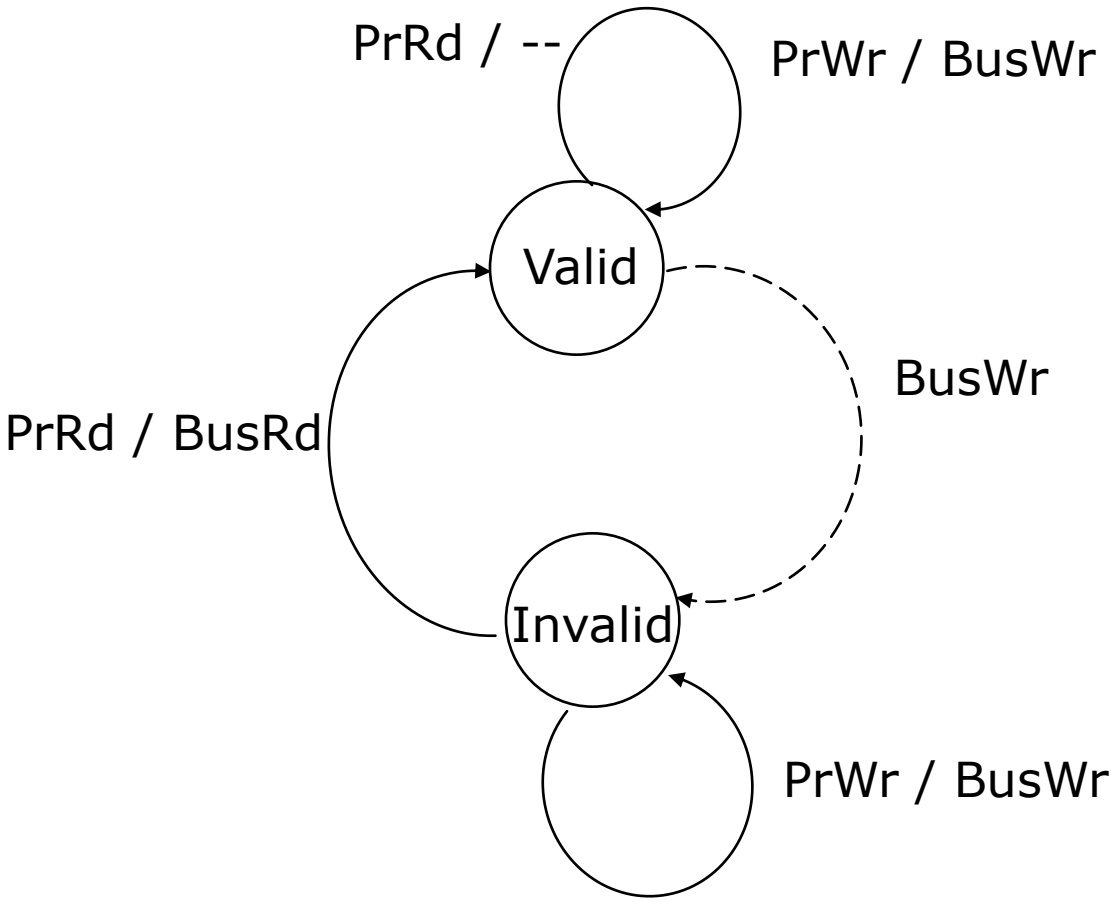
- Handling writes:
  - Write-invalidate
  - Write-update



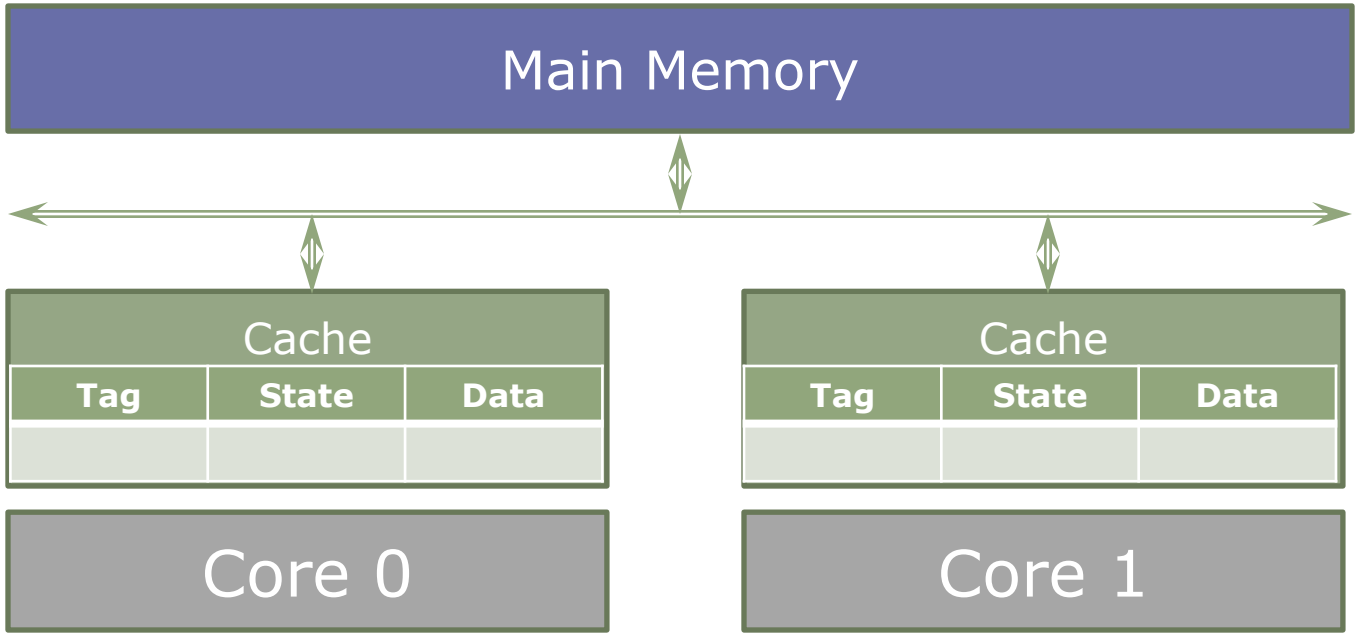
A Simple Protocol: Valid/Invalid (VI)



- Assume write-through caches

Actions
Processor Read (PrRd)
Processor Write (PrWr)
Bus Read (BusRd)
Bus Write (BusWr)

Valid/Invalid Example



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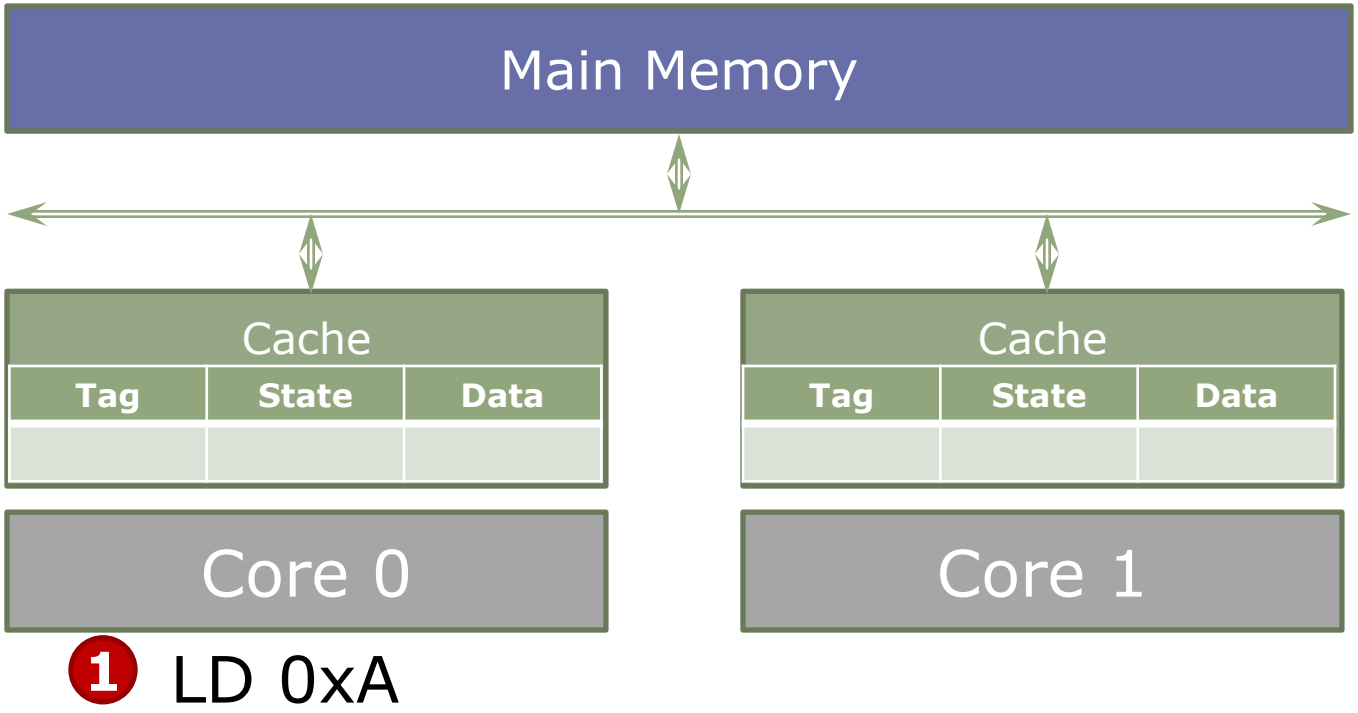

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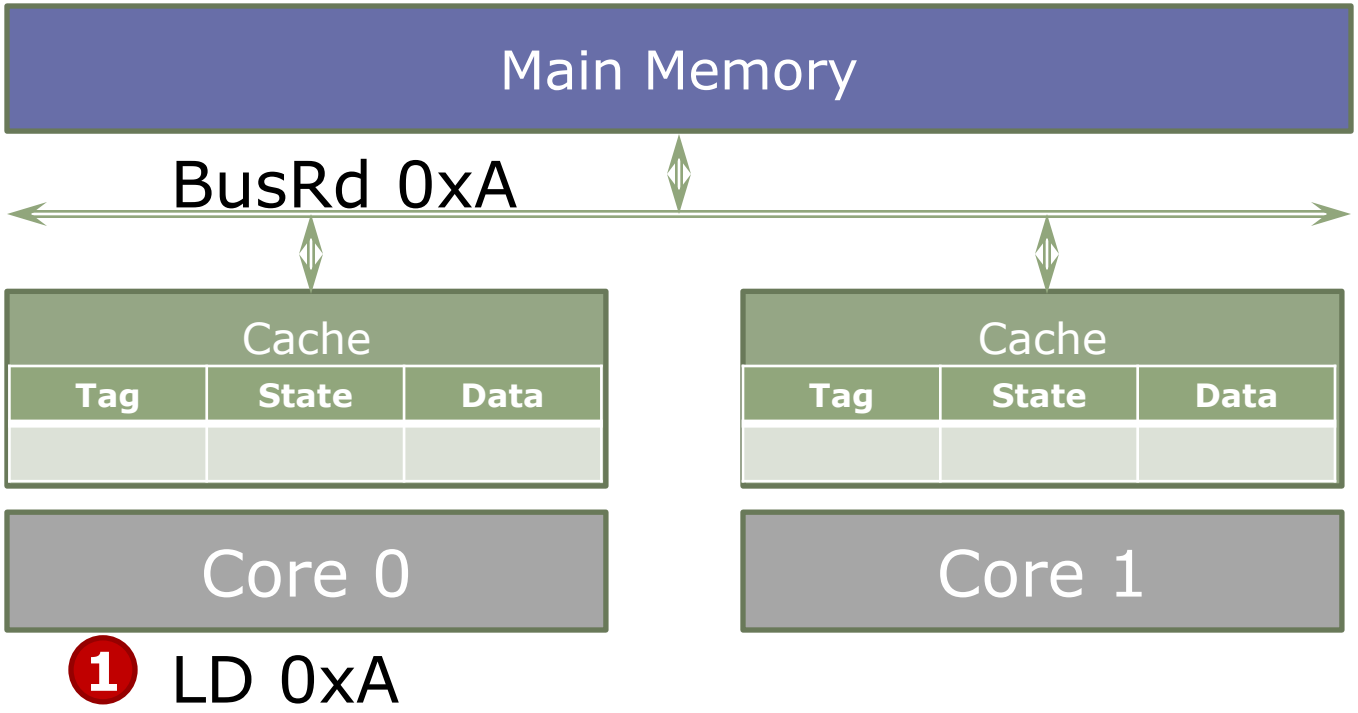

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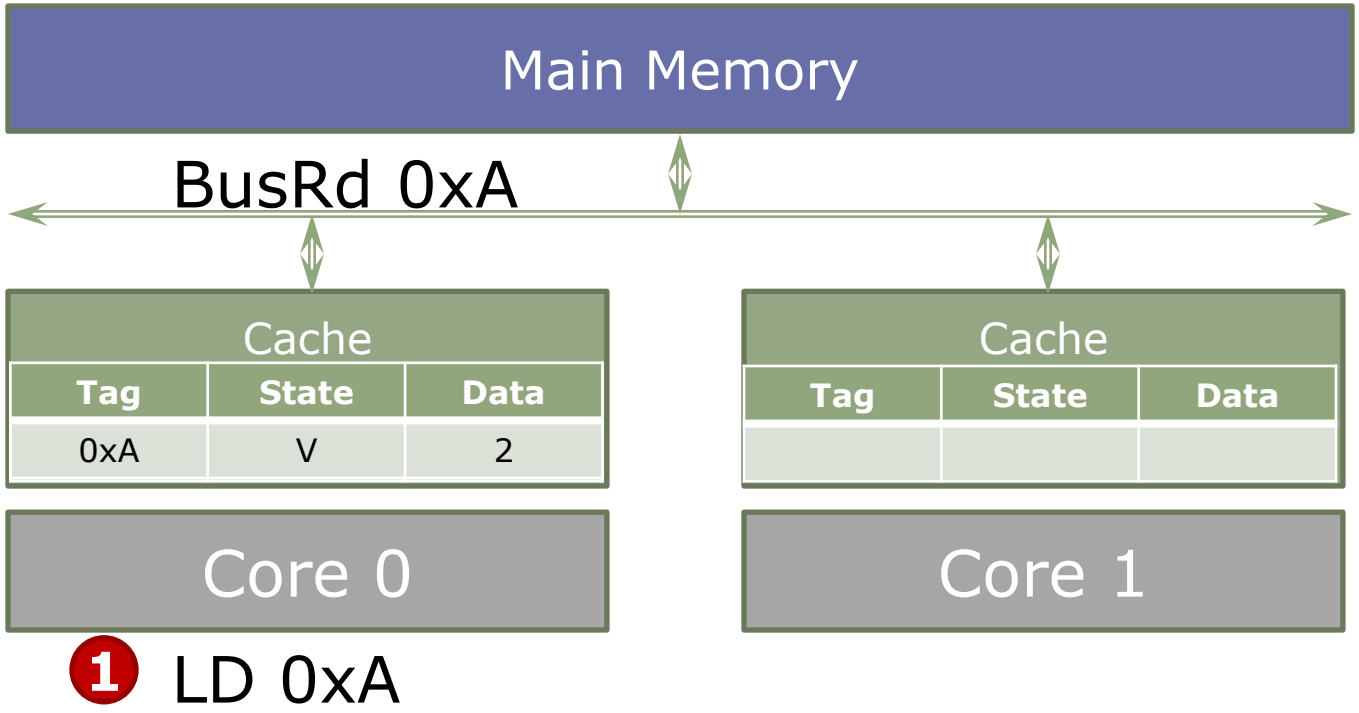

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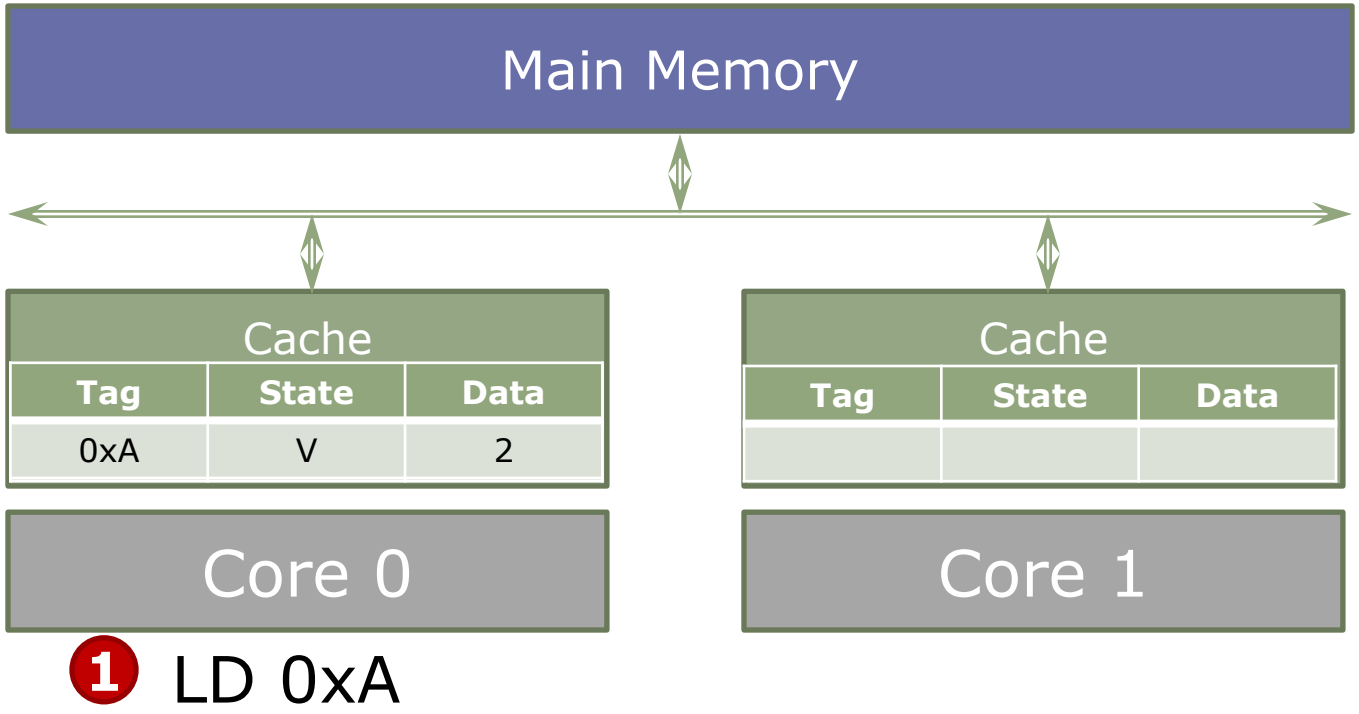

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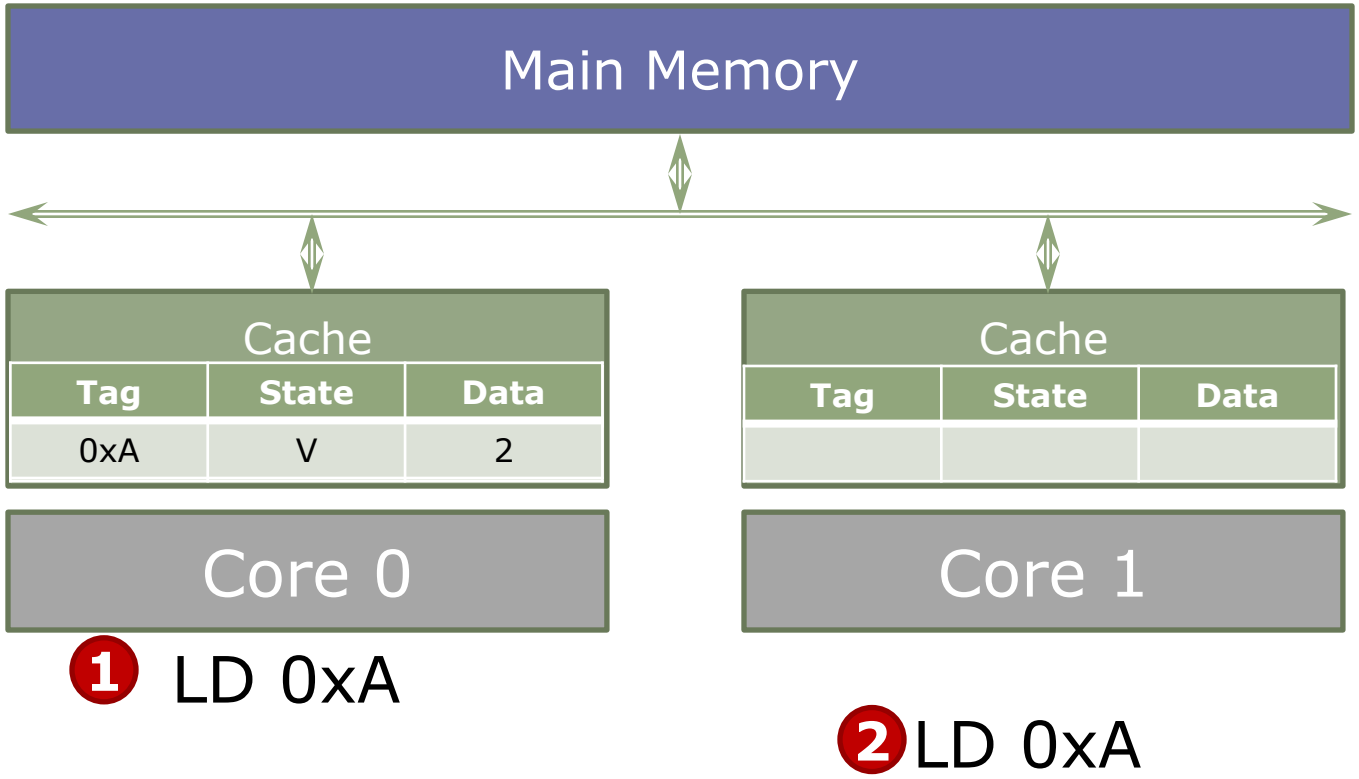

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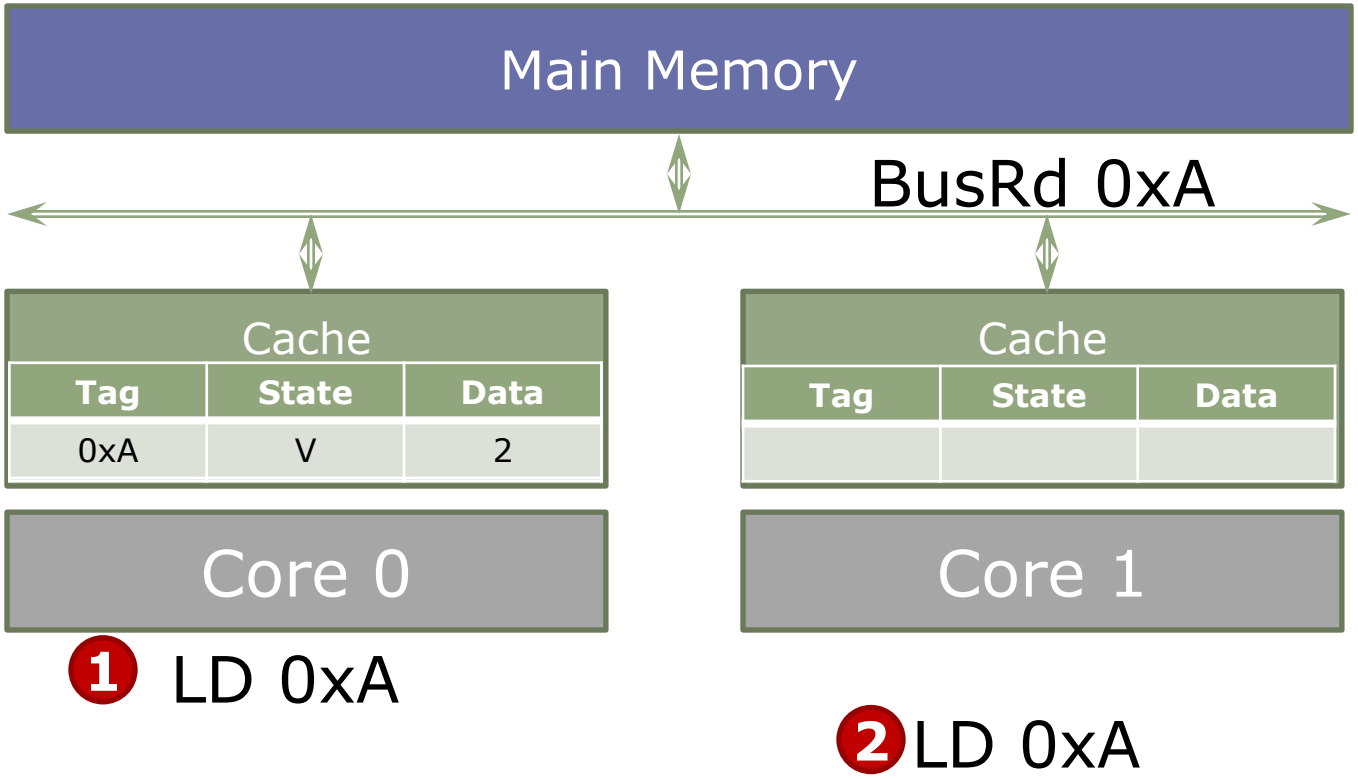

