### 6.5930/1

Hardware Architectures for Deep Learning

## Sparse Architectures - Part 1

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Joel Emer and Vivienne Sze

Massachusetts Institute of Technology Electrical Engineering \& Computer Science

## Goals of Today's Lecture

- Last lecture, we discussed how to make weights and activations of DNN models sparse
- Sparsity of DNNs on the order of 30-70\%, while existing software libraries (e.g., sparse BLAS) designed for $>99 \%$
- Need specialized hardware to exploit!
- Today and in the next lecture, we will discuss how to translate sparsity into reductions in energy consumption and processing cycles
- First, discuss the representation of sparse data
- Second, present some architectures that exploit sparsity

Resources: Course notes - Chapter 8.2 and 8.3

## Many problems use Sparse Tensors


[Hegde, et.al., ExTensor, MICRO 2019]

## Motivation

- Leverage CNN sparsity to improve energy-efficiency

[Parashar, et.al., SCNN, ISCA 2017]


## Aspects of Scheduling - Sparsity

## Gating:

Explicitly eliminate ineffectual storage accesses and computes by letting the hardware unit staying idle for the cycle to save energy


## CONV Layer



## Tensors

## Rank-0 - Scalar



Rank-2 - Matrix


## Rank-1 - Vector



## Tensor Data Terminology



- The elements of each "rank" (dimension) are identified by their "coordinates", e.g., rank H has coordinates $0,1,2$
- Each element of the tensor is identified by the tuple of coordinates from each of its ranks, i.e., a "point". So (1,2) -> "f"


## Tree-based Tensor Abstraction



## Tree-based Tensor Abstraction



## Fibertree Tensor Abstraction



## Fibertree Tensor Abstraction

Finding point $(2,1)$


## Fibertree Tensor Abstraction

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## Information in a Fiber

- Each fiber has a set of (coordinate, "payload") tuples



## Information in a Fiber

Method: maybe(payload) $=$ fiber.getPayload(coordinate $)$


## Example Fiber Representations

Each fiber has a set of (coordinate, "payload") tuples
Array


Coordinate/Payload List


Data in a fiber is accessed by its position or offset in memory

## Fiber Representation Choices

- Implicit Coordinates
- Uncompressed (no metadata required)
- Compressed - e.g., run length encoded
- Explicit Coordinates
- E.g., coordinate/payload list
- Compressed vs Uncompressed
- Compressed/uncompressed is an attribute of the representation*.
- Uncompressed means size is proportional to maximum coordinate value
- Compressed formats will have metadata overhead relative to uncompressed formats. For dense data, this may cost more than just using an uncompressed format.
- Space efficiency of a representation depends on sparsity
*Note: sparsity/density is an attribute of the data.


## Implicit Coordinates: RLE

## Example

Input: 0, 0, 12, 0, 0, 0, 0, 53, 0, 0, 22
Method 1: Run Length Coding
Rather than send zero, send "run length" of zeros
e.g., 5 bits for run length and 16 bits for non-zero value

Output: 2, 12, 4, 53, 2, 22
5b 16b

Total Number of Bits:

## Implicit Coordinates: Significance Map

Example
Input: 0, 0, 12, 0, 0, 0, 0, 53, 0, 0, 22
Method 2: Significance Map Coding
(a variant of this is referred to as bitmask coding)
Send one bit to indicate if significant (i.e., non-zero); if significant, send 16 bits for non-zero value


Total Number of Bits:

How does this compare to Run Length Coding?

## Implicit Coordinates: Huffman Encoding

## Example

Input: 0, 0, 12, 0, 0, 0, 0, 53, 0, 0, 22

## Method 3: Huffman Coding

Assign number of bits based on probability of occurrence
Message Codeword Probability

| $a_{1}$ | 0 |
| :--- | :---: |
| $a_{2}$ | 100 |
| $a_{3}$ | 110 |
| $a_{4}$ | 1110 |
| $a_{5}$ | 101 |
| $a_{6}$ | 1111 |



