

Quiz 1 Review

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Quiz 1 logistics

- Time: 1pm-2:30pm on Friday, March 6
- Location: 32-141
- Covered materials: L01-07
- Handouts will be provided
 - Handout 1 – EDSACjr
 - Handout 3 – RISC MIPS
 - Handout – SRRIP
 - Handout – Interleaved Memory
- No calculators

Pipelining

- Overlaps execution of multiple instructions
- Visualization
 - Instruction flow diagram
 - Resource usage diagram
- Hazard: an instruction cannot execute because
 - Resource is not ready: structural hazard
 - Data value is not ready: data hazard
 - PC is not ready: control hazard

Strategies to resolve hazards

- Stall
- Bypass
- Speculate
- More to cover in the coming lectures

Caches

- A small but fast storage that exploits locality
 - Temporal locality
 - Spatial locality
- Performance metrics
 - $AMAT = \text{hit time} + \text{miss rate} * \text{miss penalty}$
- Design options
 - # of sets, # of ways
 - Block size
 - Replacement policy
 - Inclusivity and exclusivity
 - More to cover in the future lectures

Interpreting a byte address



Used to distinguish lines mapped to the same set

Used to locate a set

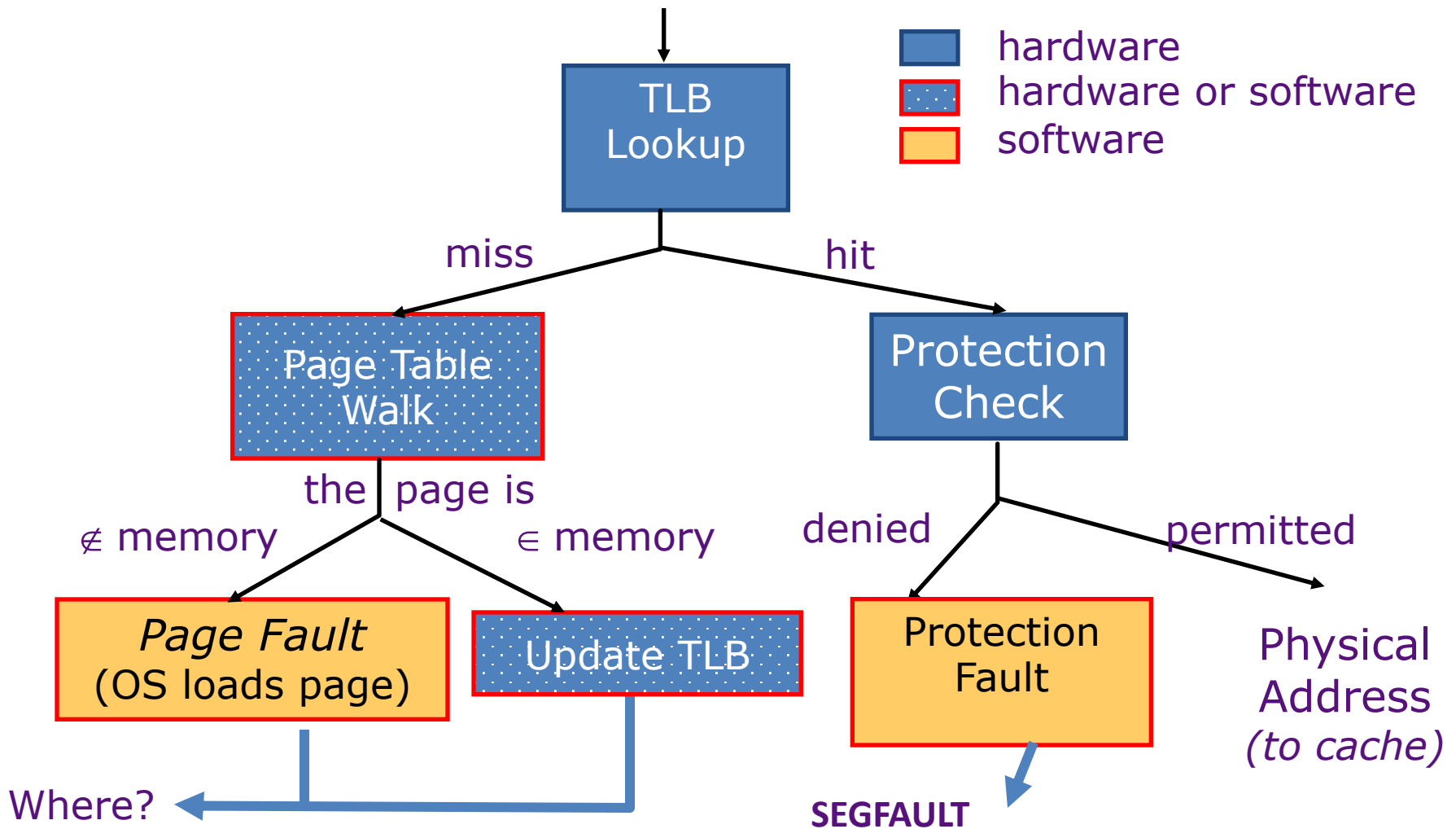
Used to locate a word/byte within the line