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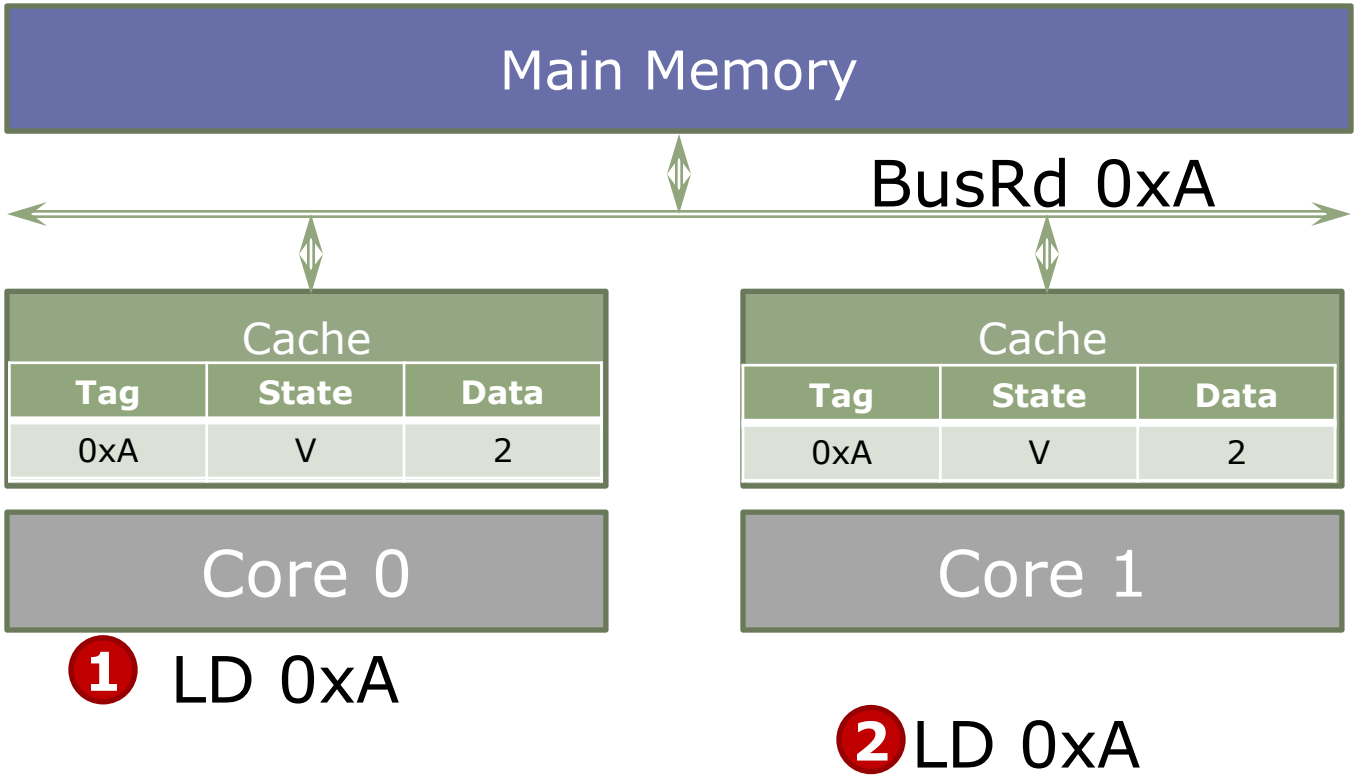

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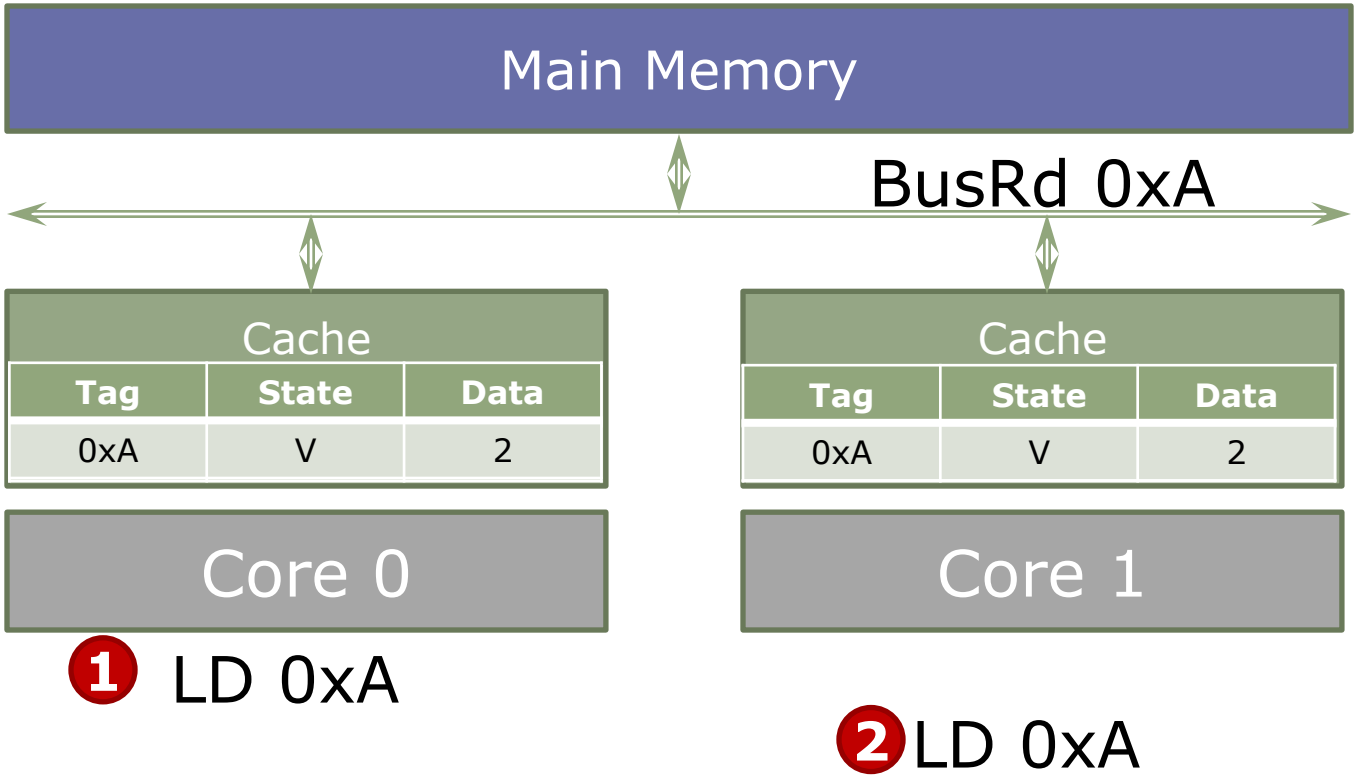

Additional loads satisfied locally, without BusRd

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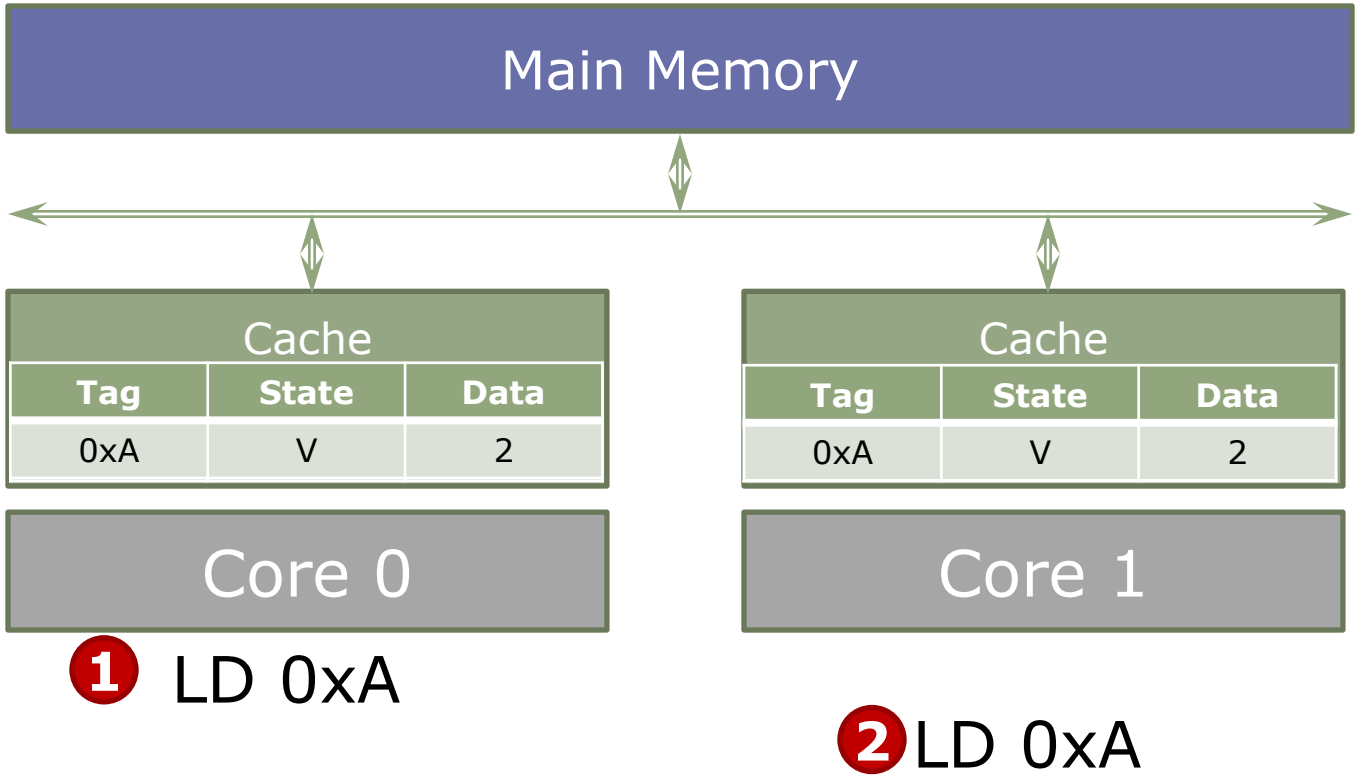

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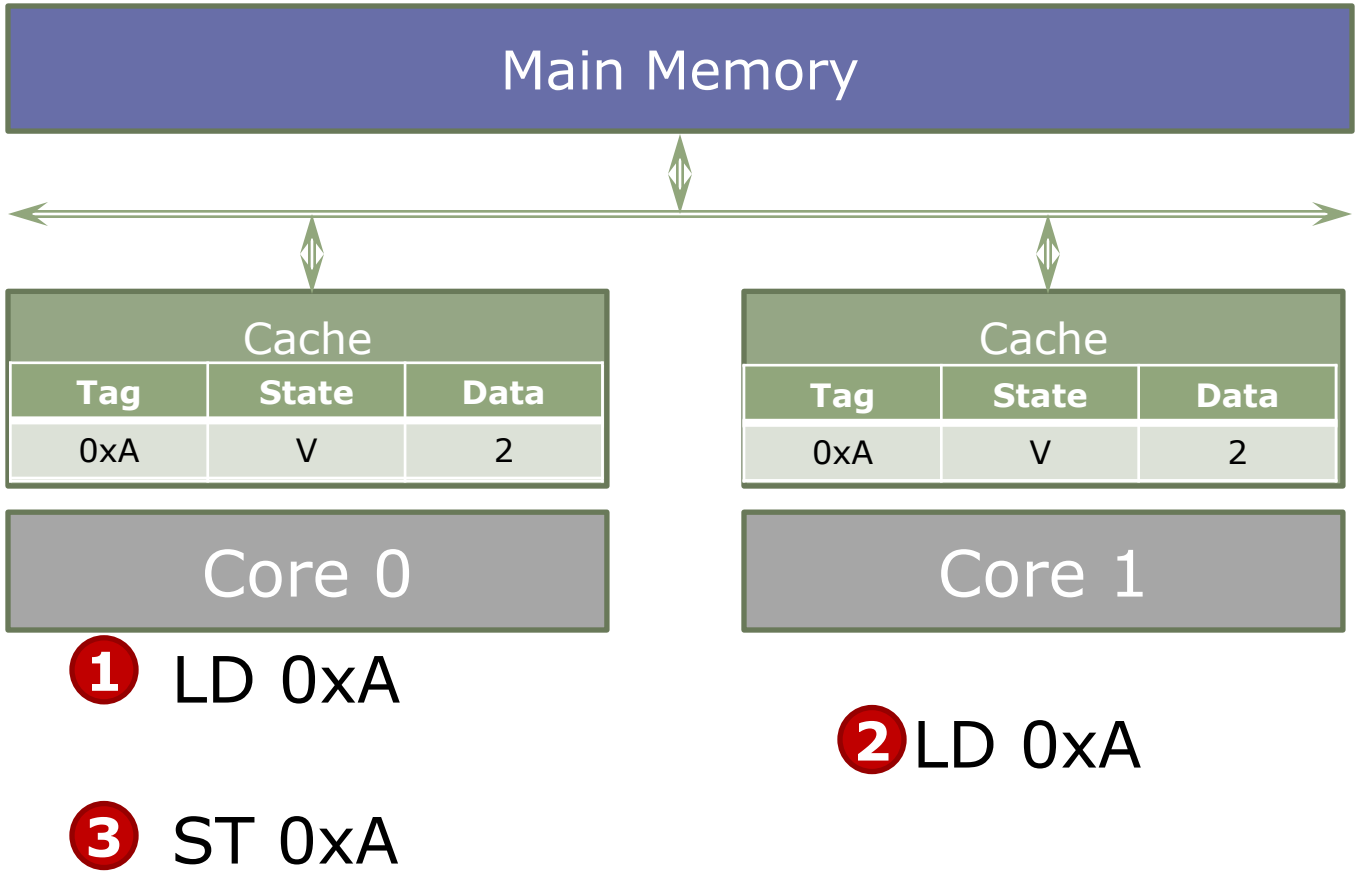

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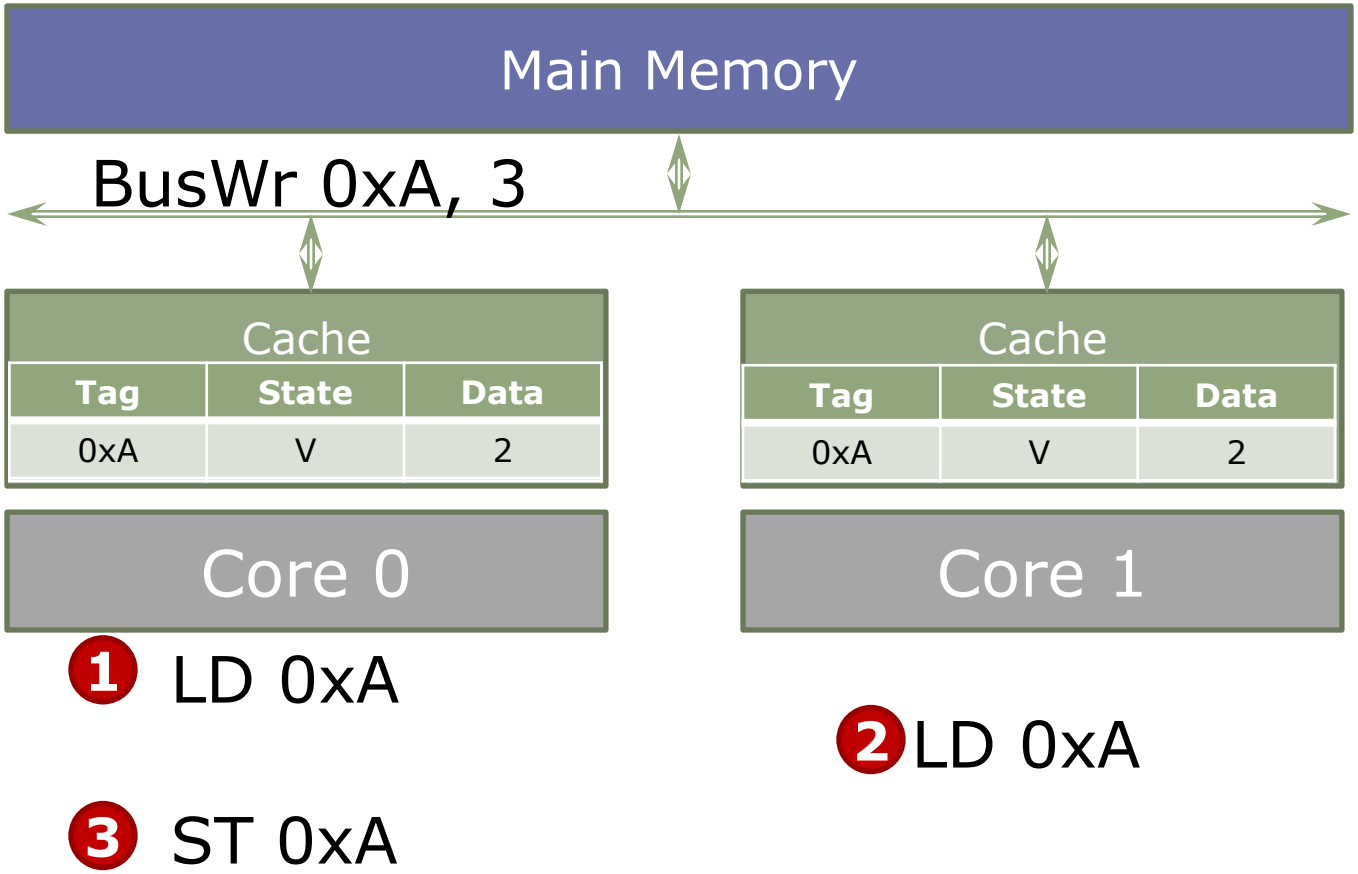
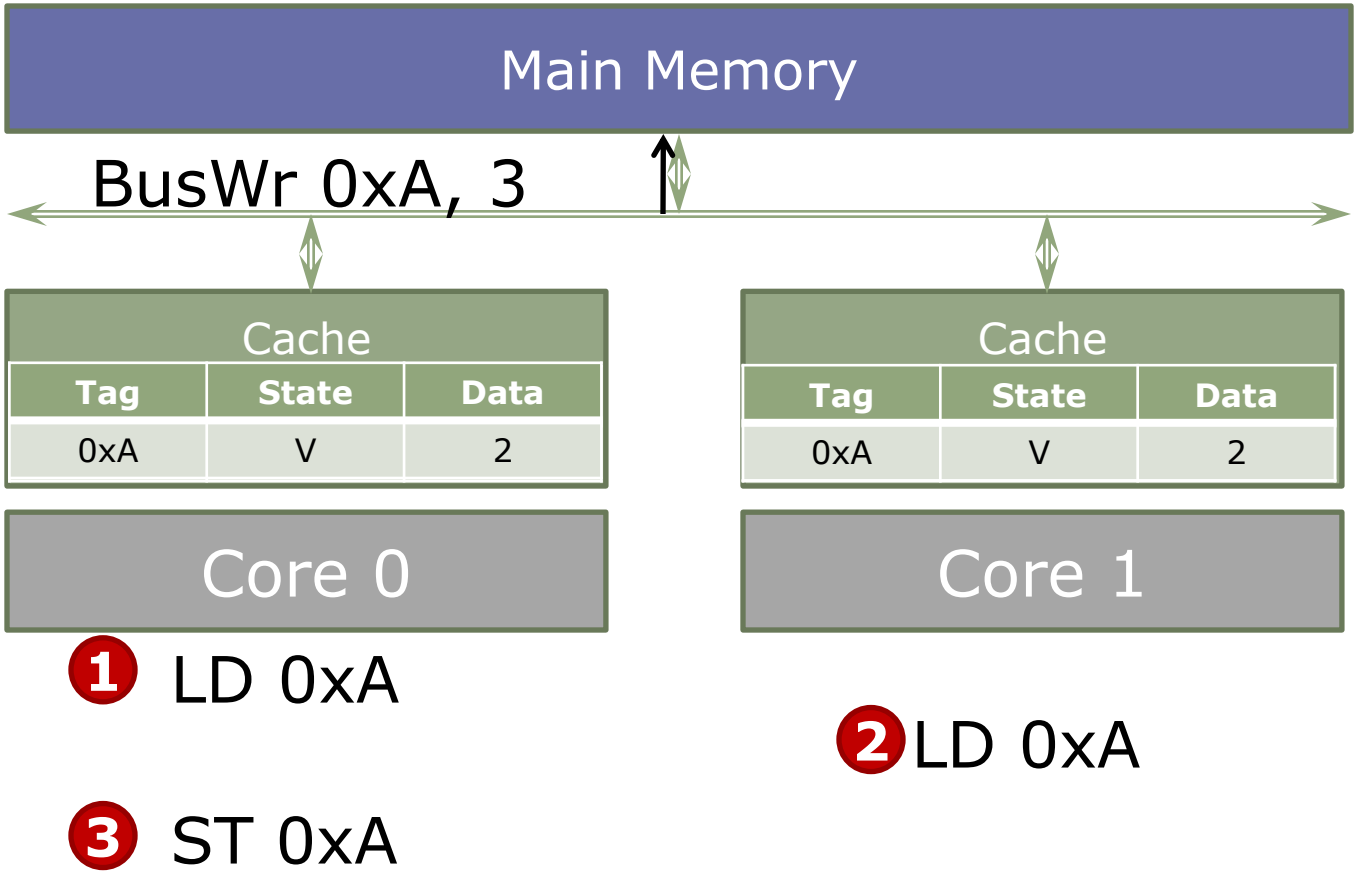
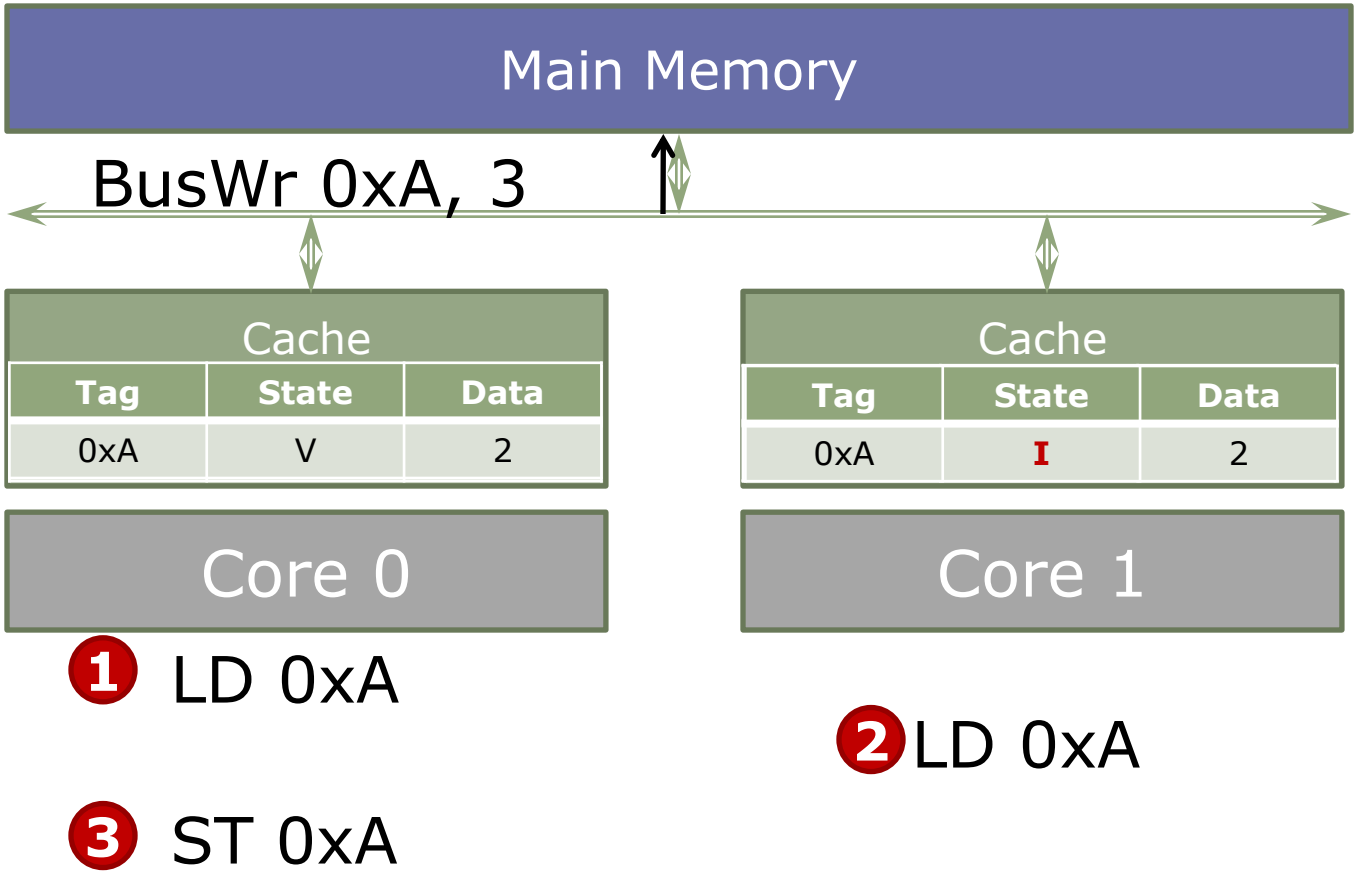