Assign codewords directly to values or to values and run-lengths

## Quantization and Compression

- Quantization + Significance Map Coding

Example:
Value: $16^{\prime}$ b0 $\rightarrow$ Compressed Code: $\left\{1^{\prime} \mathrm{b} 0\right\}$
Value: $16^{\prime} \mathrm{bx} \rightarrow$ Compressed Code: $\left\{1^{\prime} \mathrm{b} 1,16^{\prime} \mathrm{bx}\right\}$

- Tested on AlexNet $\rightarrow 2 \times$ overall BW Reduction

| Layer | Filter / Image bits ( $0 \%$ ) | Filter / Image BW Reduc. | IO / HuffiO (MB/frame) | Voltage <br> (V) | MMACs/ Frame | $\begin{aligned} & \text { Power } \\ & (\mathrm{mW}) \end{aligned}$ | $\begin{gathered} \text { Real } \\ \text { (TOPS/W) } \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| General CNN | 16 (0\%) / 16 (0\%) | 1.0x |  | 1.1 | - | 288 | 0.3 |
| AlexNet 11 | 7 (21\%) / 4 (29\%) | $1.17 \mathrm{x} / 1.3 \mathrm{x}$ | $1 / 0.77$ | 0.85 | 105 | 85 | 0.96 |
| AlexNet 12 | 7 (19\%) / 7 (89\%) | 1.15x / 5.8x | 3.2 / 1.1 | 0.9 | 224 | 55 | 1.4 |
| AlexNet 13 | 8 (11\%) / 9 (82\%) | $1.05 \mathrm{x} / 4.1 \mathrm{x}$ | 6.5 / 2.8 | 0.92 | 150 | 77 | 0.7 |
| AlexNet 14 | 9 (04\%) / 8 (72\%) | $1.00 \mathrm{x} / 2.9 \mathrm{x}$ | 5.4 / 3.2 | 0.92 | 112 | 95 | 0.56 |
| AlexNet 15 | 9 (04\%) / 8 (72\%) | $1.00 \mathrm{x} / 2.9 \mathrm{x}$ | 3.7 / 2.1 | 0.92 | 75 | 95 | 0.56 |
| Total / avg. | - | - | 19.8 / 10 | - | - | 76 | 0.94 |
| LeNet-5 11 | 3 (35\%) / 1 (87\%) | 1.40x / 5.2x | -0.008 10.007 | 0.7 | 0.3 | 25 | 1.07 |
| LeNet-5 12 | 4 (26\%) / 6 (55\%) | $1.25 \mathrm{x} / 1.9 \mathrm{x}$ | 0.050 / 0.042 | 0.8 | 1.6 | 35 | 1.75 |
| Total / avg. | - | - | 0.053 / 0.043 | - | - | 33 | 1.6 |

## I/O Compression in Eyeriss



## Off-Chip DRAM

64 bits

## Compression Reduces DRAM BW



Uncompressed Fmaps + Weights

RLE Compressed Fmaps + Weights

From information theory,

Simple RLC within $5 \%-10 \%$ of theoretical entropy limit [Chen, ISSCC 2016]

Entropy $\begin{aligned} H= & -\sum_{i=0}^{L-1} p_{i} \cdot \log _{2} p_{i}\end{aligned}$
minimum average number of bits required to code $f$

## Compressed Implicit Coordinate Representations

- "Empty" coordinate compression via zero-run encoding
- Run-length coding (RLE)
- (run-length of zeros, non-zero payload)...
- Significance map coding
- (flag to indicate if non-zero, non-zero payload)...
- Payload encoding
- Fixed length payload
- Variable length payload
- E..g., Huffman coding
- Efficiency of different traversal patterns through the tensor is affected by encoding, e.g., finding the payload for a particular coordinate...


## Compressed Explicit Coordinate Representations

- Coordinate list representation
- Struct of arrays form
(coordinate of non-zero value)...
(non-zero payload)...

- Array of structs form
(coordinate of non-zero value, non-zero payload)...
- Payload encoding
- Explicit


Black bar show scope of struct

- Immediate value
- Pointer
- Implicit
- Offset of coordinate is offset of payload


## More Explicit Coordinate Representations

- Coordinate Bitmask


Any complexity with lookupPayload()?