Memory management strategies

- Segmentation
- Paging

Address translation in paging

Virtual Address

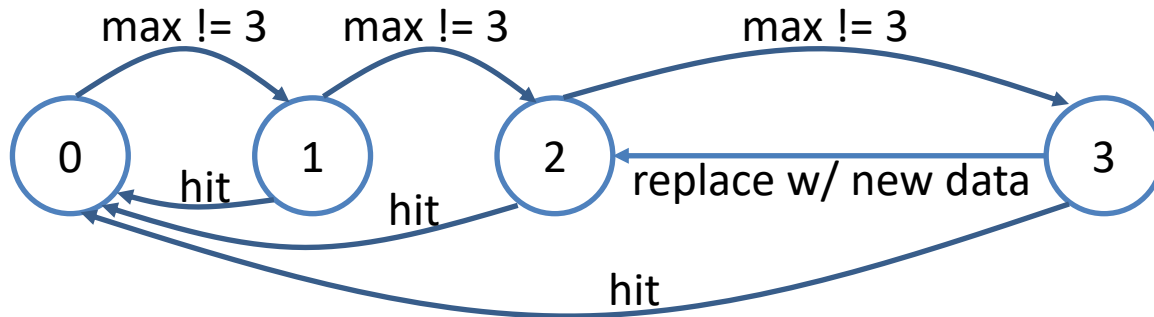


Self-modifying code

- Simplifies data path design but complicates programming

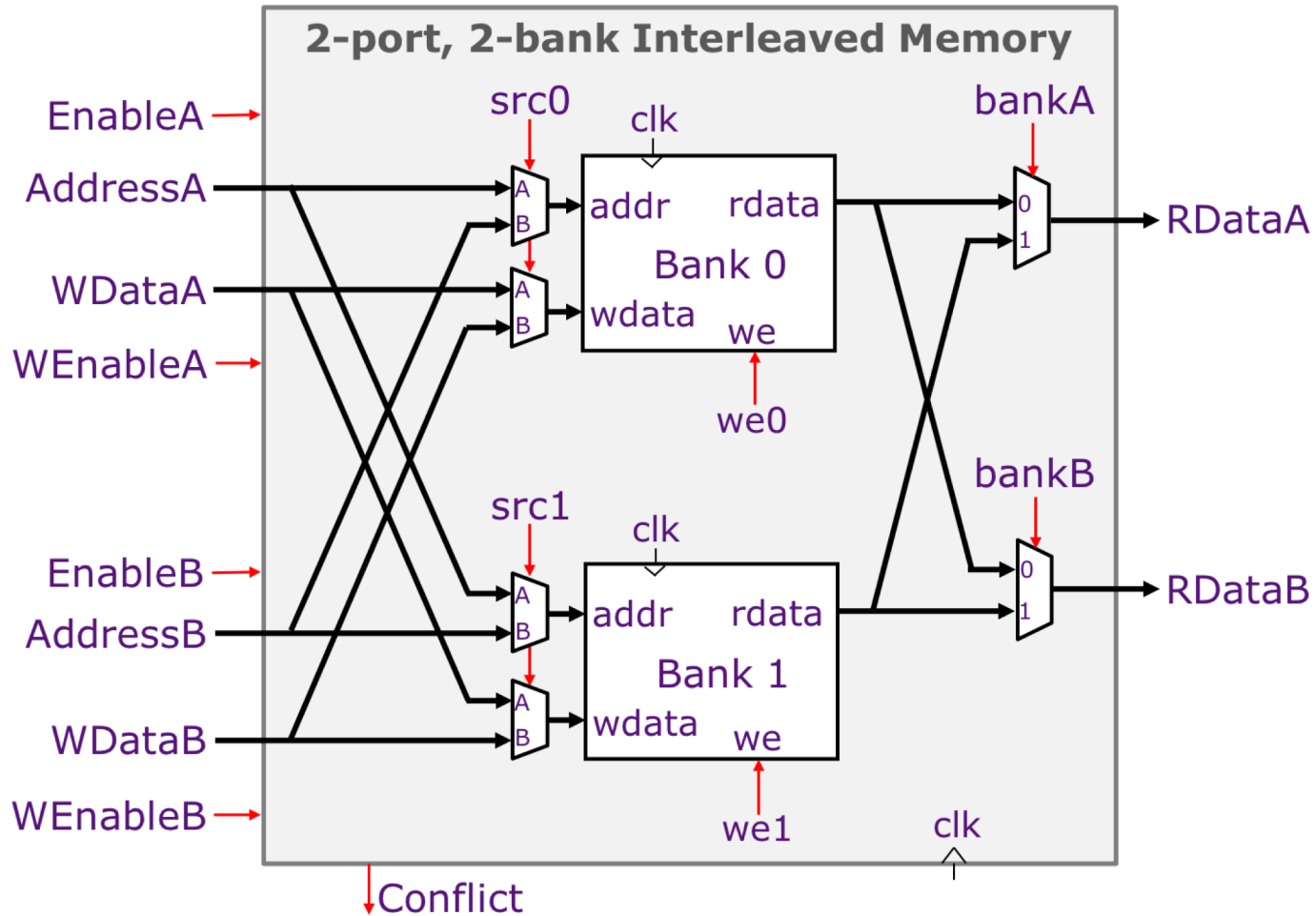
Handout - SRRIP

```
// Candidates and their RRPVs are indexed by w in [0, W - 1]
if access is a hit on way w:
    RRPVs[w] = 0
else: // access is a miss
    while (maximum value of RRPVs != 2^M - 1): // 3 if M == 2
        Increment all RRPVs by 1
    // Now there must be at least an RRPV with value 2^M - 1
    Select w as the minimum index s.t. RRPVs[w] == 2^M - 1
    // w-th candidate is selected as the victim
    Replace the old data of candidate w with the new data
    RRPVs[w] = 2^M - 2 // 2 if M == 2
```



State-transition diagram of an RRPV in 2-bit SRRIP

Handout – Interleaved Memory



Wish you all the best!