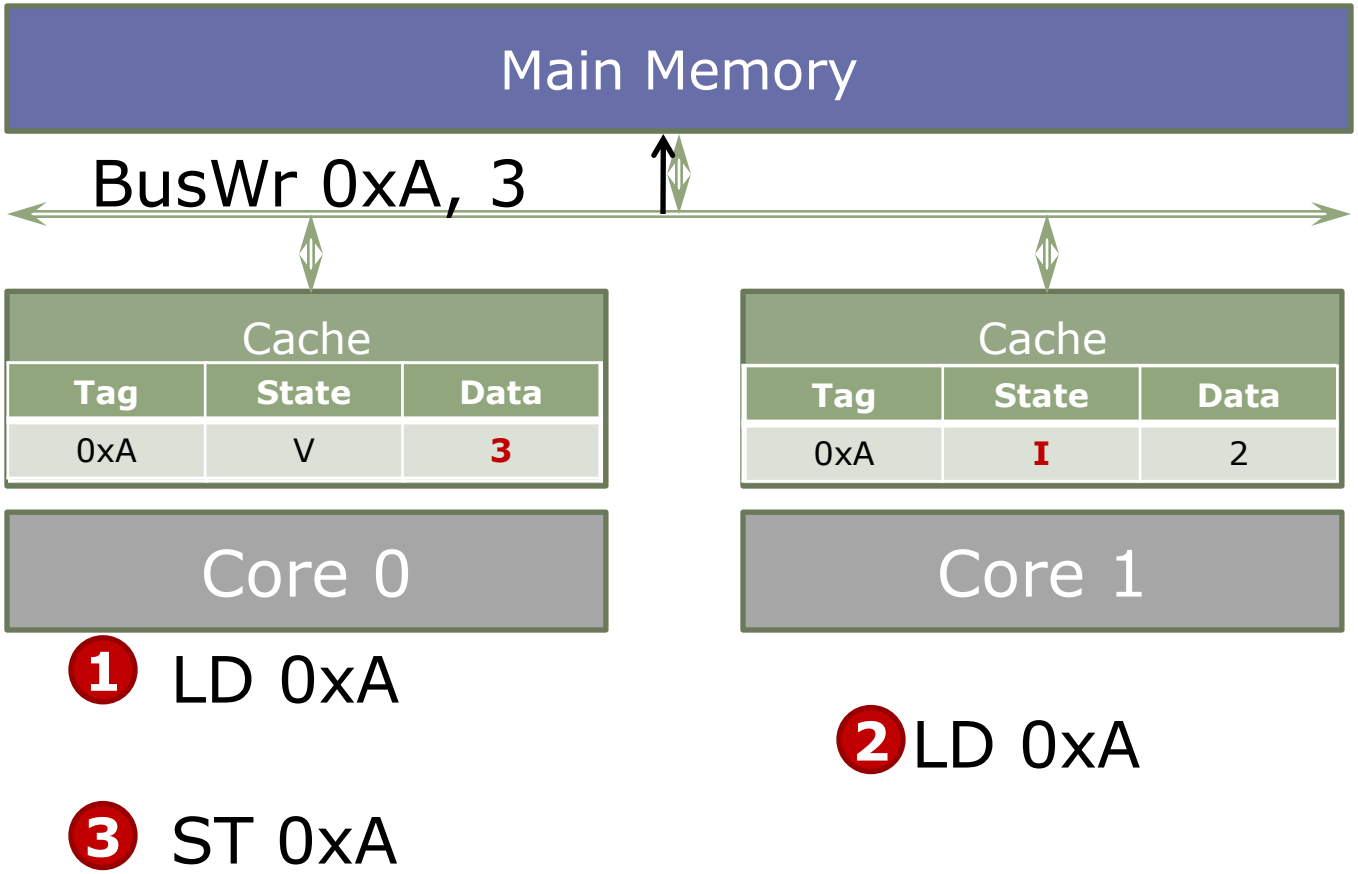
Valid/Invalid Example



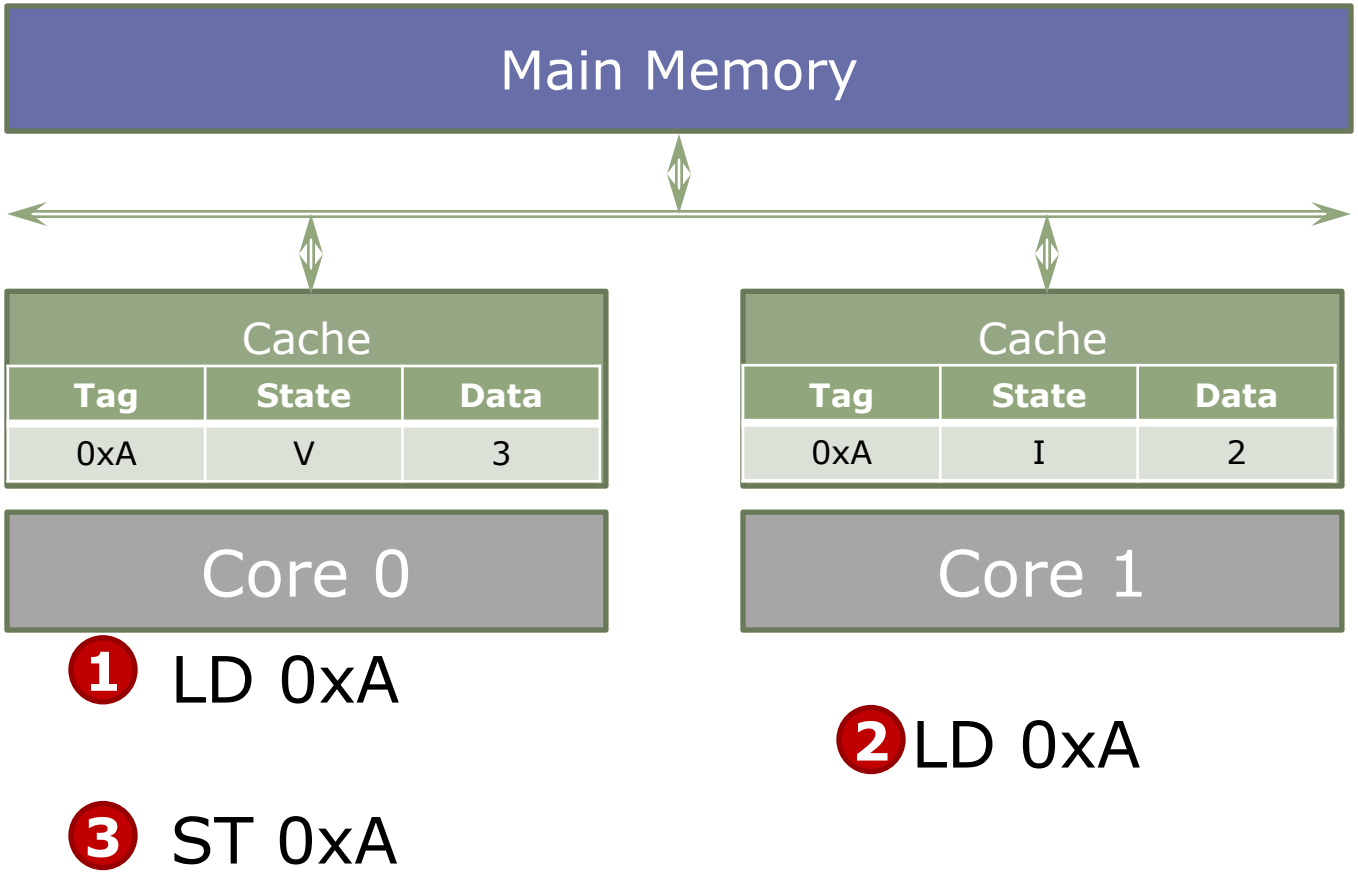
Valid/Invalid Example



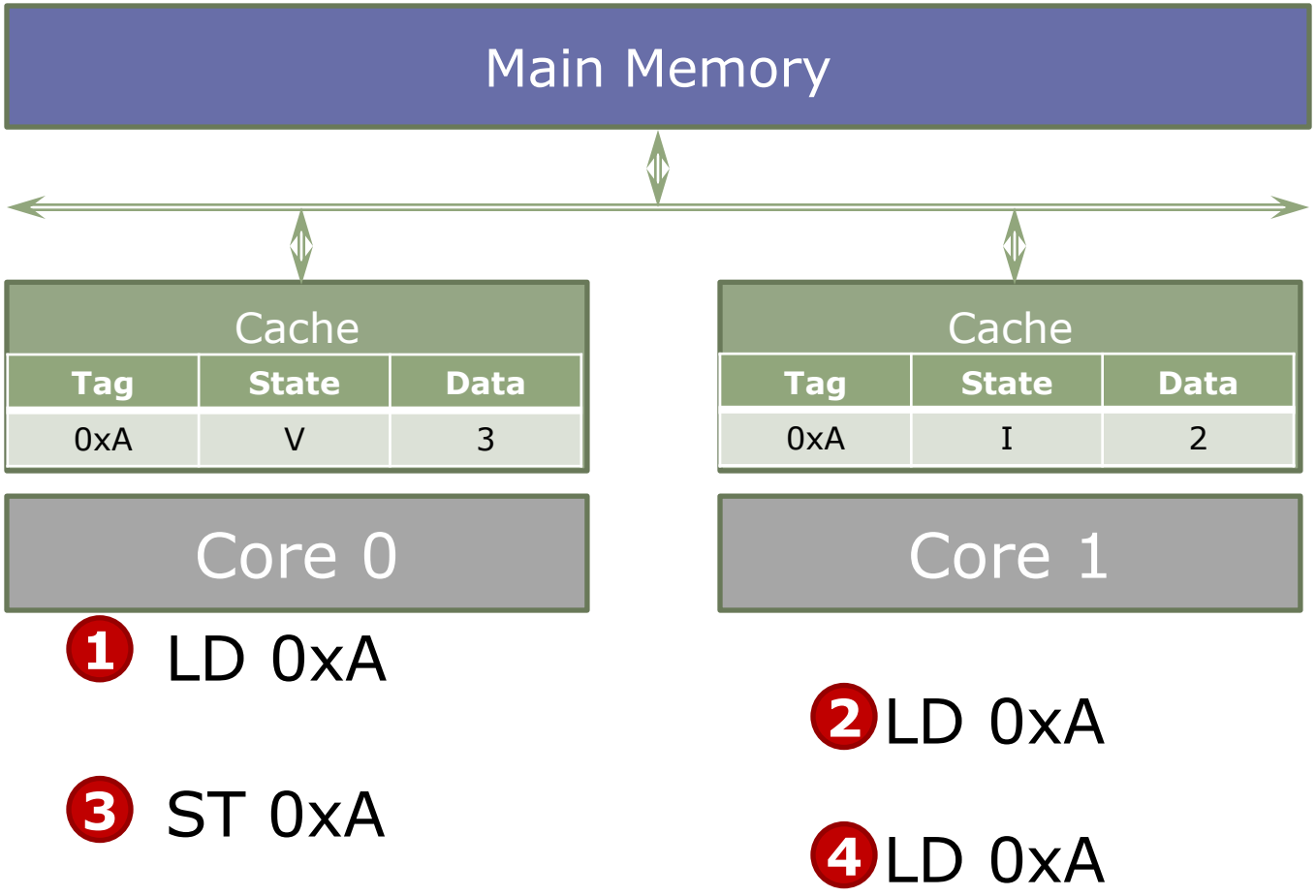
Valid/Invalid Example



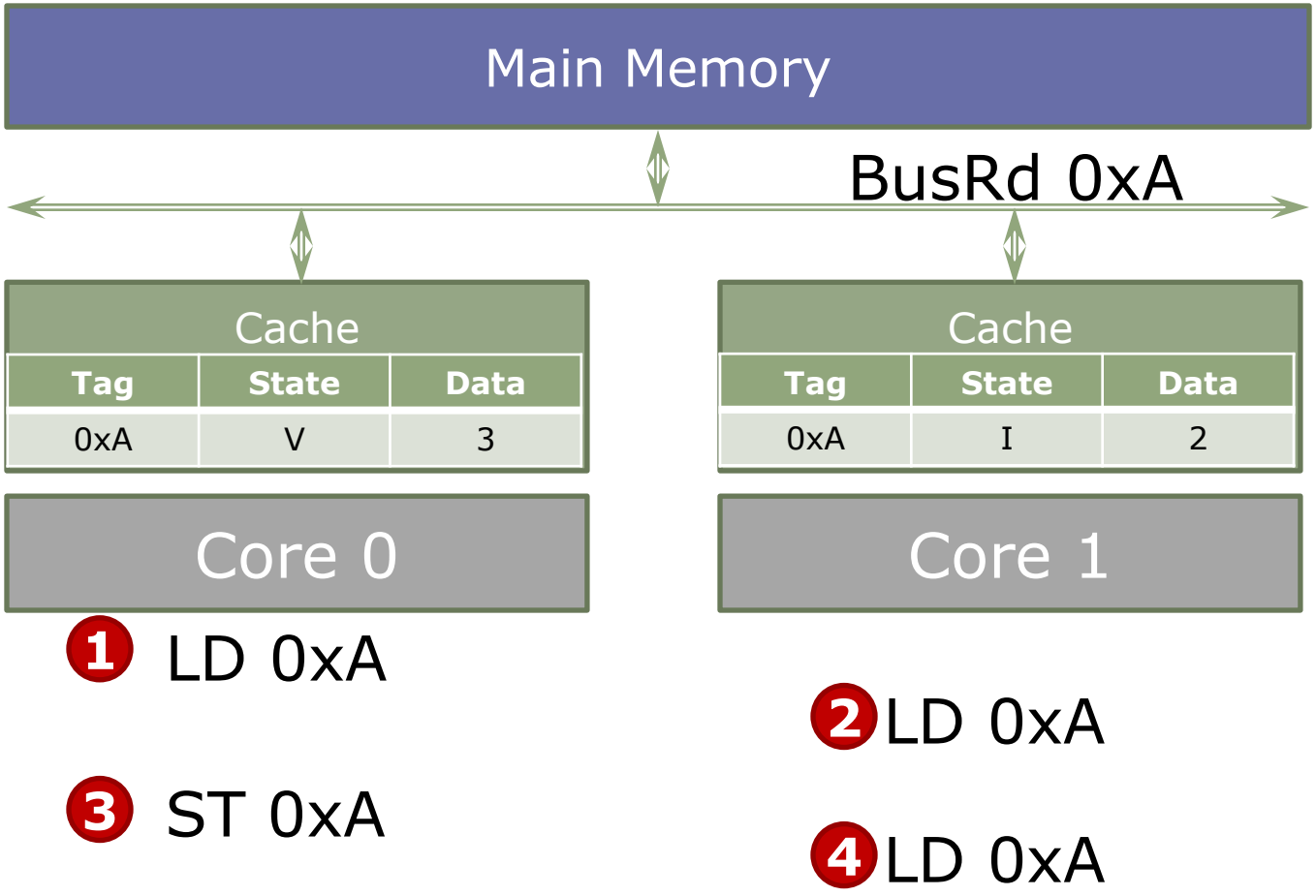
Valid/Invalid Example



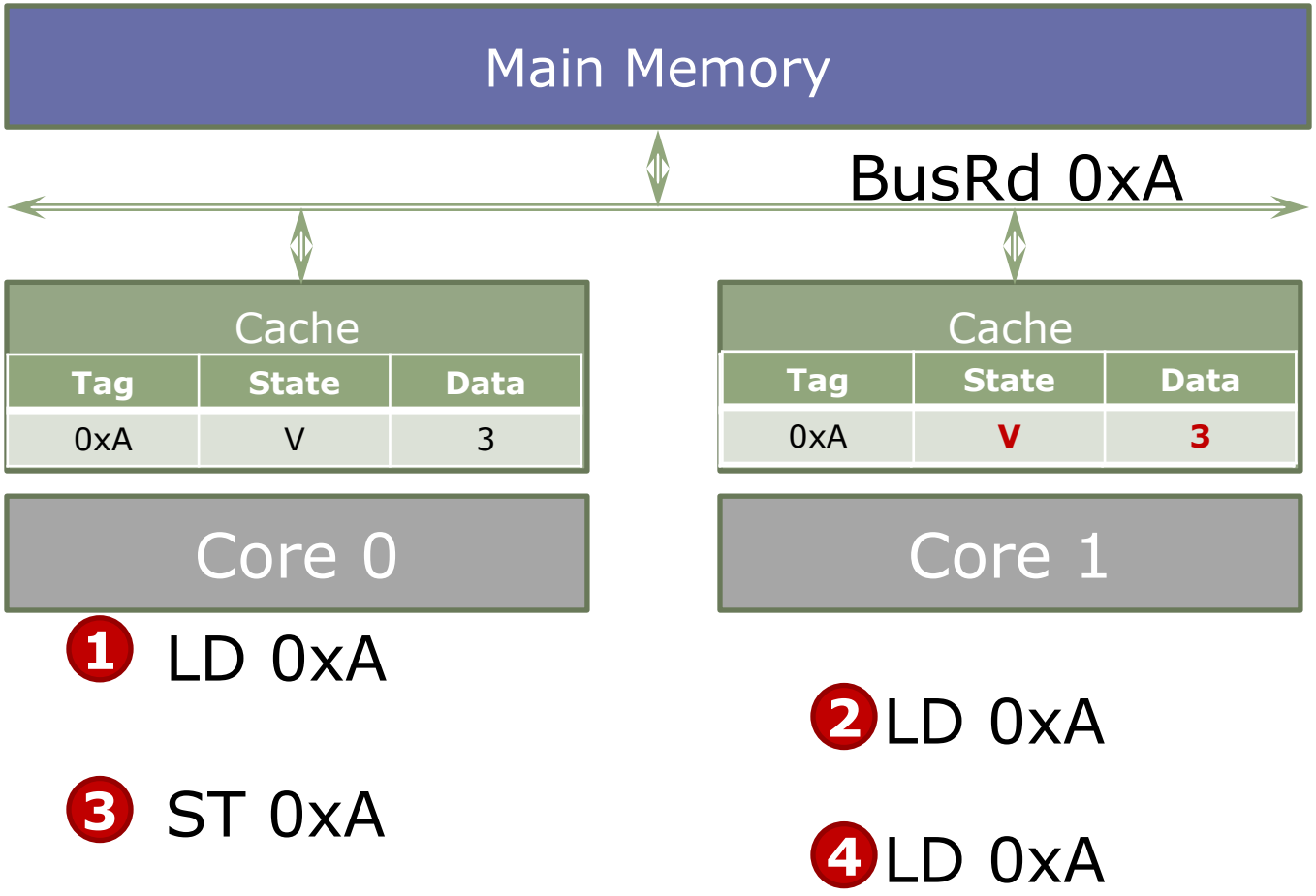
Valid/Invalid Example



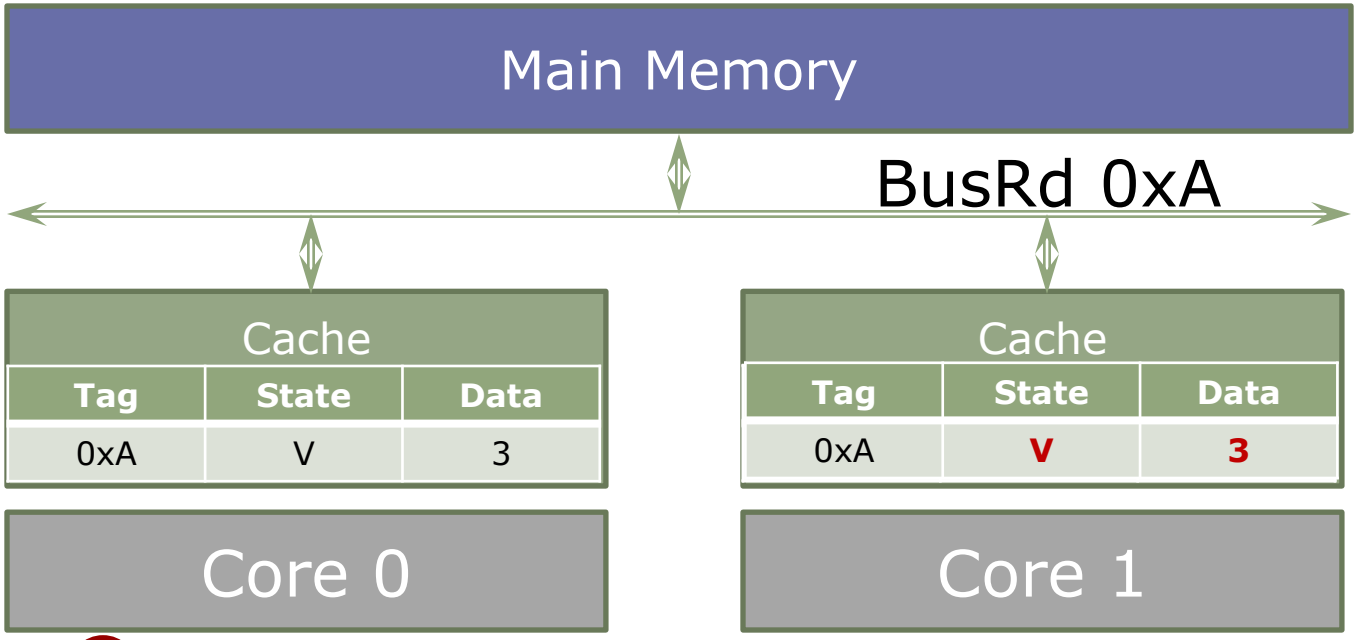
Valid/Invalid Example



Valid/Invalid Example

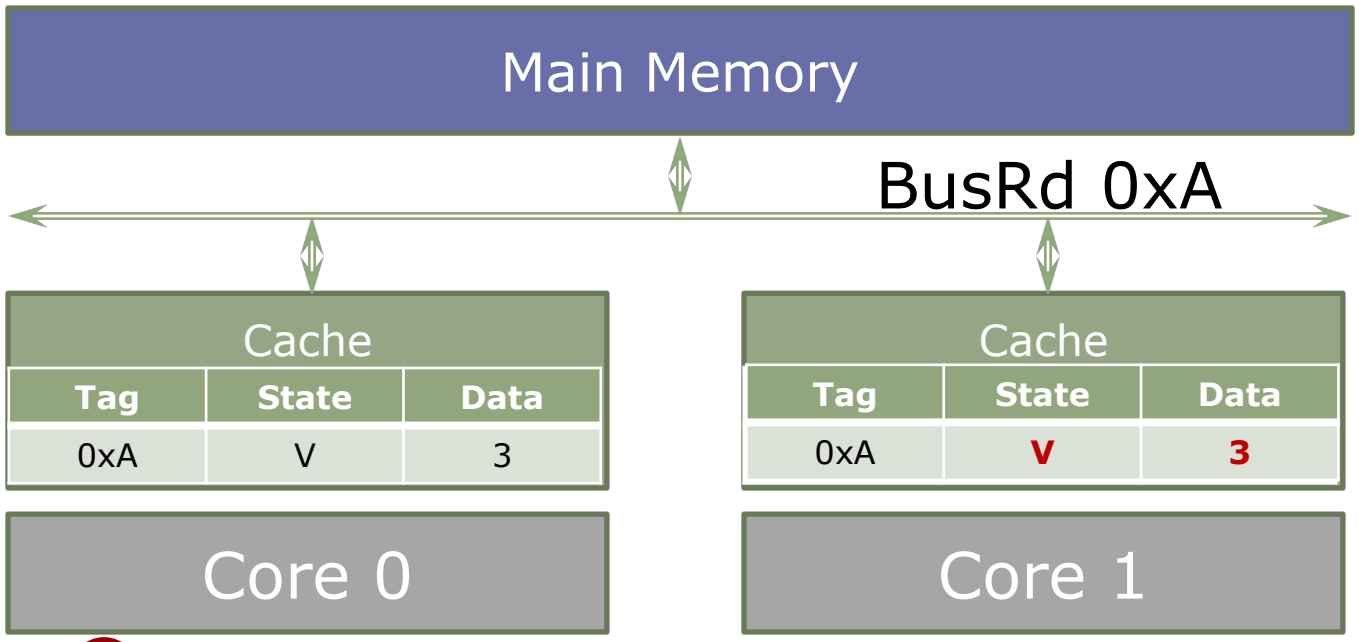


Valid/Invalid Example



VI Problems?