Have we seen a representation like this?


Is this useful even with no compression?


## Uncompressed/Compressed Representation



## Uncompressed/Compressed Representation

A specific implementation of the fibertree abstract type


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## Uncompressed/Compressed Representation



## Explicit Coordinate Representations

- Coordinate/Payload list
- (coordinate, non-zero payload)... (array of structs)
- (coordinate)..., (non-zero payload)... (struct of arrays)
- Hash table (per fiber)
- (coordinate -> payload) mapping
- Hash table (per rank)
- (fiber_id, coordinate -> payload) mapping
- Bit vector of non-zero coordinates
- Compressed or uncompressed payload


## Per Rank Tensor Representations

- Uncompressed [U]
- Run-length Encoded [R]
- Coordinate/Payload List [C]
- 
- Hash Table (per rank) [ $\mathrm{H}_{\mathrm{r}}$ ]
- Hash Table (per fiber) $\left[\mathrm{H}_{\mathrm{f}}\right]$
- Tagged union of any combination of previous types

Inspired by collaboration with Kjolstad

## Notation for CSR



## Representation of Order of Ranks

Differentiating CSR and CSC


## Traversal Efficiency

Efficiency of different traversal patterns through the tensor is affected by representation, e.g., finding the payload for a particular coordinate...

- Operations:
- maybe(payload) $=$ Fiber.getPayload(coordinate)
- (coordinate, payload) = Fiber.getNext(rank_traversal_order)

Fiber.getNext() is a useful iterator and its efficiency is highly dependent on representation, both order of ranks and representation of each rank....

## Concordant traversal orders

CSR and CSC each has a natural (or "concordant"*) traversal order

Row-major order
Processing
Order

Original Matrix



Compressed
Sparse Row (CSR)


Column-major order


Compressed
Sparse Column (CSC)


## Example Traversal Efficiency

- Efficiency of getPayload():
- Uncompressed - direct reference - $\mathrm{O}(1)$
- Run length encoded - linear search - O(n)
- Hash table - multiple references and compute - $O(1)$
- Coordinate/Payload list - binary search - O(log n)
- Efficiency of getNext() - (concordant traversal)
- Uncompressed - sequential reference, good spatial locality - O(1)
- Run length encoded - sequential reference - O(1)
- Coordinate/Payload list - same as uncompressed
- Efficiency of getNext() (discordant traversal)
- Essentially as good (or bad) as getPayload-method....


## Traversing a Sparse Tensor



## Traversing a Sparse Tensor



## Traversing a Sparse Tensor



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## Tensor Traversal (2-D)



## Tensor Traversal (2-D)

$$
Z=T_{h, w}
$$



| t_pos | $\mathbf{h}$ | $\mathbf{t}$ _h_pos | $\mathbf{w}$ | $\mathbf{t}$ _val |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | $?$ | $?$ | $?$ |
| 0 | 0 | 0 | 0 | a |
| 0 | 0 | 1 | 2 | c |
| 1 | 2 | $?$ | $?$ | $?$ |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |

## Tensor Traversal (2-D)



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## Tensor Traversal (2-D)



## Tensor Traversal (2-D)



## Abstraction versus Implementation

- Abstraction
- An interface and semantics
- Attributes: No implementation, data layout or timing
- Use: implementation-agnostic understanding
- Examples:
- Fibers
- Fibertree
- Implementation
- Specific implementation of an abstract spec
- Attributes: Concrete implementation, data layout and timing
- Examples:
- Fibers $\rightarrow$ uncompressed array, coordinate/payload list
- Fiber-tree $\rightarrow$ CSR, CSC, CSF, COO...


## Tensor Traversal (CSR Style)

```
# 2-D Tensor Traversal (CSR)
t_segs = Array(H)
t_coords = Array(W)
t_vals = Array(W)
    sum = 0
for t_h_pos in [0,H):
    h = t_h_pos
    t_w_start = t_segs[t_h_pos]
    t_w_len = t_segs[t_h_pos+1]-t_w_start
    for t_w_pos in [t_w_start, t_w_len):
        h = t_coords[t_w_pos]
        t_val = t_vals[t_w_pos]
            sum += t_val
```

For uncompressed rank coordinate equals position

## Merging Ranks

For efficiency one can form new representations where the data structure for two or more ranks are combined.