Valid/Invalid Example



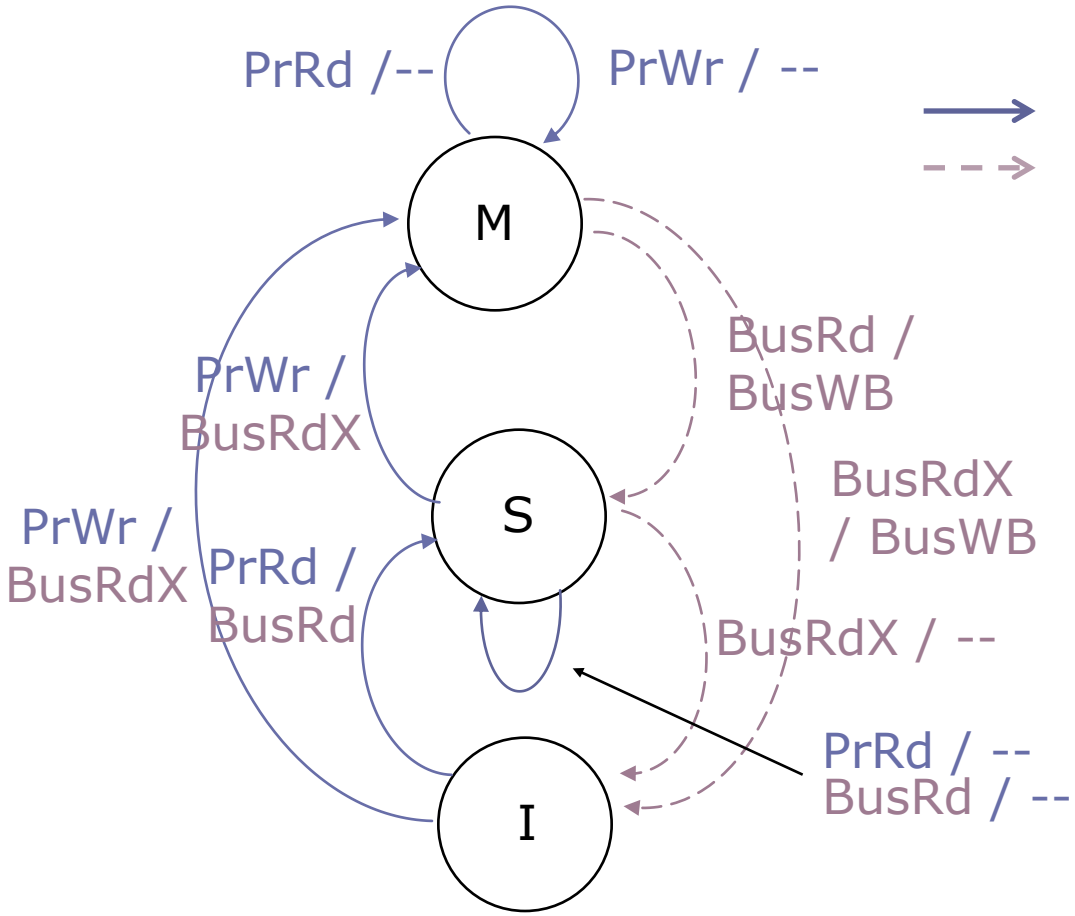
VI Problems?

Every write updates main memory

Every write requires broadcast & snoop

Modified/Shared/Invalid (MSI) Protocol

- Allows writeback caches + satisfying writes locally

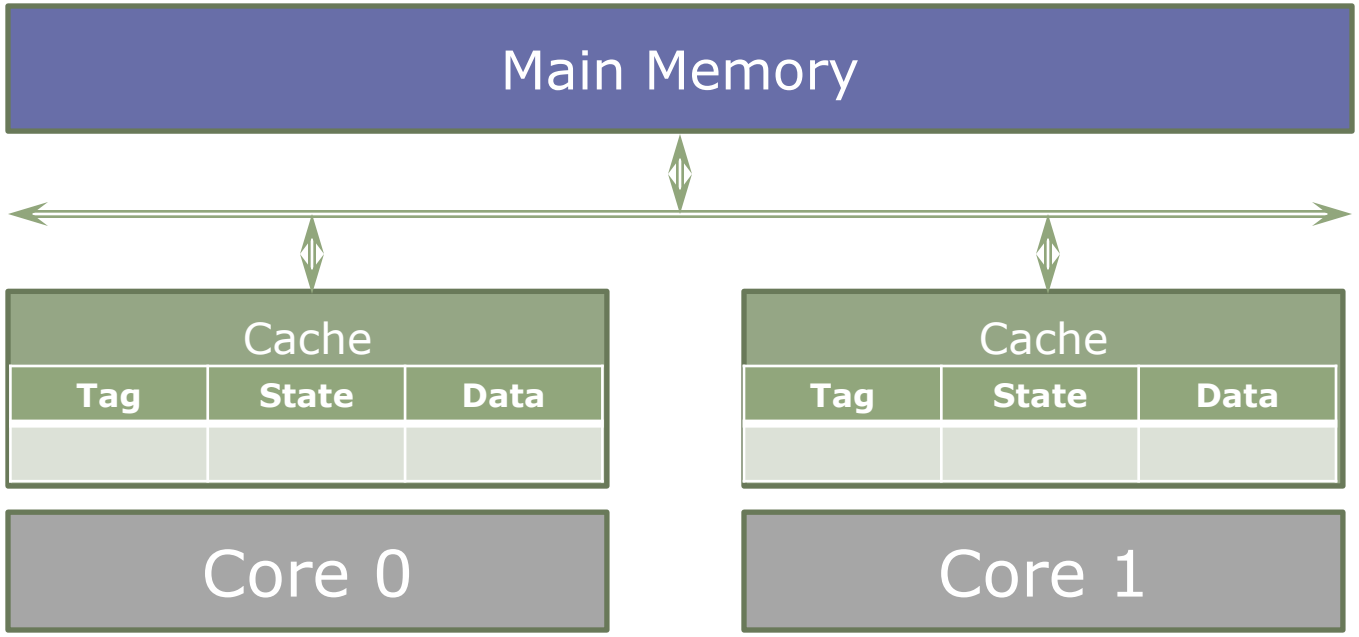


Processor-initiated transitions

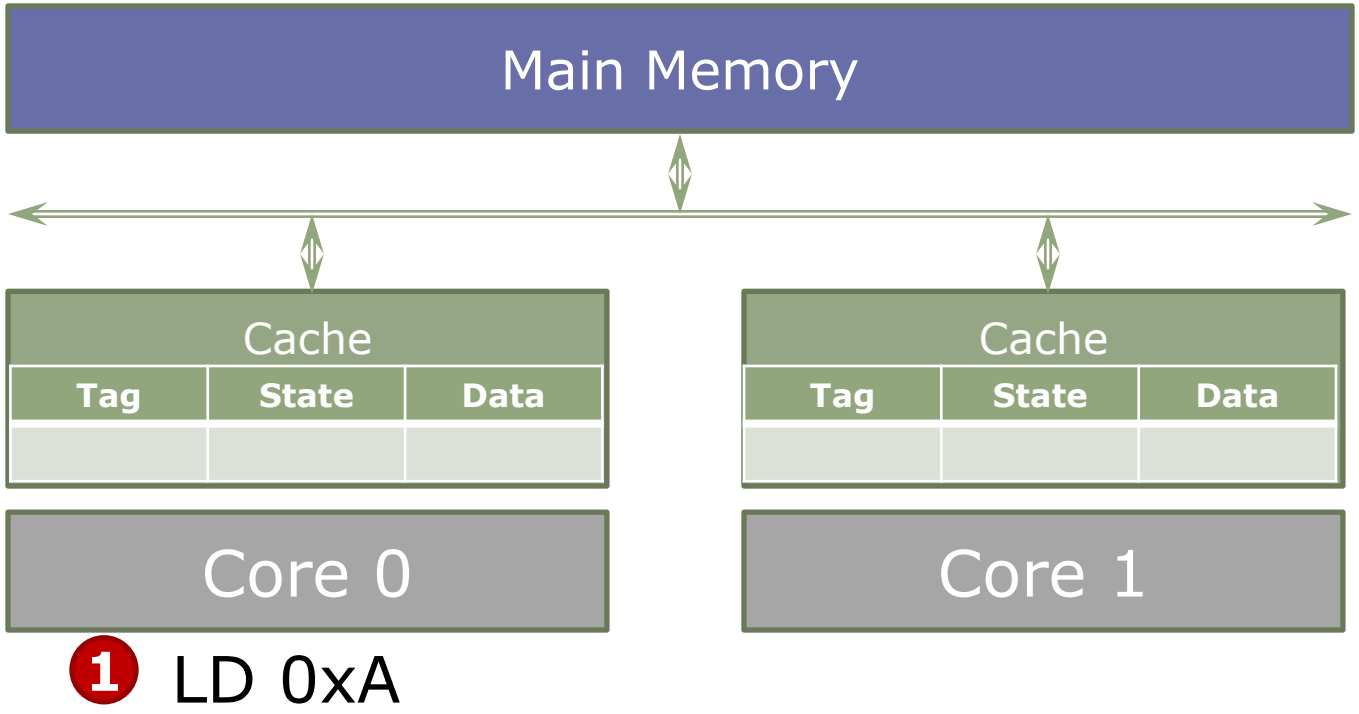
Bus-initiated transitions

Actions
Processor Read (PrRd)
Processor Write (PrWr)
Bus Read (BusRd)
Bus Read Exclusive (BusRdX)
Bus Writeback (BusWB)

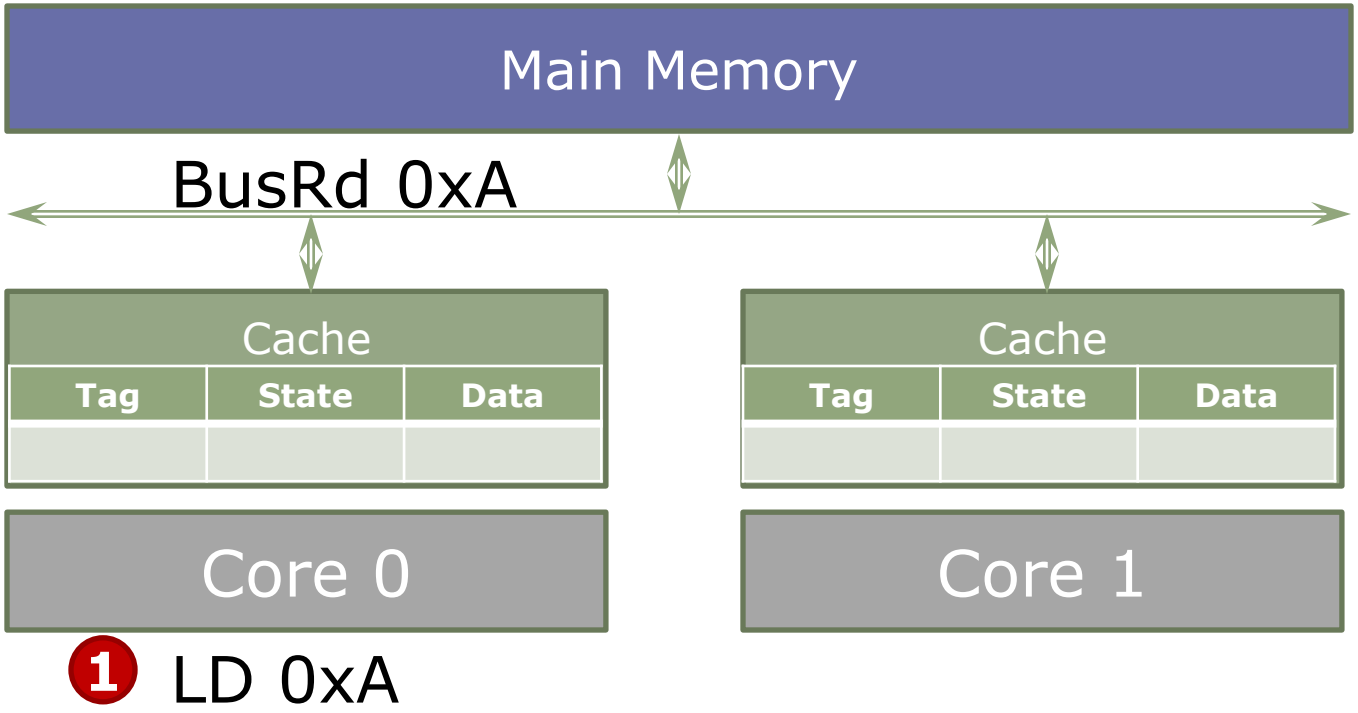
MSI Example


Main Memory
Cache
Tag
State
Data
Cache
Tag
State
Data
Core 0
Core 1

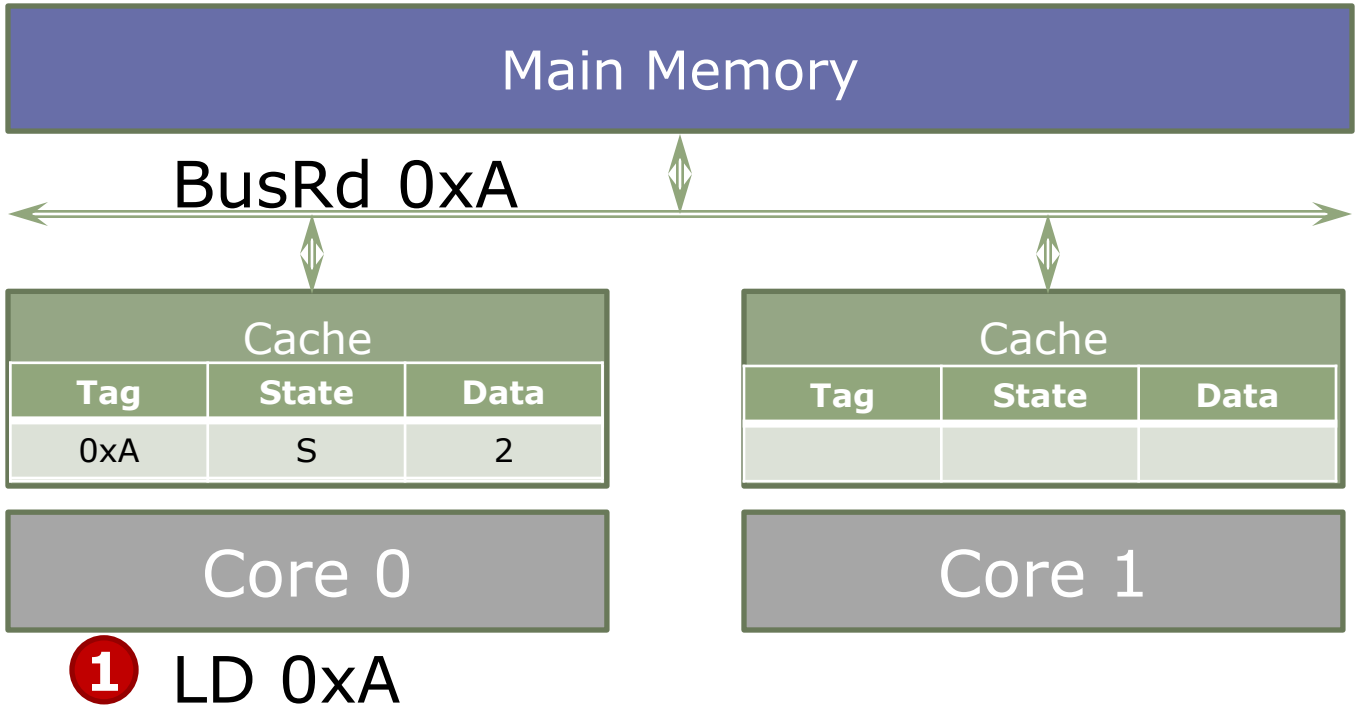
MSI Example


Main Memory
Cache
Tag
State
Data
Cache
Tag
State
Data
Core 0
Core 1
1
LD 0xA

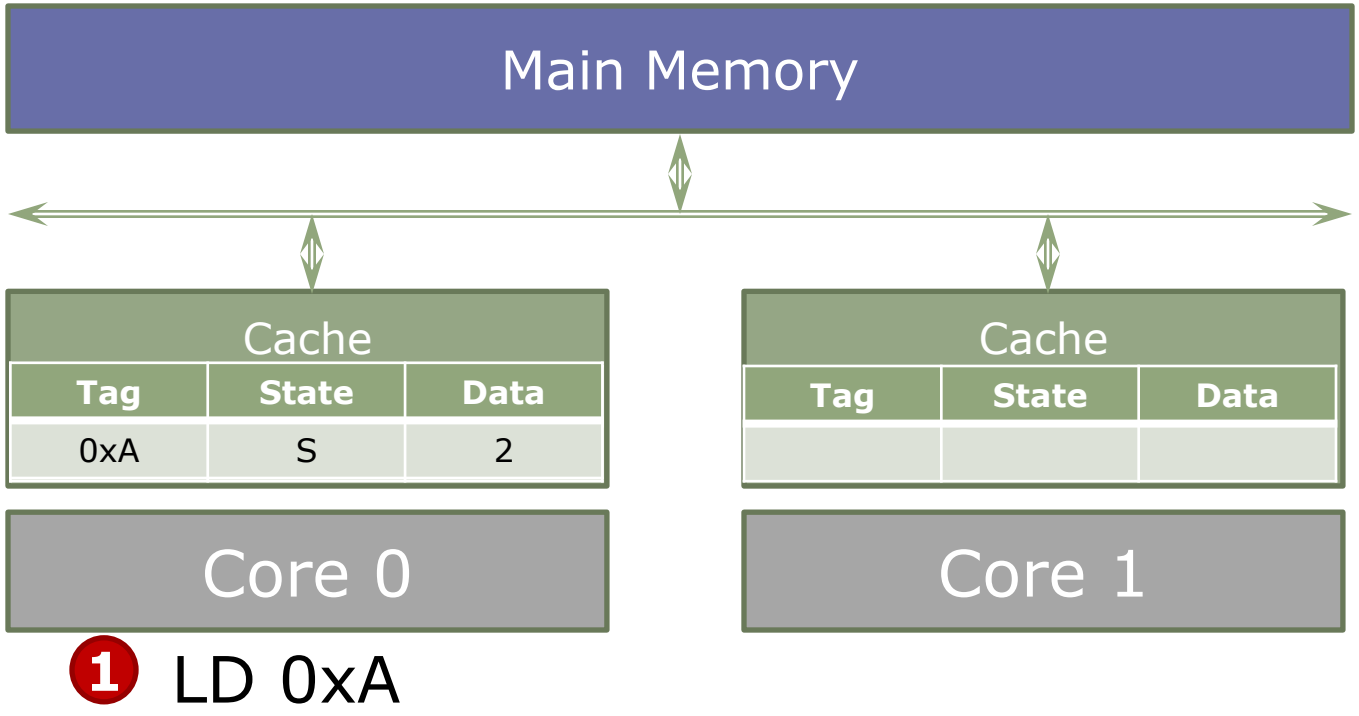
MSI Example


Main Memory
BusRd 0xA
Cache
Tag
State
Data
Cache
Tag
State
Data
Core 0
Core 1
1
LD 0xA

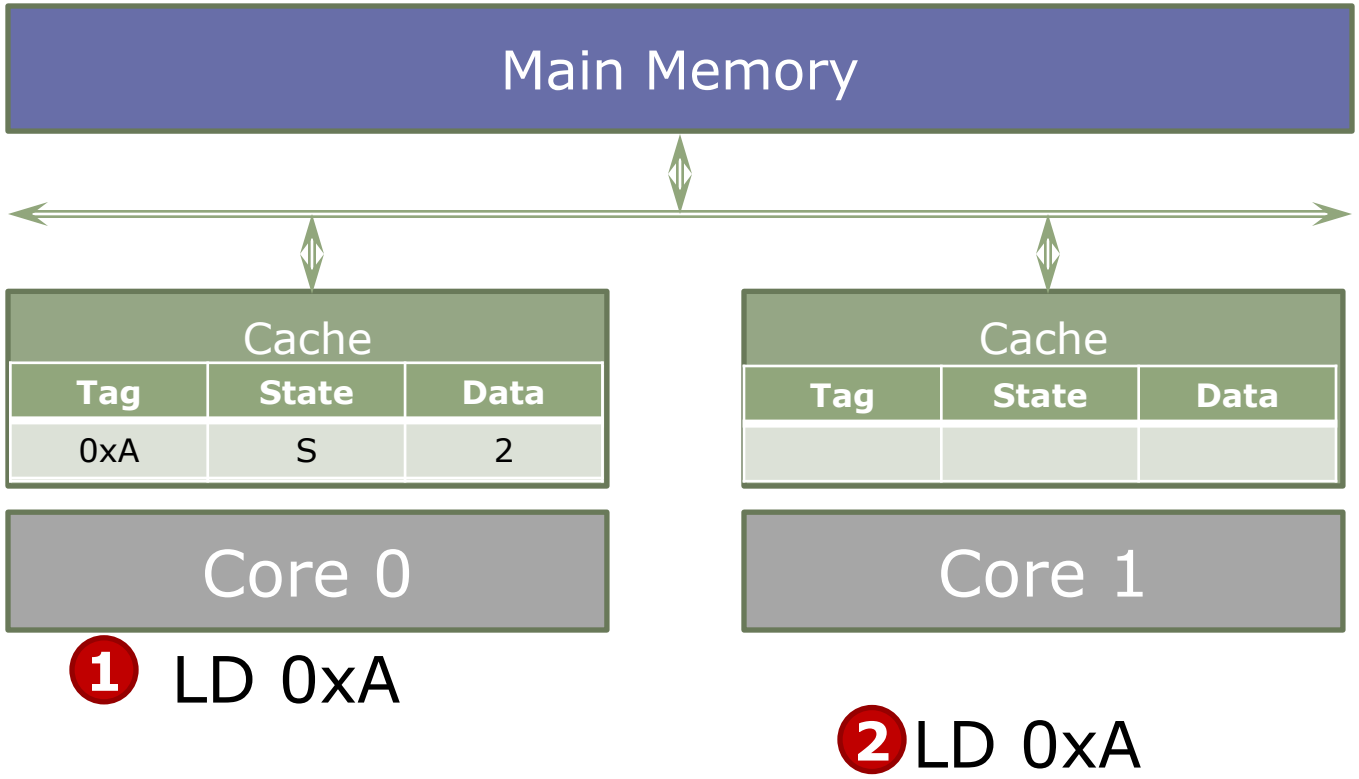
MSI Example


Main Memory
BusRd 0xA
Cache
Tag
State
Data
0xA
S
2
Cache
Tag
State
Data
Core 0
Core 1
1
LD 0xA

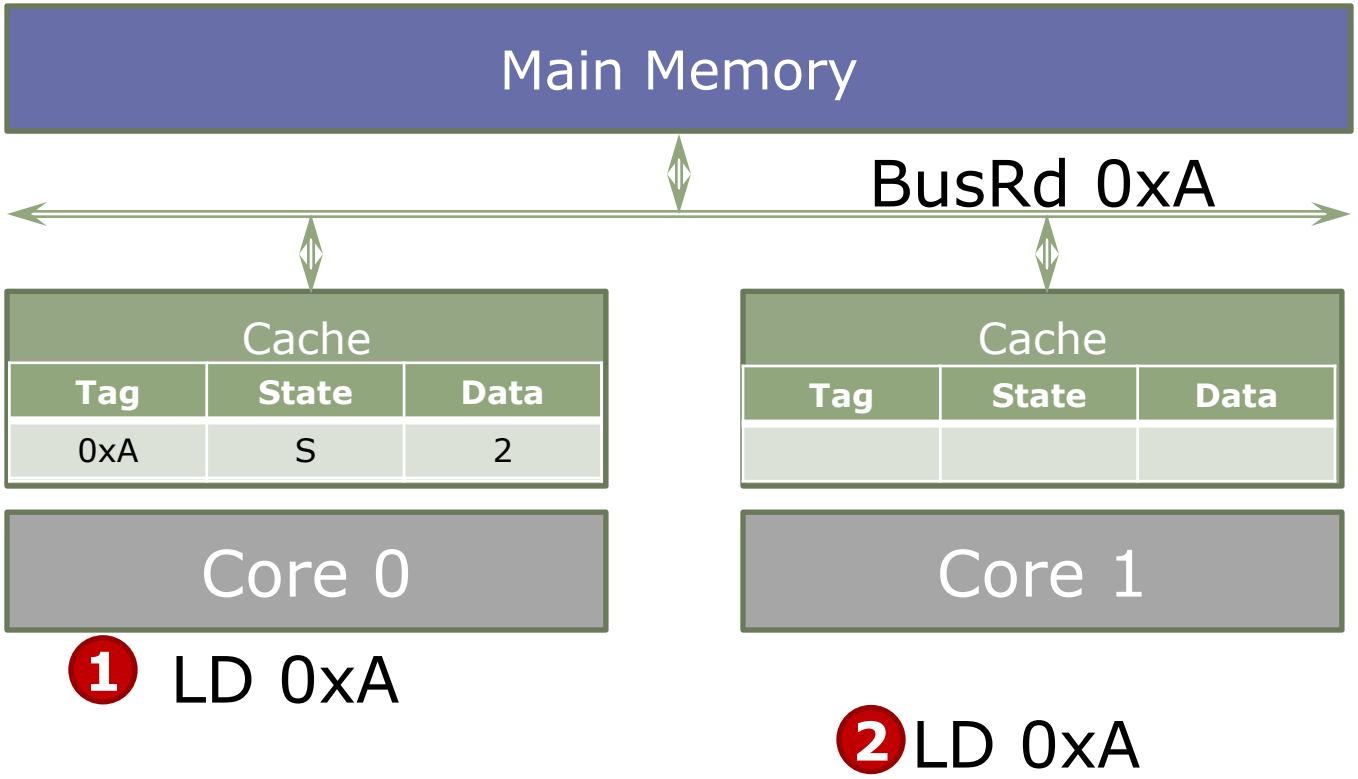
MSI Example


Main Memory
Cache
Tag
State
Data
0xA
S
2
Cache
Tag
State
Data
Core 0
Core 1
1
LD 0xA

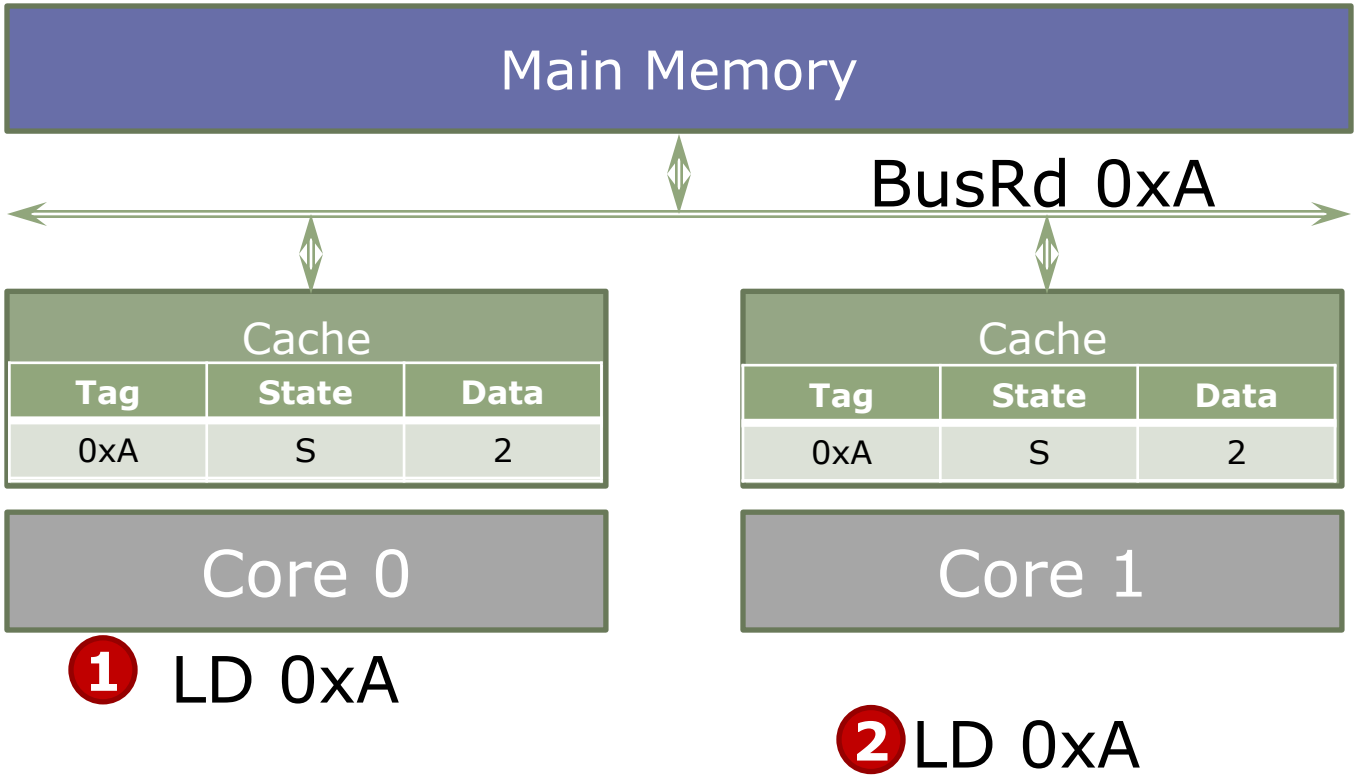
MSI Example


Main Memory
Cache
Tag
State
Data
0xA
S
2
Cache
Tag
State
Data
Core 0
Core 1
1
LD 0xA
2
LD 0xA

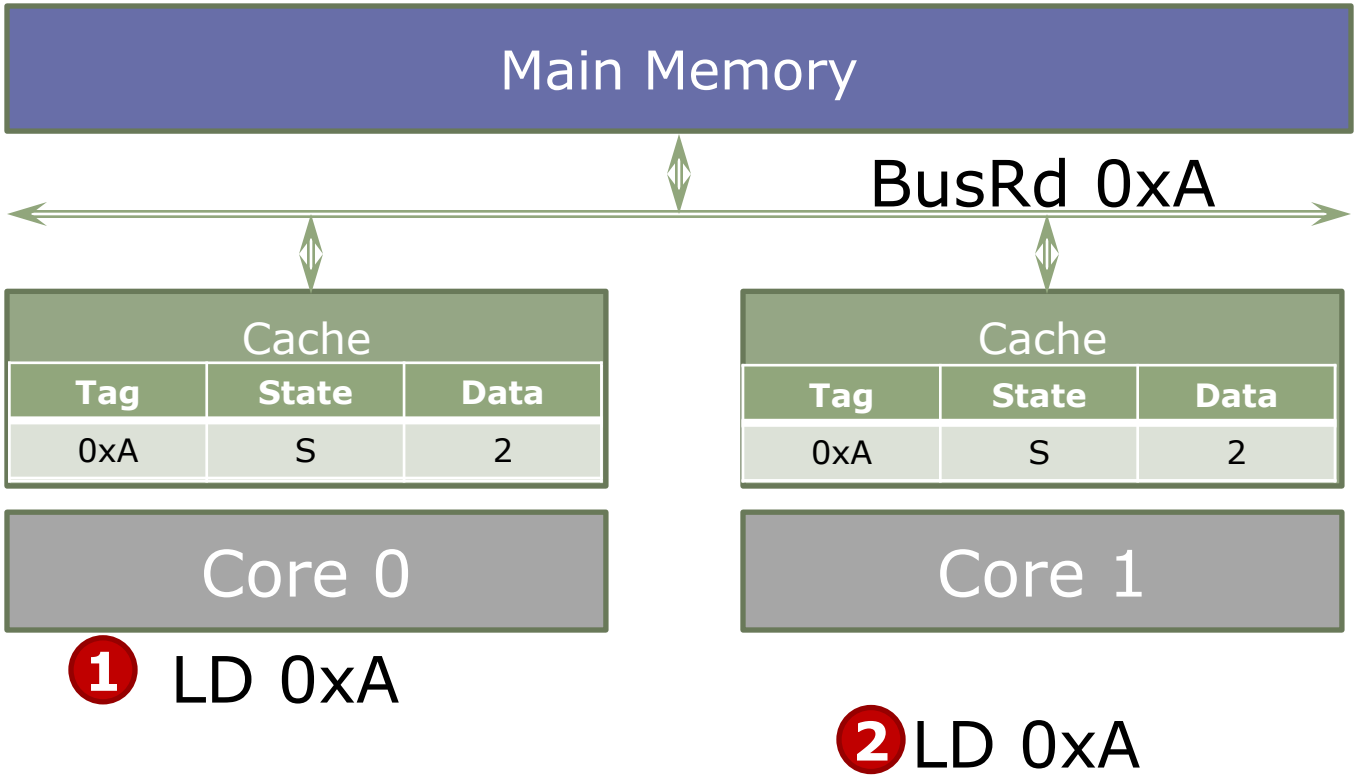
MSI Example


Main Memory
BusRd 0xA
Cache
Tag
State
Data
0xA
S
2
Cache
Tag
State
Data
Core 0
Core 1
1
LD 0xA
2
LD 0xA

MSI Example

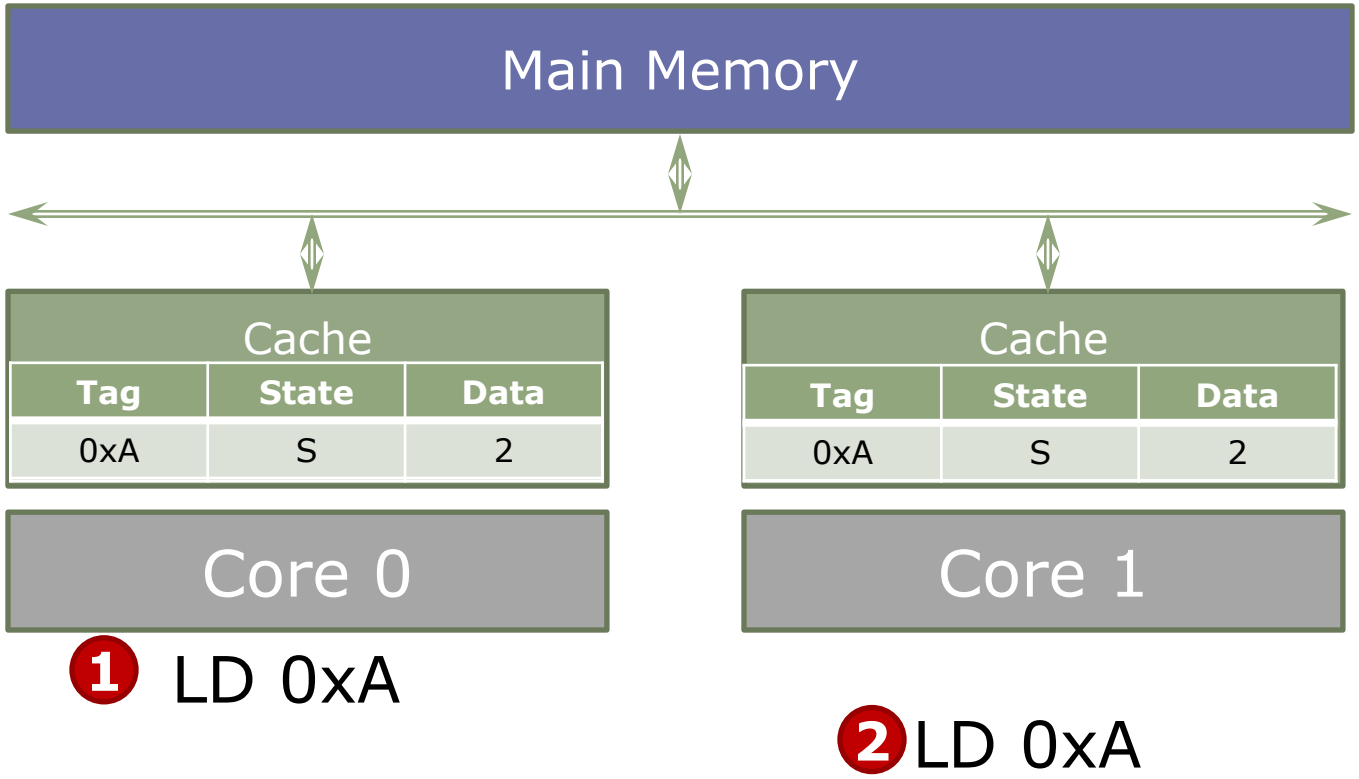

Main Memory
BusRd 0xA
Cache
Tag
State
Data
0xA
S
2
Cache
Tag
State
Data
0xA
S
2
Core 0
Core 1
1
LD 0xA
2
LD 0xA

MSI Example

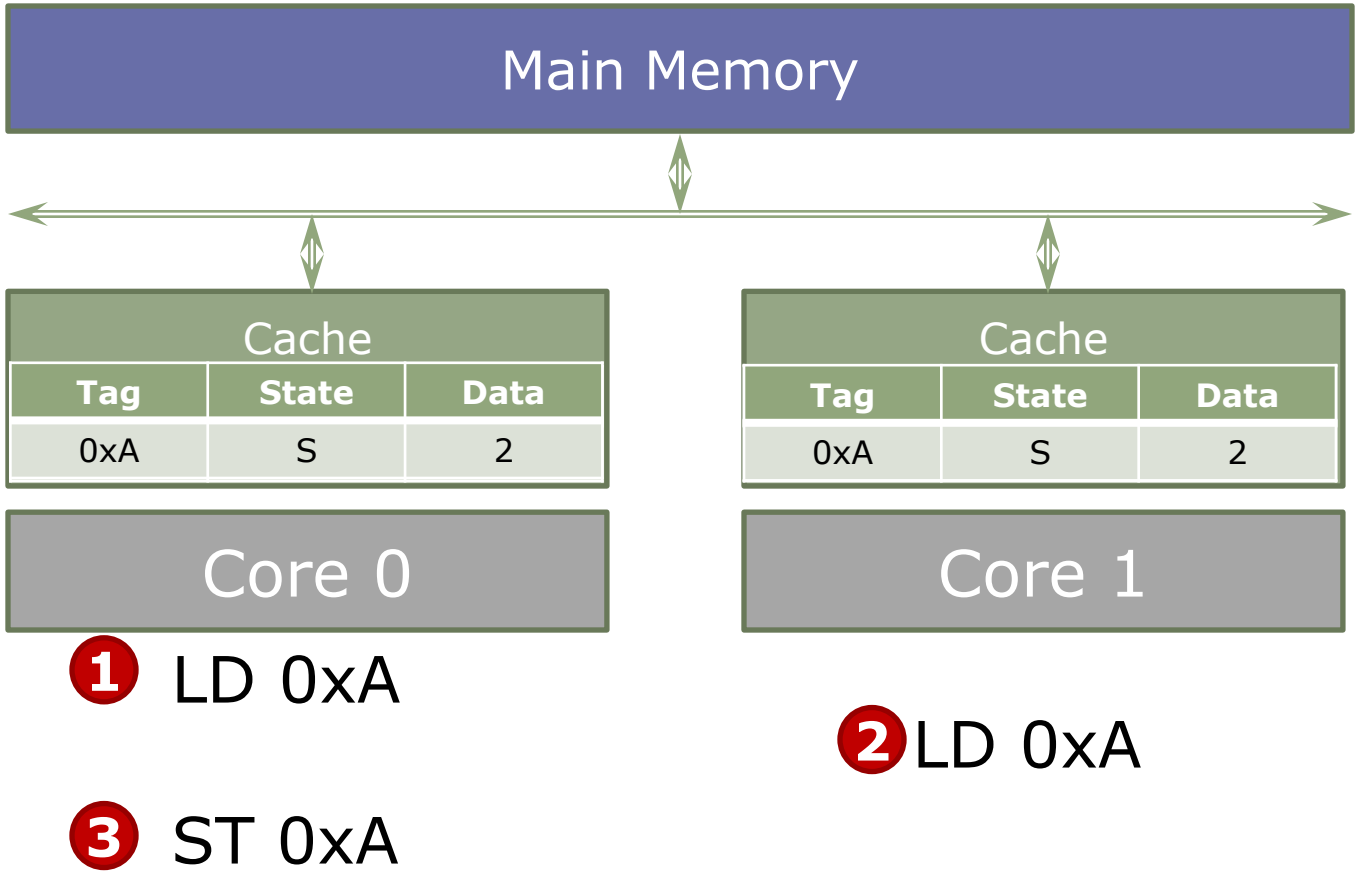


Additional loads satisfied locally, without BusRd (like in VI)

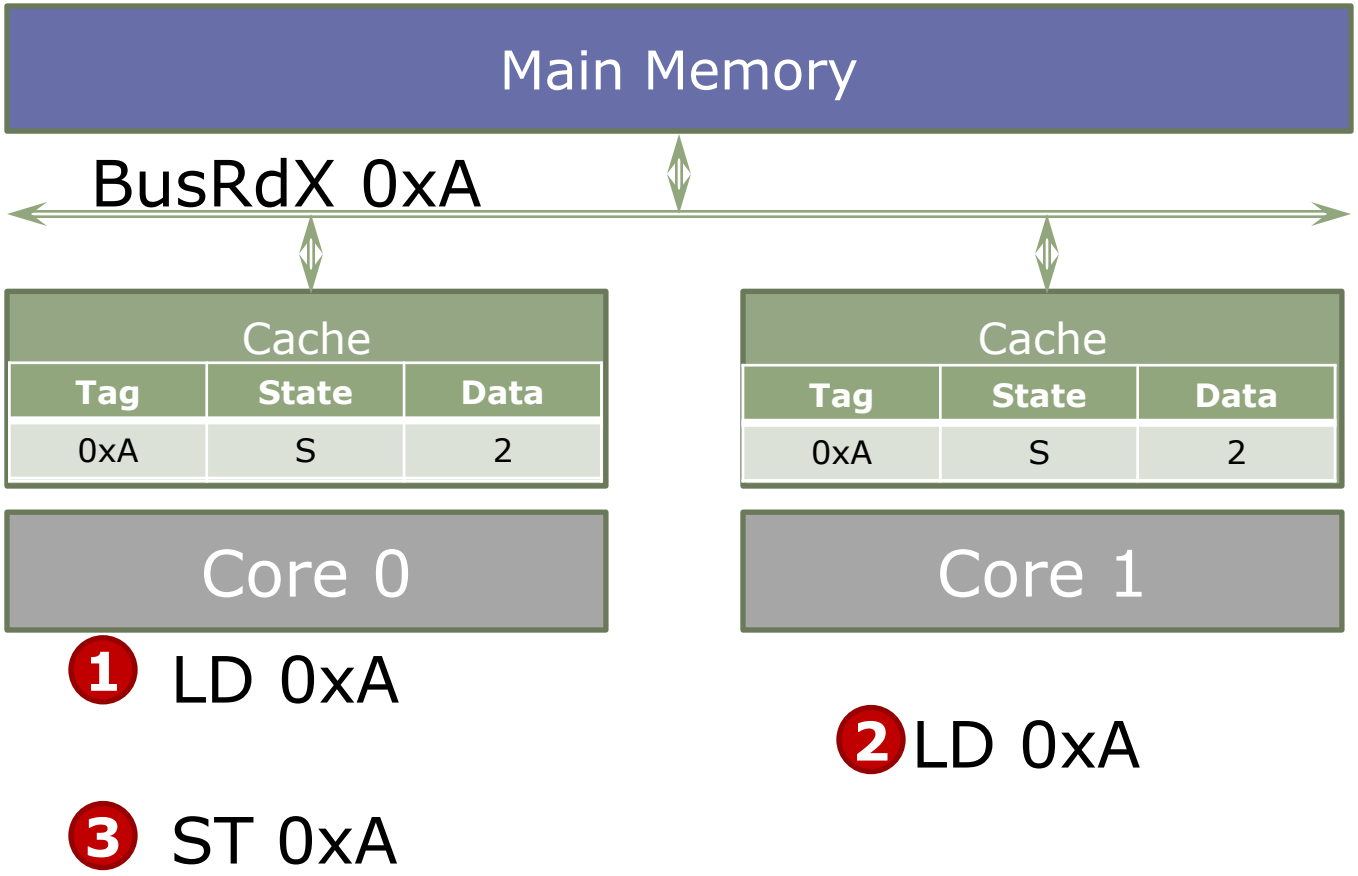
MSI Example


Main Memory
Cache
Tag
State
Data
0xA
S
2
Cache
Tag
State
Data
0xA
S
2
Core 0
Core 1
1
LD 0xA
2
LD 0xA

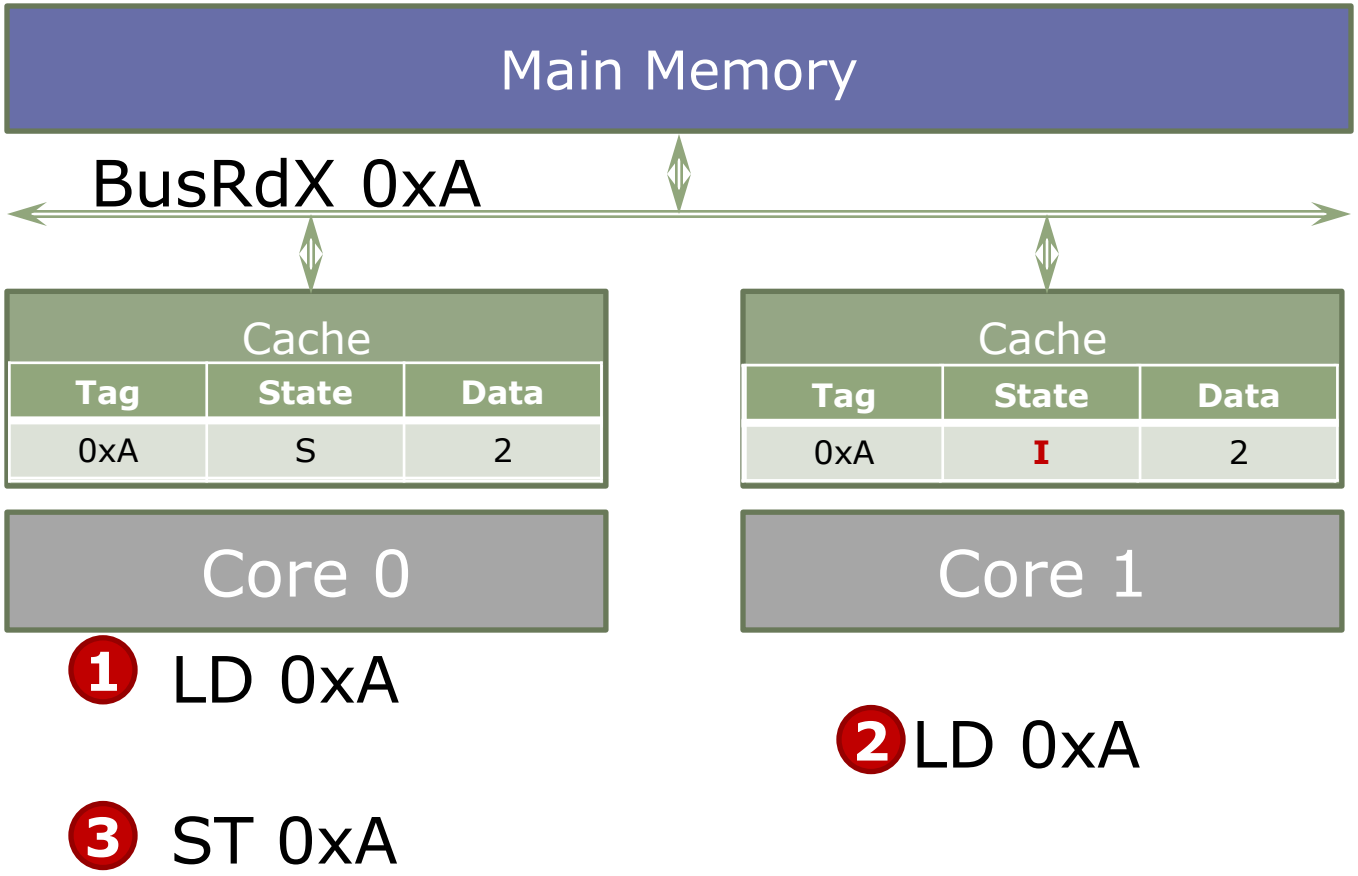
MSI Example


Main Memory
Cache
Tag
State
Data
0xA
S
2
Cache
Tag
State
Data
0xA
S
2
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA