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## Merging Ranks

- For efficiency one can form new representations where the data structure for two or more ranks are combined:
- Examples:
- Tensor-(C²)

List of (coordinate tuple, payload) - COO

- Tensor-(H ${ }^{2}$ )
- Hash table with coordinate tuple as key
- Tensor-(U2)
- Flattened array
- Coordinates can be recovered with modulo arithmetic on "position"
- Tensor-( $\mathrm{R}^{2}$ )
- Flattened run-length encoded sequence


## Splitting Fibers - Coordinate Space



## Splitting Fibers - Coordinate Space



Split uniformly by coordinates (groups of 8 coordinates)


## Splitting Fibers - Coordinate Space



Split uniformly by coordinates (groups of 8 coordinates)


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Split uniformly by coordinates (groups of 8 coordinates)


## Splitting Fibers - Coordinate Space



Split uniformly by coordinates (groups of 8 coordinates)


## Splitting Fibers - Coordinate Space



Split uniformly by coordinates (groups of 8 coordinates)


## Splitting Fibers - Position Space



## Splitting Fibers - Position Space



Split evenly by occupancy (groups of 4)

## Splitting Fibers - Position Space



Split evenly by occupancy (groups of 4)


## Splitting Fibers - Position Space



Split evenly by occupancy (groups of 4)


## Splitting Fibers - Position Space



Split evenly by occupancy (groups of 4)


## Splitting Fibers - Position Space



Split evenly by occupancy (groups of 4)


## Splitting Fibers - Position Space



Split evenly by occupancy (groups of 4)


## Splitting Fibers - Position Space



Split evenly by occupancy (groups of 4)


## Splitting Fibers - Position Space



Split evenly by occupancy (groups of 4)

## Fibertree Representation of Weight Pruning



Each dimension in the original tensor is represented as a rank in the tree

## Specification of Channel-based Sparsity



## Flattening: Specification of Sub-kernel Sparsity



Fibertree Representation of Dense Tensor


## Flattening: Specification of Sub-kernel Sparsity



Fibertree Representation of Dense Tensor


Flattening: Specification of Unstructured Sparsity


Fibertree Representation of Dense Tensor

## Flattening: Specification of Unstructured Sparsity



Fibertree Representation of Dense Tensor


## Flattening: Specification of Unstructured Sparsity



Fibertree Representation of Dense Tensor


## Reordering \& Partitioning: Specification of 2:4 Sparsity



Fibertree Representation of Dense Tensor

Reordering \& Partitioning: Specification of 2:4 Sparsity


Fibertree Representation of Dense Tensor


## Reordering \& Partitioning: Specification of 2:4 Sparsity



## Reordering \& Partitioning: Specification of 2:4 Sparsity



Fibertree Representation of Dense Tensor


## Hierarchical Structured Sparsity (HSS)



Fibertree Representation of Dense Tensor
(1) Reorder Ranks
(2) Partition Rank C into N ranks ( $\mathrm{N}>=2$ ) , e.g., $\mathrm{N}=3$ as shown below
(3) Apply Per-rank Pruning Rule


## Hierarchical Structured Sparsity (HSS)


(1) Reorder Ranks
(2) Partition Rank C into N ranks ( $\mathrm{N}>=2$ ) , e.g., $\mathrm{N}=3$ as shown below
(3) Apply Per-rank Pruning Rule

- N-1 rank HSS defined as

$$
\begin{aligned}
\cdot R S \rightarrow C_{N-1} \rightarrow C_{N-2}\left(G_{N-2}: H_{N-2}\right) \\
\\
\rightarrow \ldots \rightarrow C_{1}(3: 4) \rightarrow C_{0}(2: 4)
\end{aligned}
$$

- HSS qualitative difference: allows pruning rules for more thank one ranks
- HSS provides a systematic and modularized way to represent a large number of sparsity degrees



## Next: Sparse Architectures