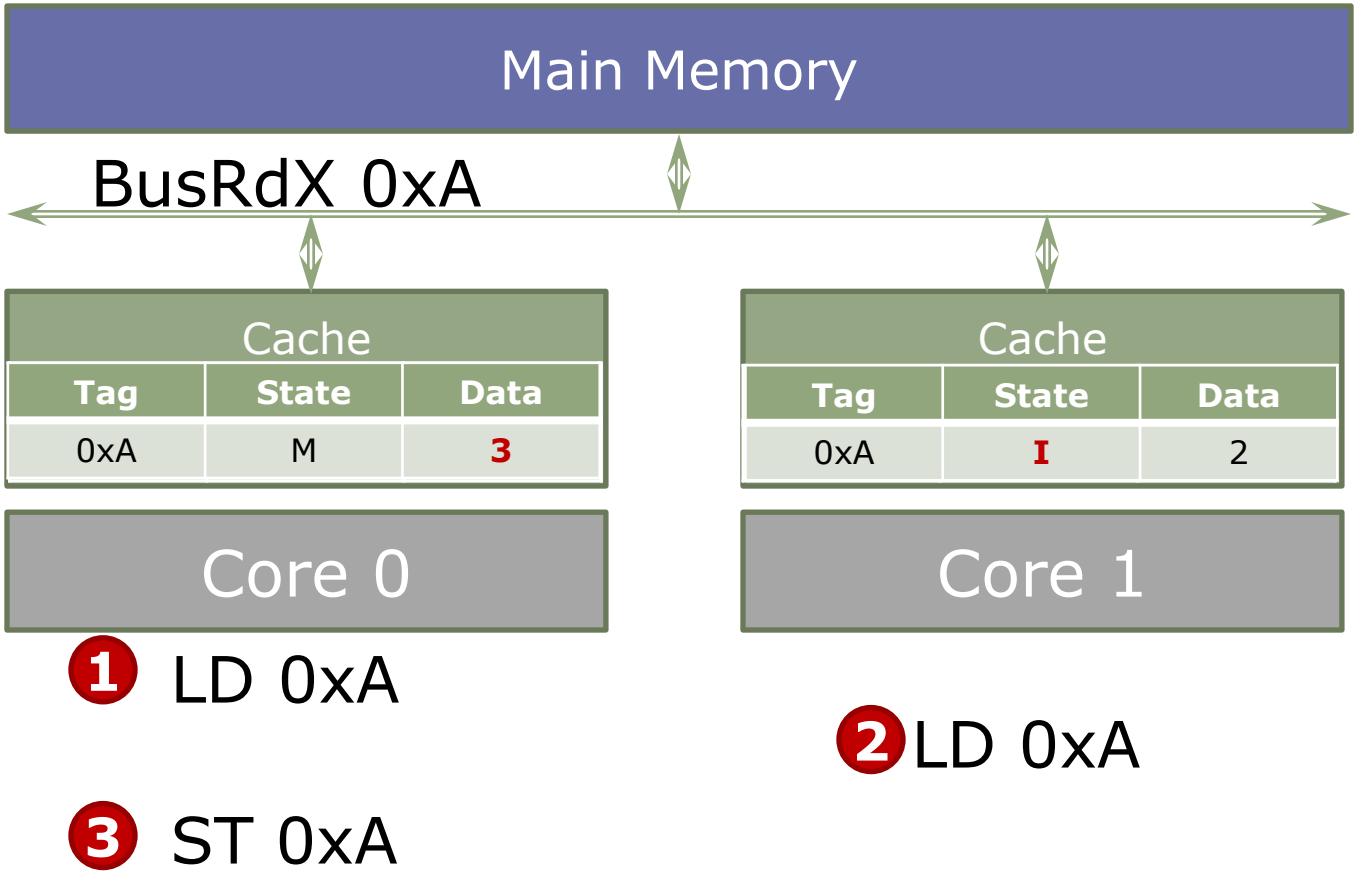
MSI Example


Main Memory
BusRdX 0xA
Cache
Tag
State
Data
0xA
S
2
Cache
Tag
State
Data
0xA
S
2
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA

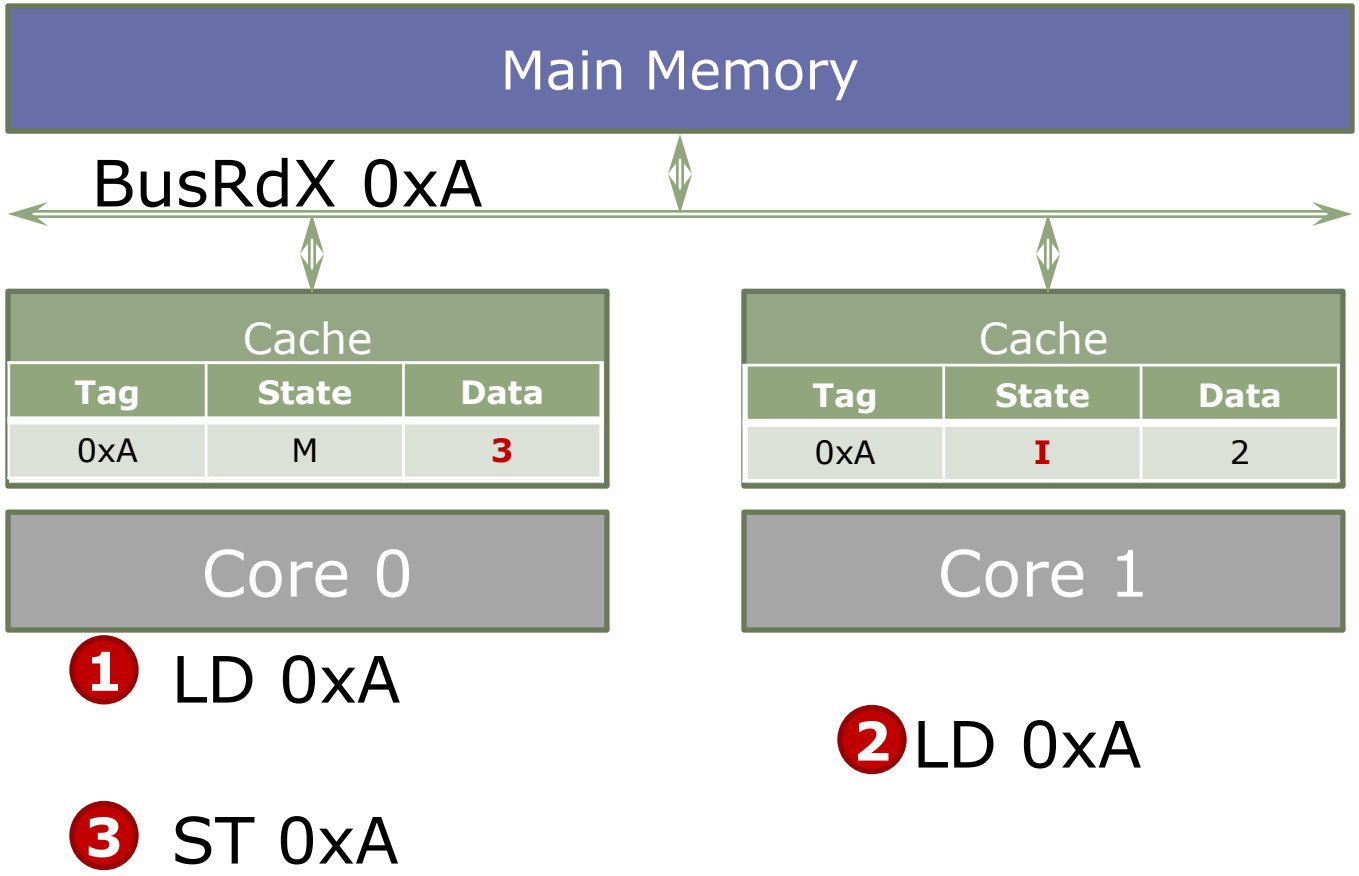
MSI Example


Main Memory
BusRdX 0xA
Cache
Tag
State
Data
0xA
S
2
Cache
Tag
State
Data
0xA
I
2
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA

MSI Example

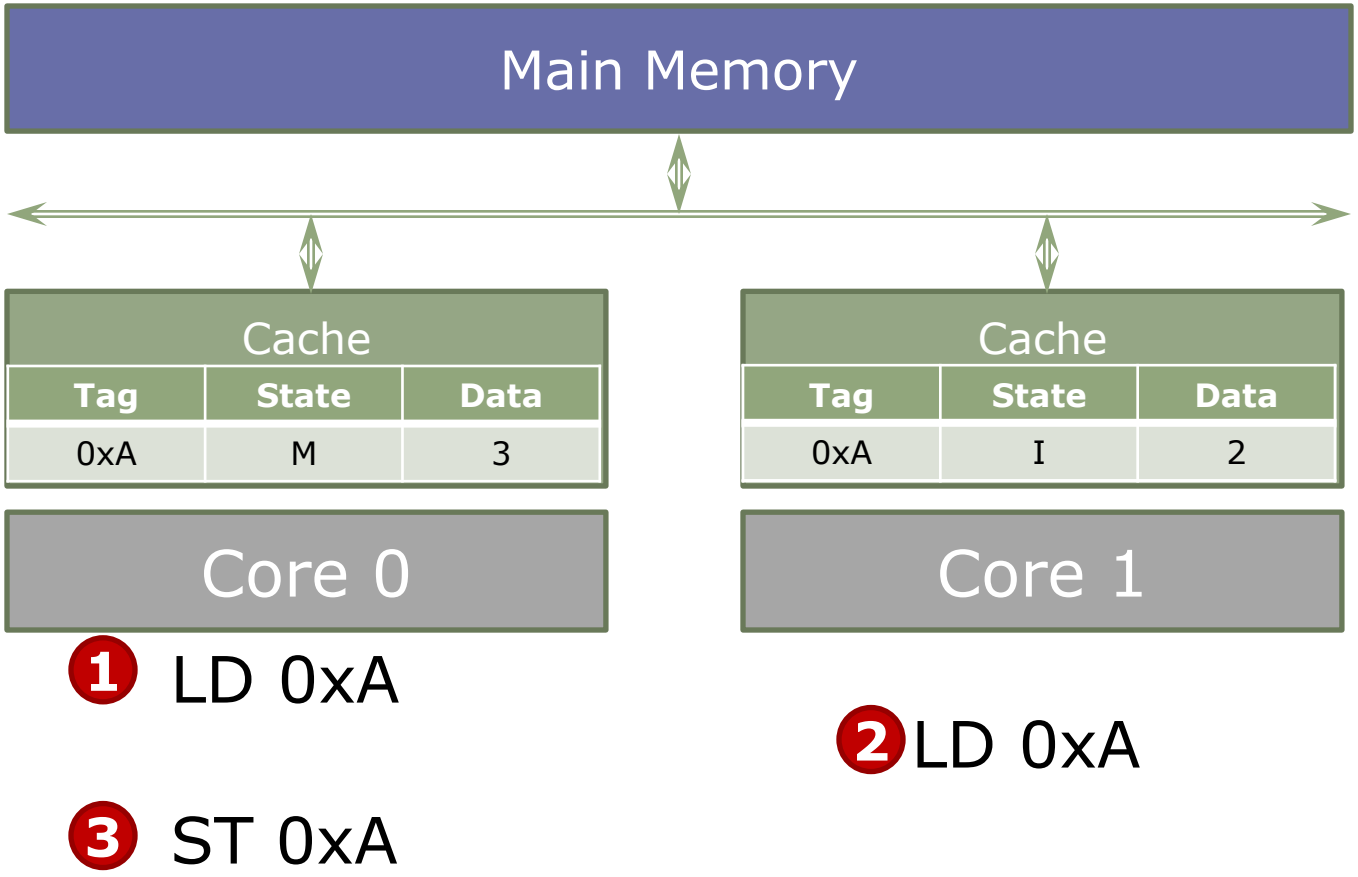

Main Memory
BusRdX 0xA
Cache
Tag
State
Data
0xA
M
3
Cache
Tag
State
Data
0xA
I
2
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA

MSI Example

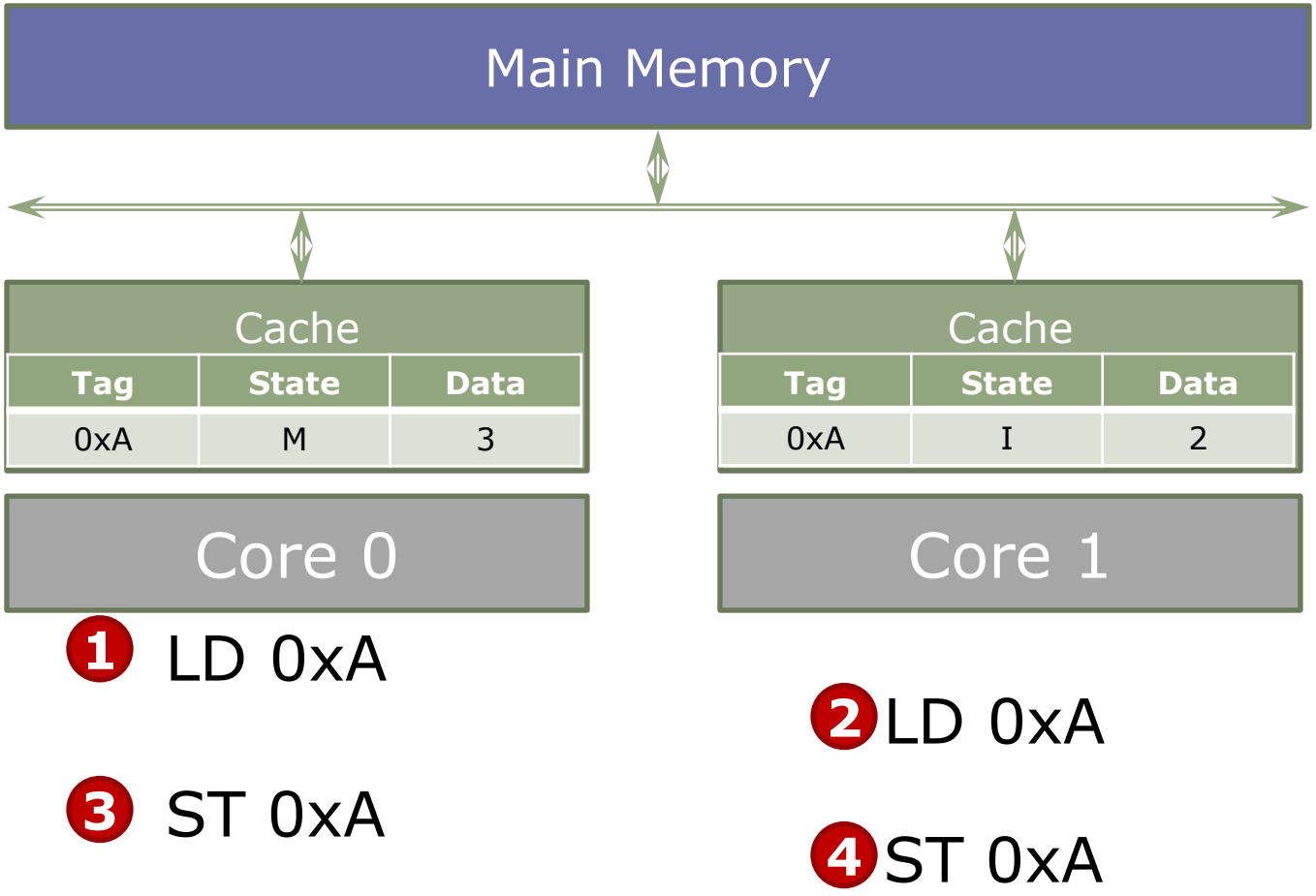


Additional loads *and stores* from core 0 satisfied locally, without bus transactions (unlike in VI)

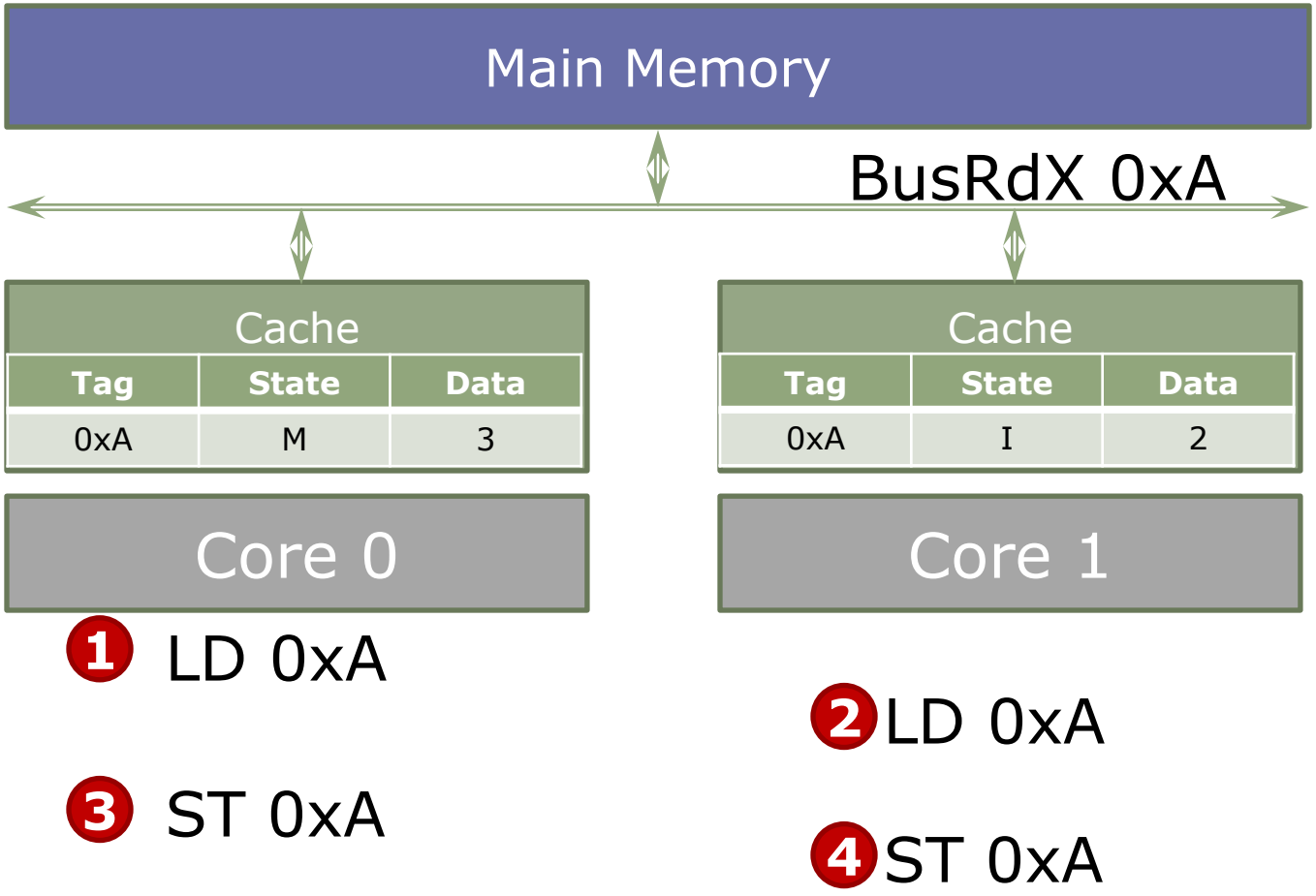
MSI Example


Main Memory
Cache
Tag
State
Data
0xA
M
3
Cache
Tag
State
Data
0xA
I
2
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA

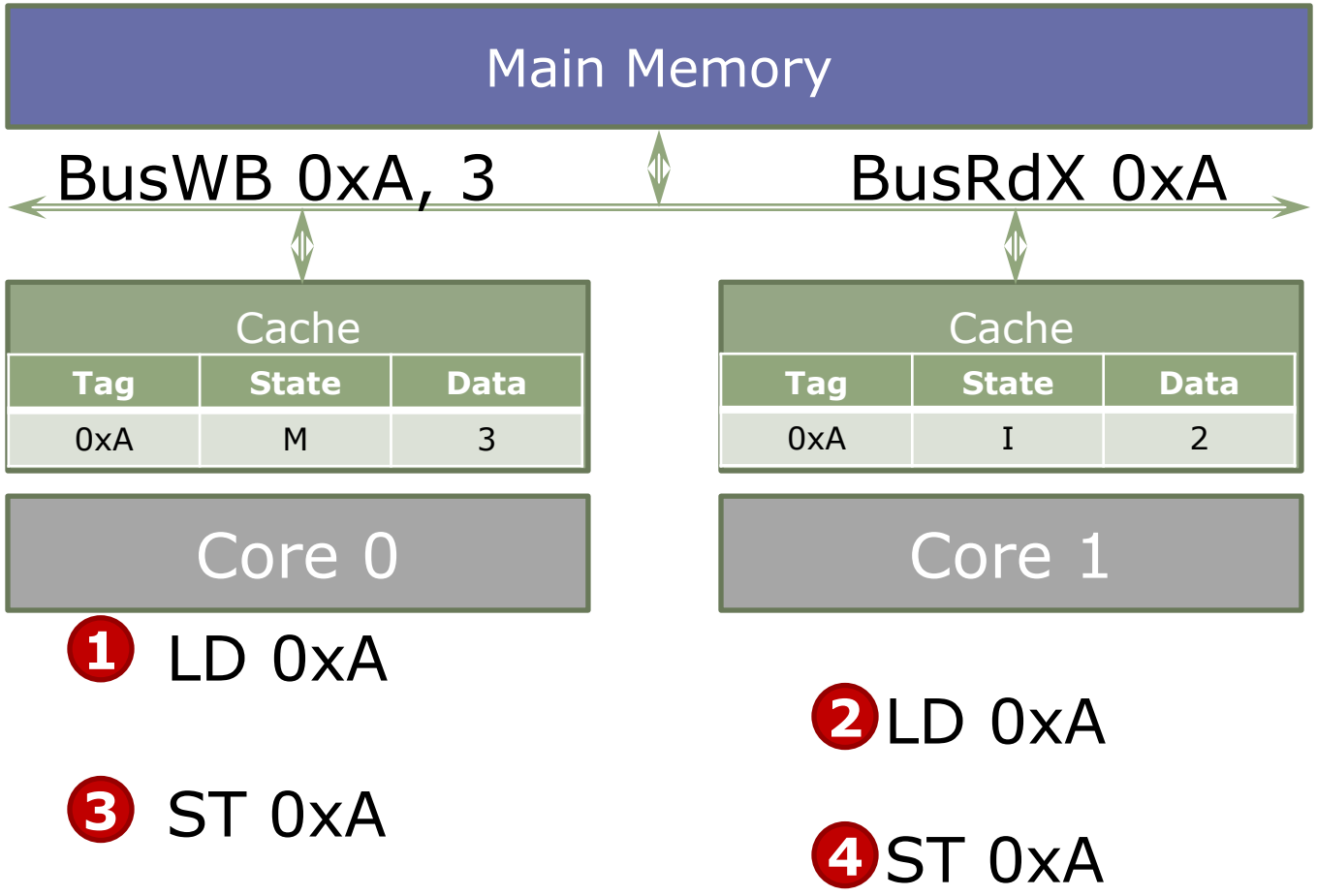
MSI Example


Main Memory
Cache
Tag
State
Data
0xA
M
3
Cache
Tag
State
Data
0xA
I
2
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA
4 ST 0xA

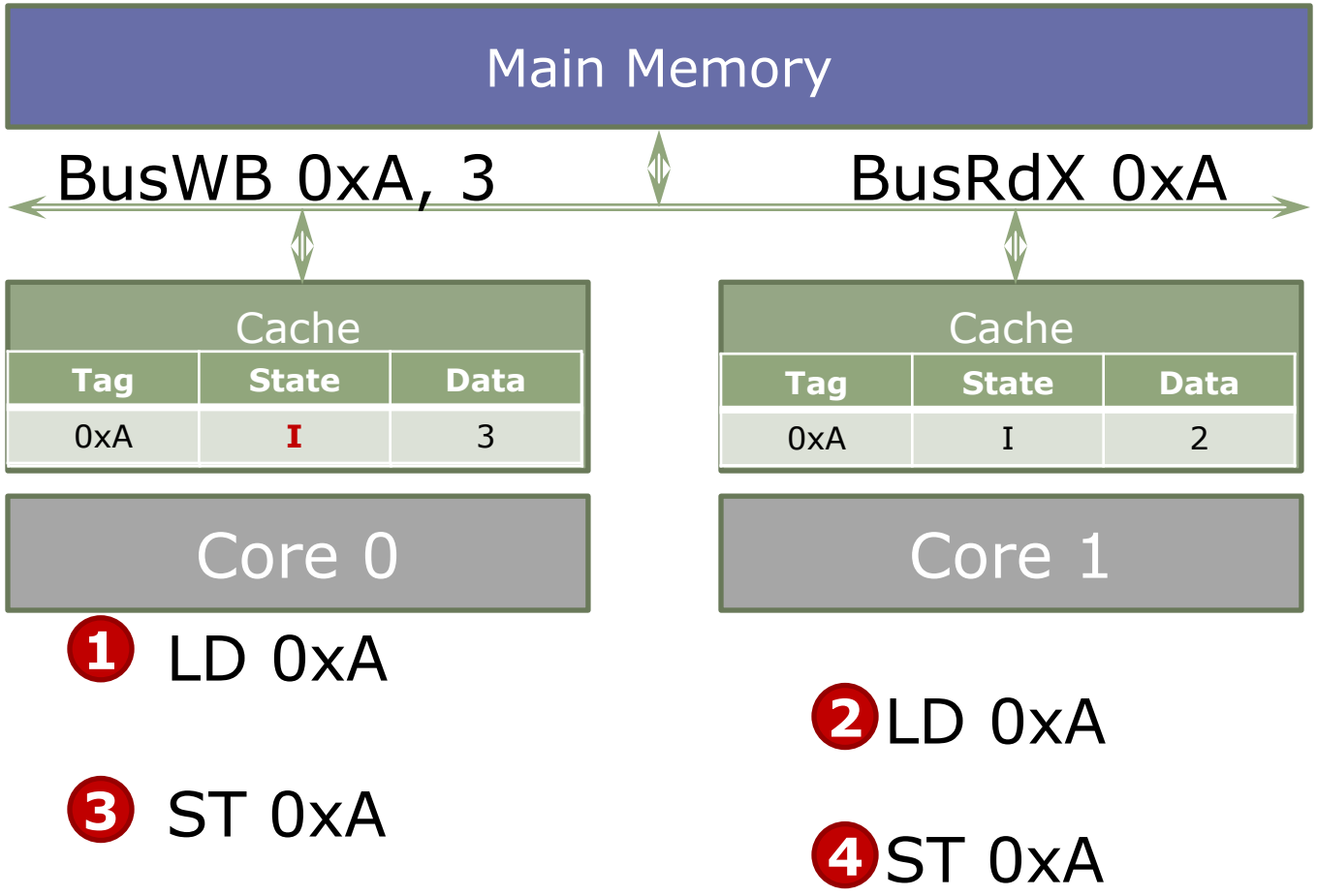
MSI Example


Main Memory
BusRdX 0xA
Cache
Tag
State
Data
0xA
M
3
Cache
Tag
State
Data
0xA
I
2
Core 0
Core 1
1 LD 0xA
3 ST 0xA
2 LD 0xA
4 ST 0xA

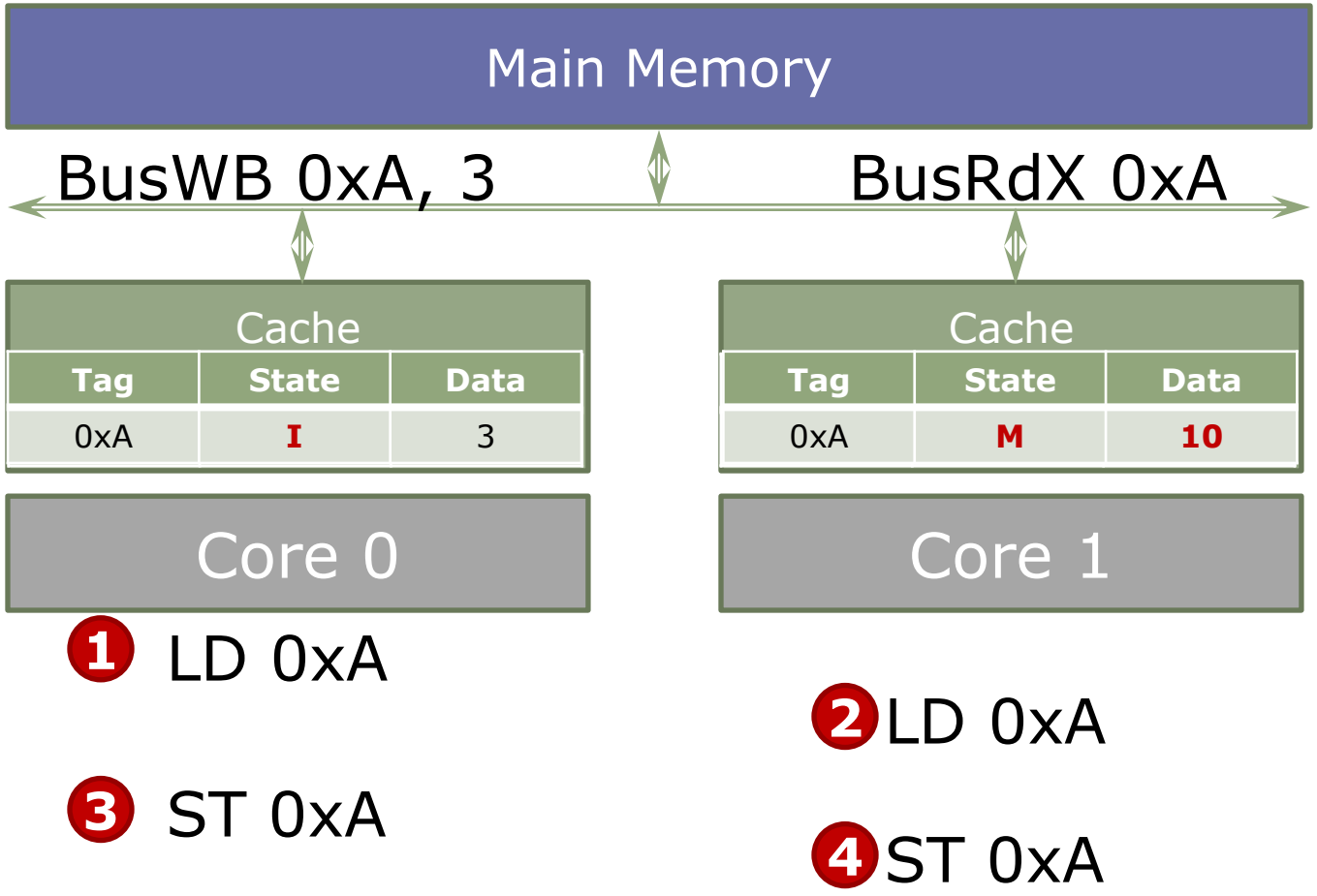
MSI Example


Main Memory
BusWB 0xA, 3
BusRdX 0xA
Cache
Tag
State
Data
0xA
M
3
Cache
Tag
State
Data
0xA
I
2
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA
4 ST 0xA

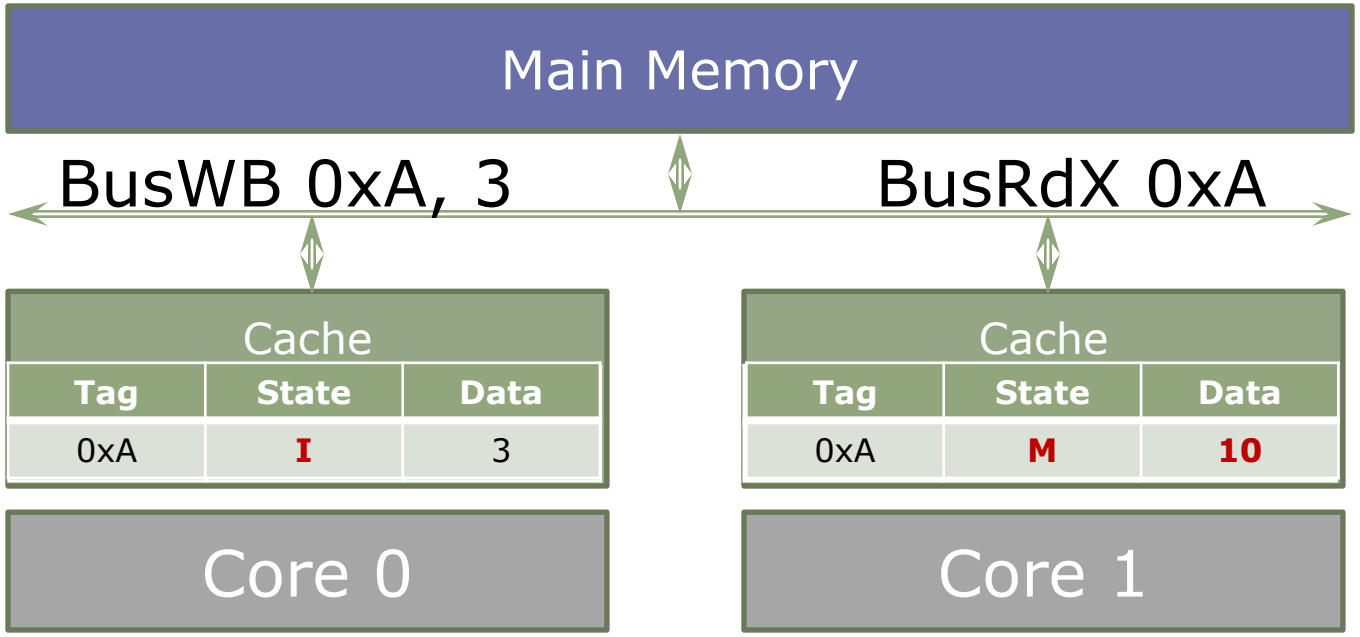
MSI Example


Main Memory
BusWB 0xA, 3
BusRdX 0xA
Cache
Tag
State
Data
0xA
I
3
Cache
Tag
State
Data
0xA
I
2
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA
4 ST 0xA

MSI Example

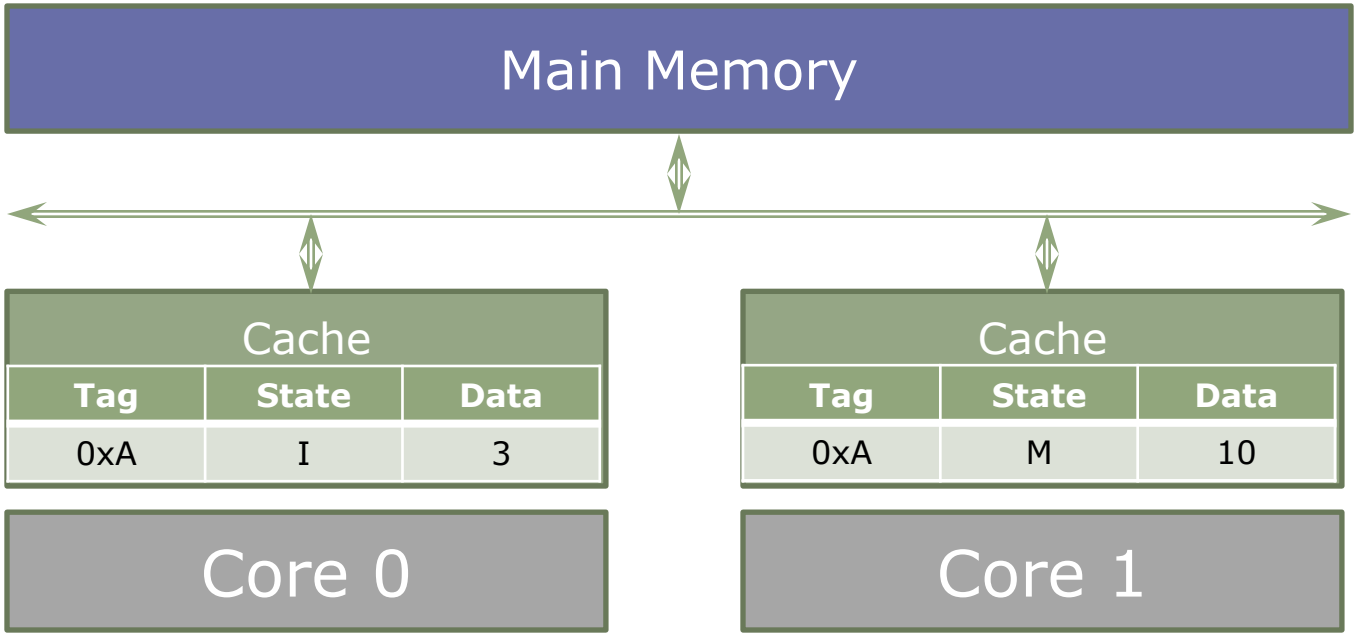

Main Memory
BusWB 0xA, 3
BusRdX 0xA
Cache
Tag
State
Data
0xA
I
3
Cache
Tag
State
Data
0xA
M
10
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA
4 ST 0xA

Cache interventions

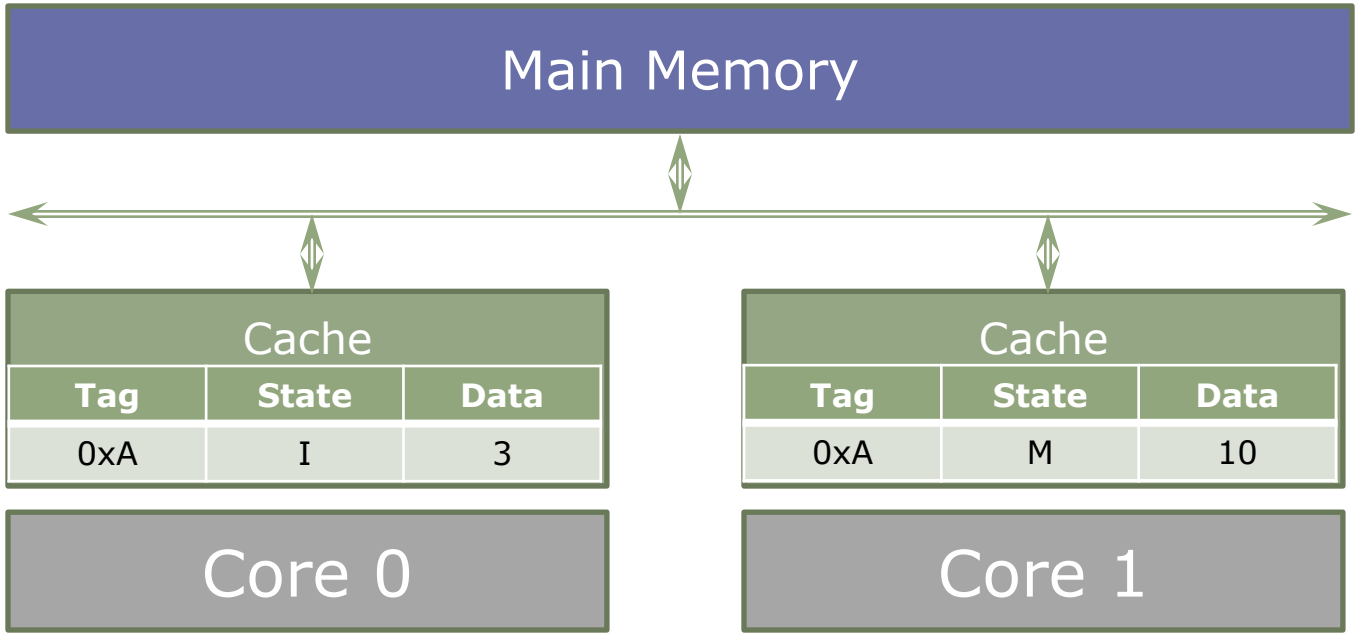


- MSI allows caches to serve writes without updating memory, so main memory can have stale data
  - Core 0's cache needs to supply data
  - But main memory may also respond!
- Cache must override response from main memory

MSI Example

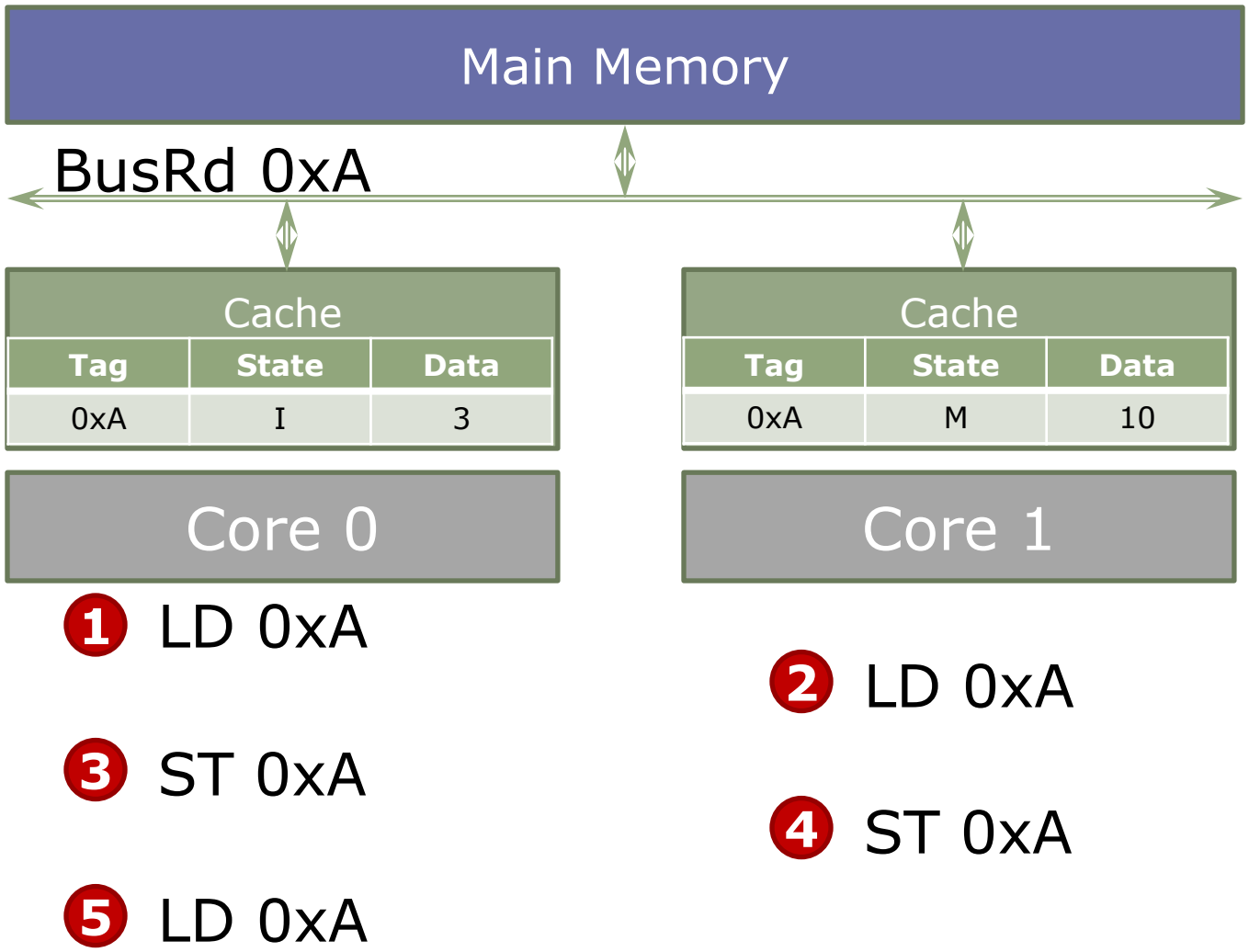

Main Memory
Cache
Tag
State
Data
0xA
I
3
Cache
Tag
State
Data
0xA
M
10
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA
4 ST 0xA

MSI Example

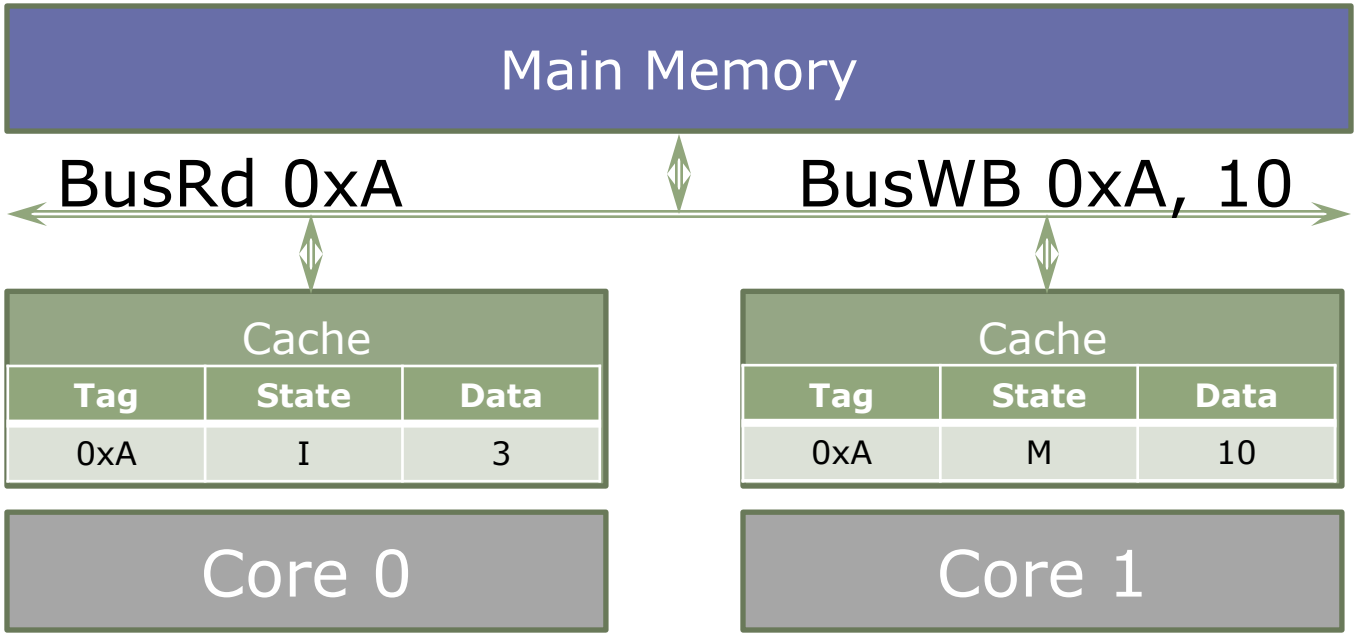


1. LD 0xA
2. LD 0xA
3. ST 0xA
4. ST 0xA
5. LD 0xA

MSI Example

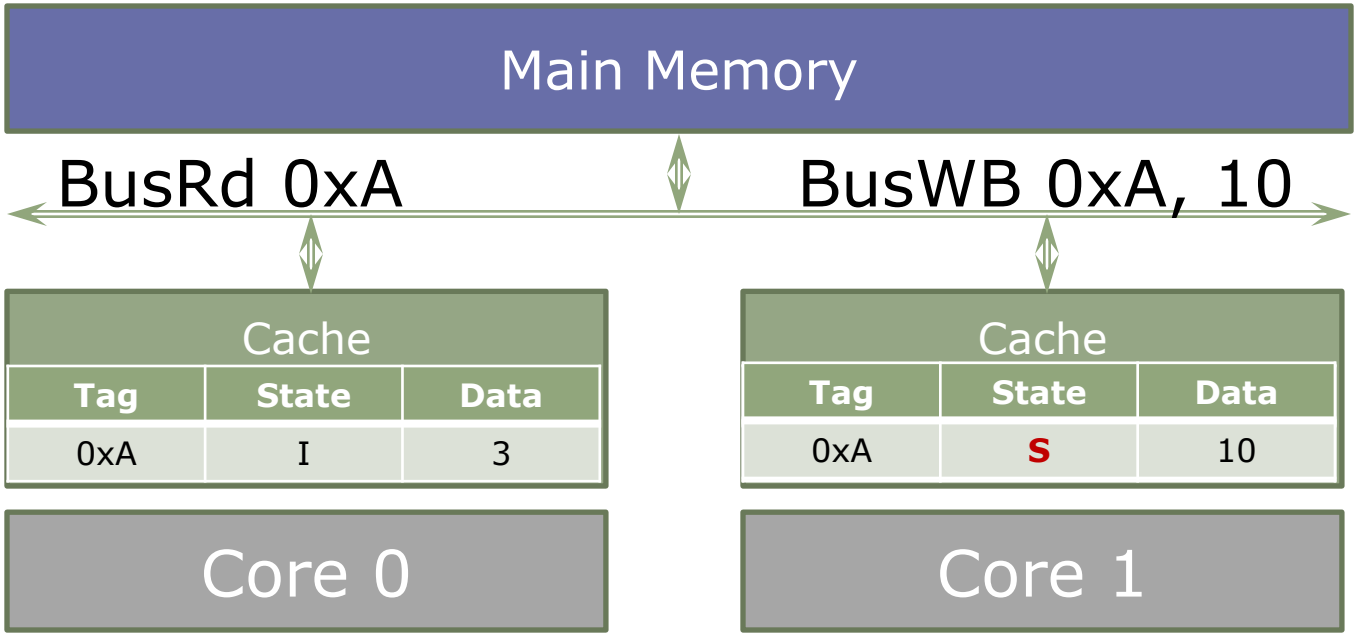

Main Memory
BusRd 0xA
Cache
Tag
State
Data
0xA
I
3
Cache
Tag
State
Data
0xA
M
10
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA
4 ST 0xA
5 LD 0xA

MSI Example



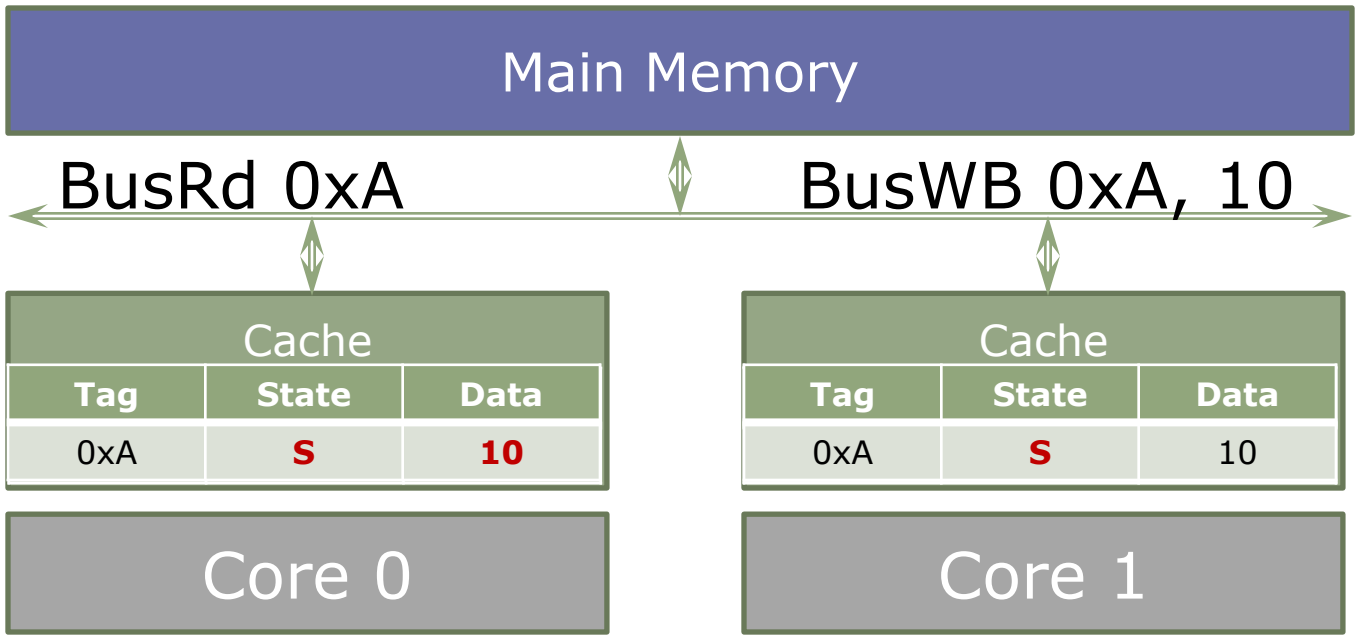
1. LD 0xA
2. LD 0xA
3. ST 0xA
4. ST 0xA
5. LD 0xA

MSI Example



1. LD 0xA
2. LD 0xA
3. ST 0xA
4. ST 0xA
5. LD 0xA

MSI Example


Main Memory
BusRd 0xA
BusWB 0xA, 10
Cache
Tag
State
Data
0xA
S
10
Cache
Tag
State
Data
0xA
S
10
Core 0
Core 1
1 LD 0xA
2 LD 0xA
3 ST 0xA
4 ST 0xA
5 LD 0xA

MSI Optimizations: Exclusive State

- Observation: Doing read-modify-write sequences on private data is common
  - What's the problem with MSI?

MSI Optimizations: Exclusive State

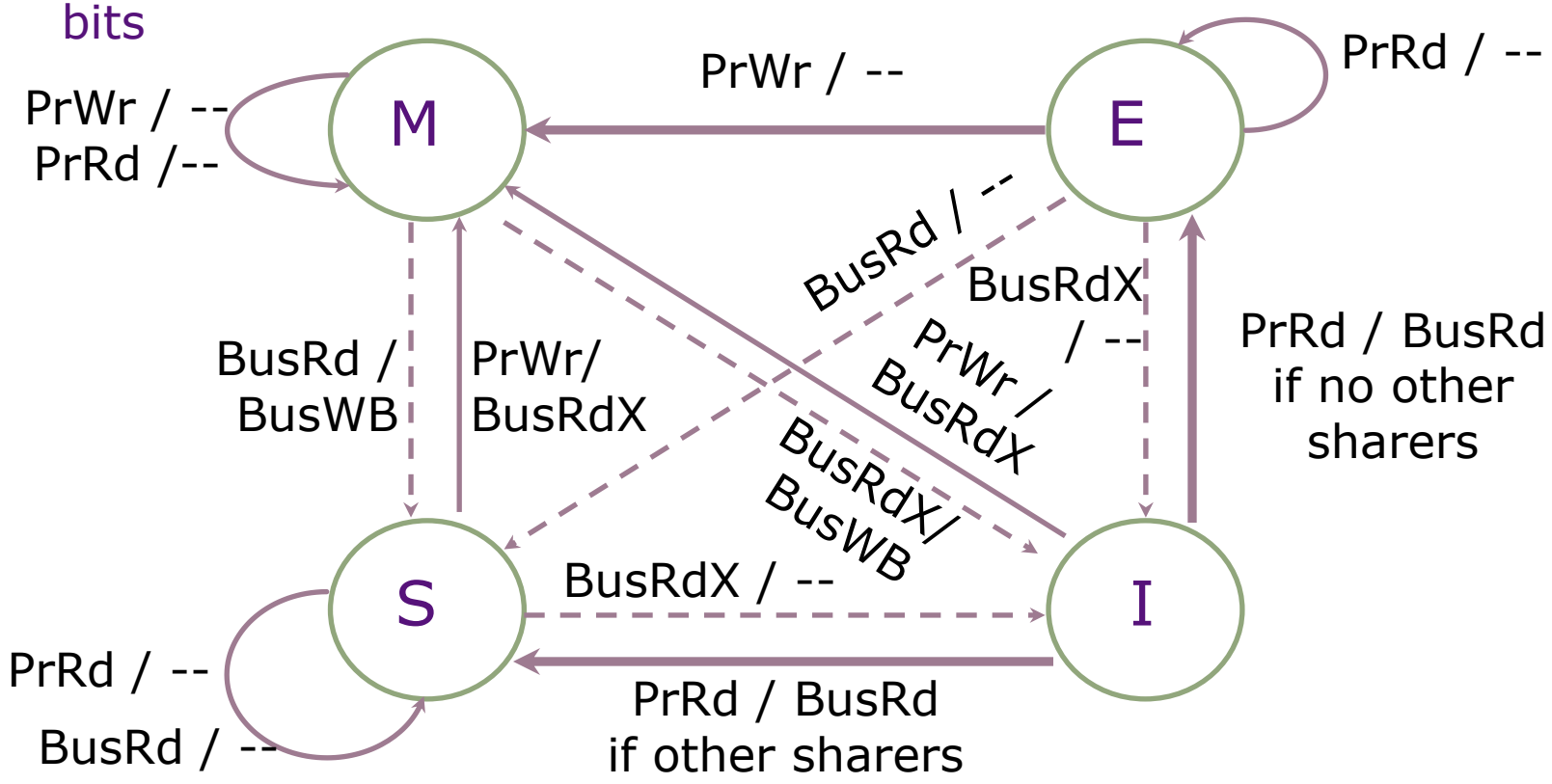
- Observation: Doing read-modify-write sequences on private data is common
  - What's the problem with MSI?
- Solution: E state (exclusive, clean)
  - If no other sharers, a read acquires line in E instead of S
  - Writes silently cause E→M (exclusive, dirty)

MESI: An Enhanced MSI protocol

increased performance for private read-write data

Each cache line has a tag

M: Modified Exclusive
E: Exclusive, unmodified
S: Shared
I: Invalid



MSI Optimizations: Owner State

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  - What happens with frequent read-write sharing?
  - Can we defer the write after S?

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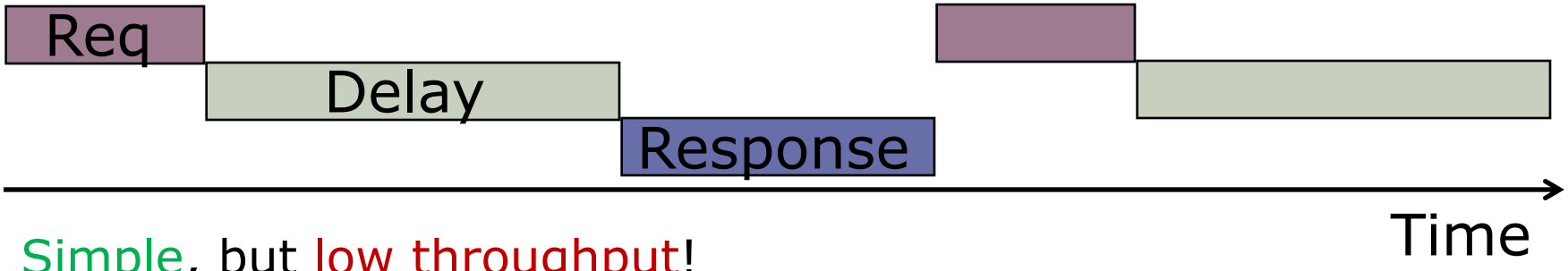
- MSI, MESI, MOSI, MOESI…
  - Typically E if private read-write >> shared read-only (common)
  - Typically O only if writebacks are expensive (main mem vs L3)

Split-Transaction and Pipelined Buses

Atomic Transaction Bus



Simple, but low throughput!

Split-Transaction and Pipelined Buses

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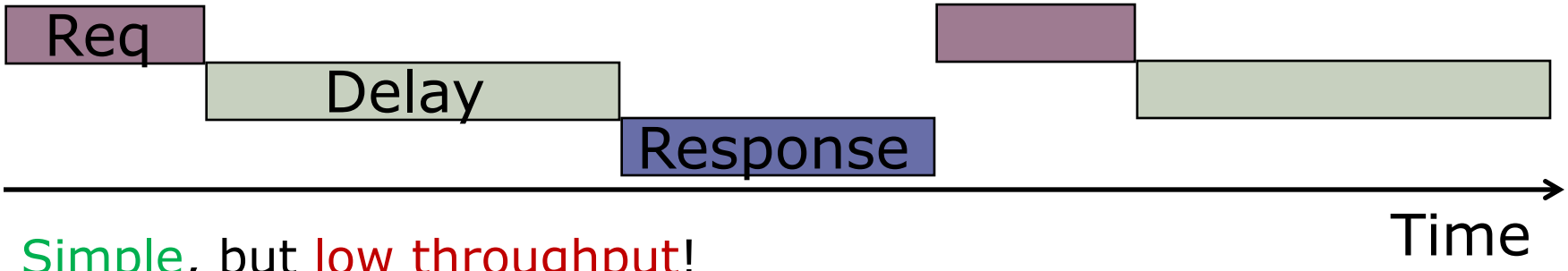

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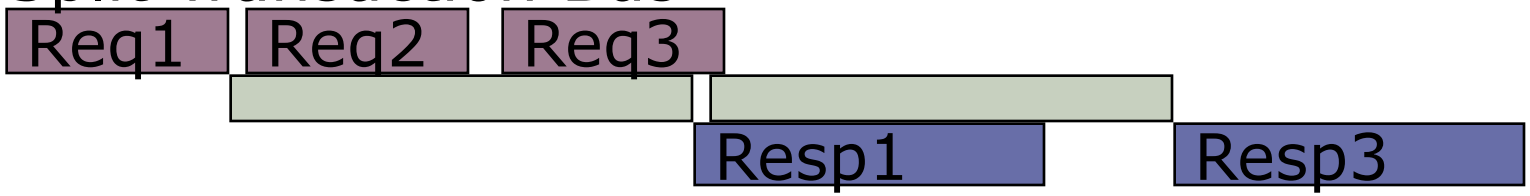
Split-Transaction Bus

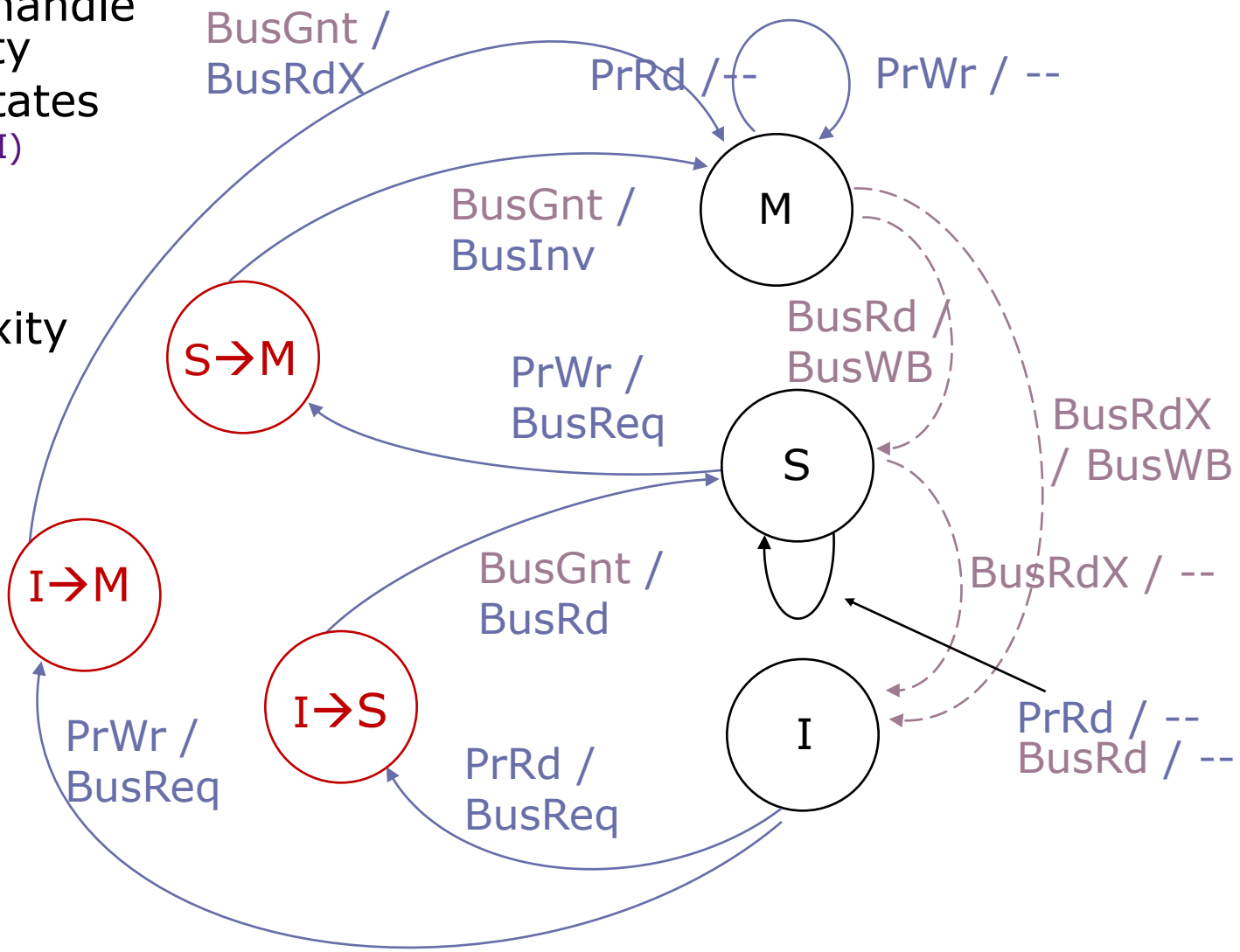


- Supports multiple simultaneous transactions
  - Higher throughput
  - Responses may arrive out of order
- Often implemented as multiple buses (req+resp)

Non-Atomicity → Transient States

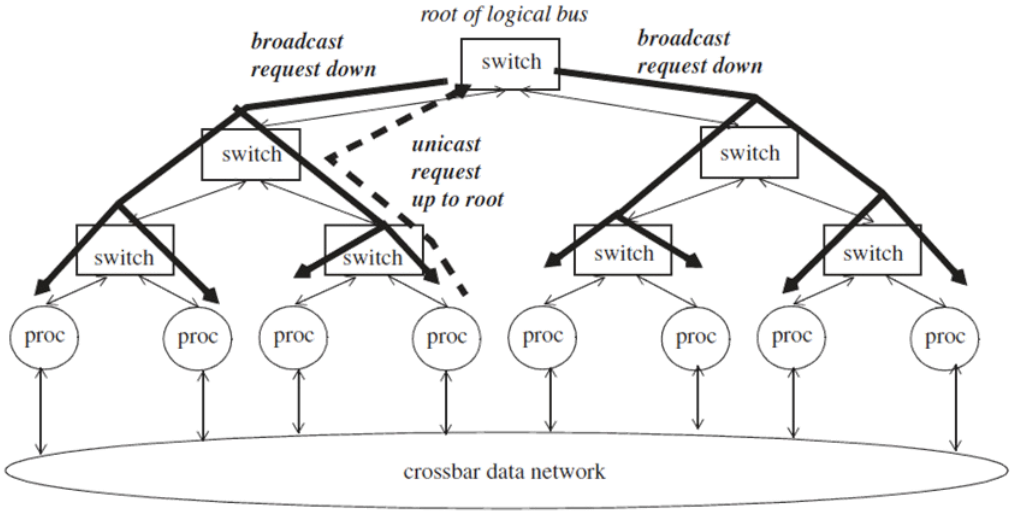
- Protocol must handle lack of atomicity
- Two types of states
  - Stable (e.g. MSI)
  - Transient
- Split + race transitions
- Higher complexity

Actions
Bus Request (BusReq)
Bus Grant (BusGnt)



Scaling Cache Coherence

- Can implement ordered interconnects that scale better than buses…



Starfire E10000 (drawn with only eight processors for clarity). A coherence request is unicast up to the root, where it is serialized, before being broadcast down to all processors

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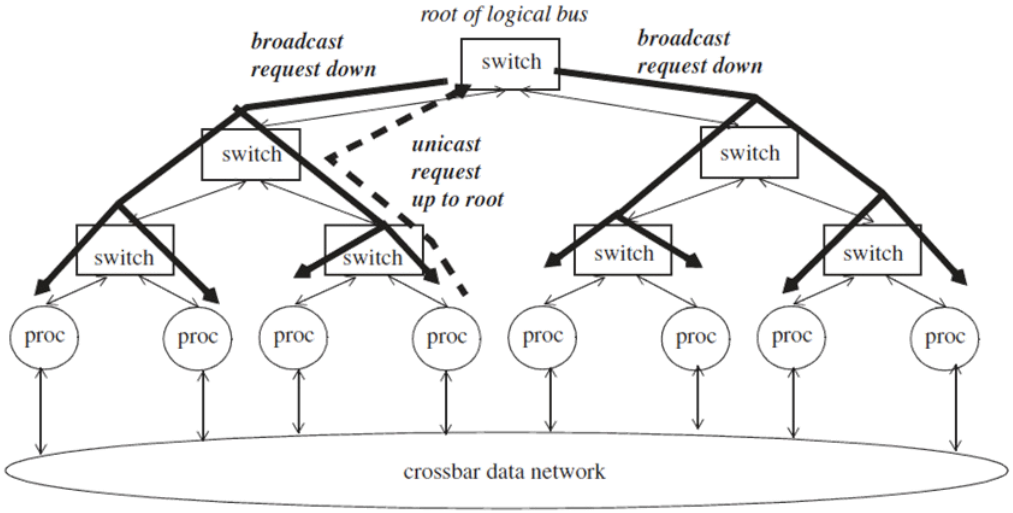
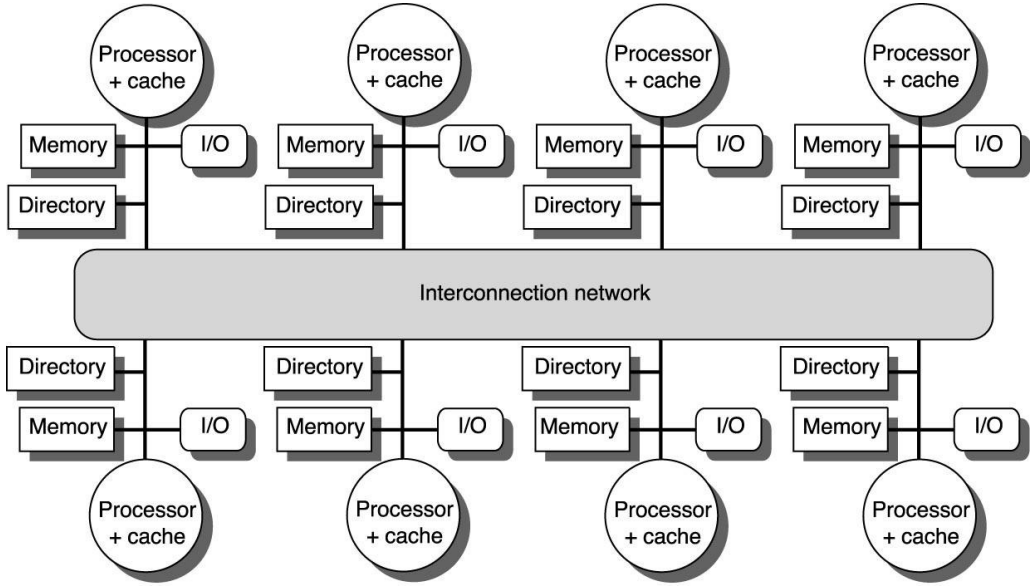


Starfire E10000 (drawn with only eight processors for clarity). A coherence request is unicast up to the root, where it is serialized, before being broadcast down to all processors

- … but broadcast is fundamentally unscalable
  - Bandwidth, energy of transactions with 100s of cache snoops?

Directory-Based Coherence



- Route all coherence transactions through a directory
  - Tracks contents of private caches → No broadcasts
  - Serves as ordering point for conflicting requests → Unordered networks

(more on next lecture)

CC and False Sharing
Performance Issue - 1

A cache block contains more than one word and cache coherence is done at the block-level and not word-level

Suppose P_1 writes $word_i$ and P_2 writes $word_k$ and both words have the same block address.

What can happen?

CC and False Sharing
Performance Issue - 1

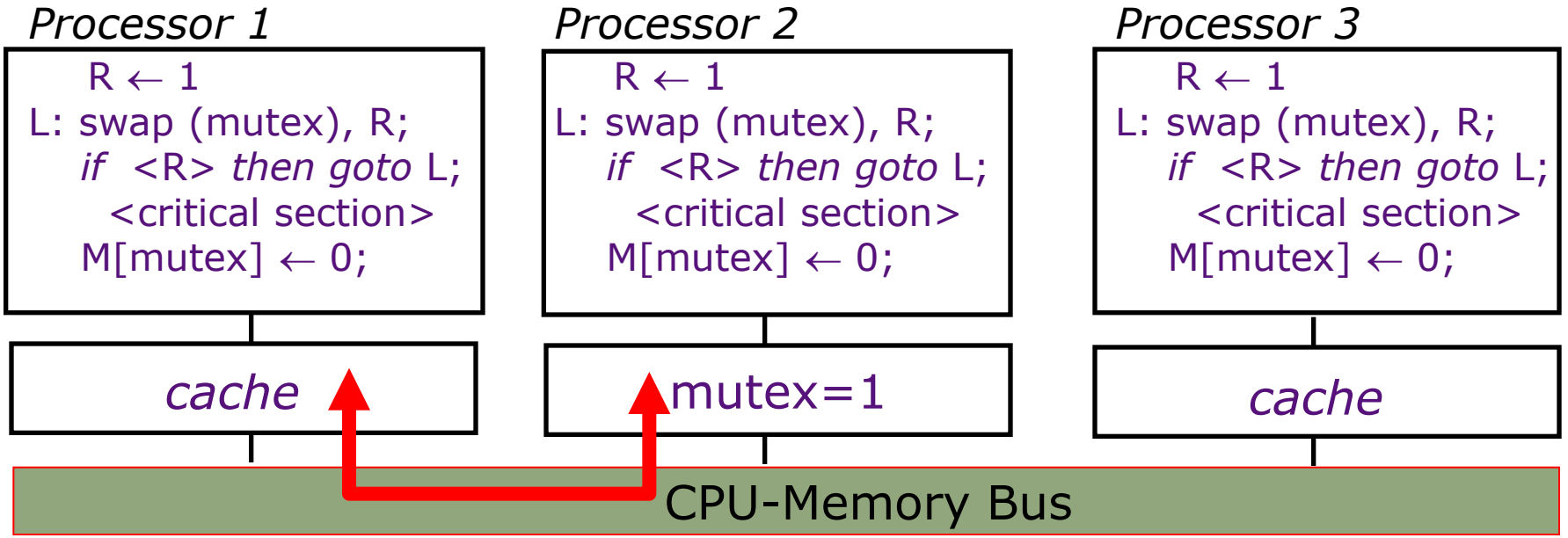
A cache block contains more than one word and cache coherence is done at the block-level and not word-level

Suppose P_1 writes $word_i$ and P_2 writes $word_k$ and both words have the same block address.

What can happen? The block may be invalidated (ping-pong) many times unnecessarily because addresses are in the same block.

CC and Synchronization

Performance Issue - 2



Cache coherence protocols will cause mutex to *ping-pong* between P1's and P2's caches.

Ping-ponging can be reduced by first reading the mutex location *(non-atomically)* and executing a swap only if it is found to be zero (test&test&set).

CC and Bus Occupancy
Performance Issue - 3

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In a multiprocessor setting, bus needs to be locked for the entire duration of the atomic read and write operation

$\Rightarrow$ expensive for simple buses

$\Rightarrow$ *very expensive* for split-transaction buses

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modern processors use

load-reserve

store-conditional

Load-reserve & Store-conditional

Special register(s) to hold reservation flag and address, and the outcome of store-conditional

```
Load-reserve R, (a):
    <flag, adr> ← <1, a>;
    R ← M[a];
```

```
Store-conditional (a), R:
    if <flag, adr> == <1, a>
    then  cancel other procs'
          reservation on a;
          M[a] ← <R>;
          status ← succeed;
    else  status ← fail;
```

If the snooper sees a store transaction to the address in the reserve register, the reserve bit is set to 0

- Several processors may reserve 'a' simultaneously
- These instructions are like ordinary loads and stores with respect to the bus traffic

Performance:
Load-reserve & Store-conditional

The total number of memory (bus) transactions is not necessarily reduced, but splitting an atomic instruction into load-reserve & store-conditional:

- *increases bus utilization* (and reduces processor stall time), especially in split-transaction buses
- *reduces cache ping-pong effect* because processors trying to acquire a mutex do not have to perform stores each time