

6.S078 - Computer Architecture:  
A Constructive Approach

# Combinational ALU

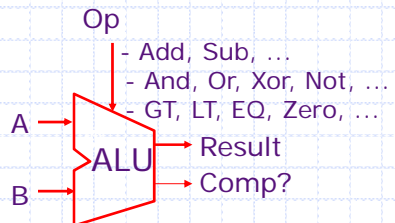
Arvind  
Computer Science & Artificial Intelligence Lab.  
Massachusetts Institute of Technology

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-1

# Arithmetic-Logic Unit (ALU)



ALU performs all the arithmetic  
and logical functions

We will first implement individual functions  
and then combine them to form an ALU

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-2

## A w-bit Ripple-Carry Adder

```
function Bit#(TAdd#(w,1)) addN(Bit#(w) x, Bit#(w) y,  
                               Bit#(1) c0);  
  Bit#(w) s; Bit#(TAdd#(w,1)) c; c[0] = c0;  
  let valw = valueOf(w);  
  for(Integer i=0; i<valw; i=i+1)  
  begin  
    let cs = fa(x[i],y[i],c[i]);  
    c[i+1] = cs[1]; s[i] = cs[0];  
  end  
  return {c[valw],s};  
endfunction
```

Structural interpretation of a loop – unfold it to generate an acyclic graph

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-3

## Integer versus Int#(32)

- ◆ In mathematics integers are unbounded but in computer systems integers always have a fixed size
- ◆ Bluespec allows us to express both types of integers, though unbounded integers are used only as a programming convenience

```
for(Integer i=0; i<valw; i=i+1)  
begin  
  let cs = fa(x[i],y[i],c[i]);  
  c[i+1] = cs[1]; s[i] = cs[0];  
end
```

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-4

## Static Elaboration phase

- ◆ When Bluespec program are compiled, first type checking is done and then the compiler gets rid of many constructs which have no direct hardware meaning, like Integers, loops

```
for(Integer i=0; i<valw; i=i+1) begin
  let cs = fa(x[i],y[i],c[i]);
  c[i+1] = cs[1]; s[i] = cs[0];
end
```

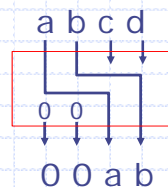
```
cs0 = fa(x[0], y[0], c[0]); c[1]=cs0[1]; s[0]=cs0[0];
cs1 = fa(x[1], y[1], c[1]); c[2]=cs1[1]; s[1]=cs1[0];
...
csw = fa(x[valw-1], y[valw-1], c[valw-1]);
c[valw] = csw[1]; s[valw-1] = csw[0];
```

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-5

## Logical right shift by 2



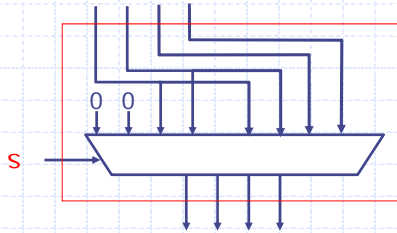
- ◆ Fixed size shift operation is cheap in hardware – just wire the circuit appropriately
- ◆ Rotate, sign-extended shifts – all are equally easy

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-6

## Conditional operation: shift versus no-shift



- ◆ We need a Mux to select the appropriate wires: if **s** is one the Mux will select the wires on the left otherwise it would select wires on the right

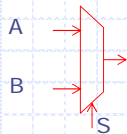
```
(s==0)?{a,b,c,d}:{0,0,a,b};
```

February 10, 2012

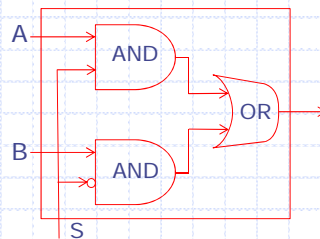
<http://csg.csail.mit.edu/6.S078>

L2-7

## A 2-way multiplexer



```
(s==0)?A:B
```



Gate-level implementation

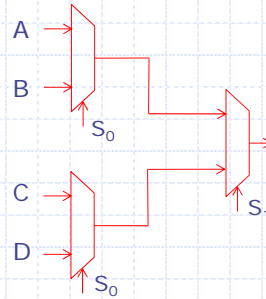
February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-8

## A 4-way multiplexer

```
case {s1,s0} matches
  0: A;
  1: B;
  2: C;
  3: D;
endcase
```



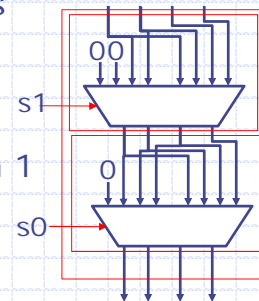
February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-9

## Logical right shift by $n$

- ◆ Shift  $n$  can be broken down in several steps of fixed-length shifts of 1, 2, 4, ...
- ◆ Thus a shift of size 3 can be performed by first doing a shift of length 2 and then a shift of length 1
- ◆ We need a Mux to select whether we need to shift at all
  - The whole structure can be expressed in a bunch of nested conditional expressions



You will write a Bluespec program to produce a variable size shifter in Lab 1

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-10

## Type synonyms

```
typedef Bit [7:0] Byte;  
  
typedef Bit#(8) Byte;  
  
typedef Bit [31:0] Word;  
  
typedef Tuple2#(a,a) Pair#(type a);  
  
typedef Int#(n) MyInt#(type n);  
  
typedef Int#(n) MyInt#(numeric type n);
```

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-11

## Enumerated types

```
typedef enum {Red, Blue, Green}  
Color deriving(Bits, Eq);  
  
typedef enum {Eq, Neq, Le, Lt, Ge, Gt}  
BrType deriving(Bits, Eq);  
  
typedef enum {Add, Sub, And, Or, Xor, Nor, Slt, Sltu,  
              LShift, RShift, Sra}  
AluFunc deriving(Bits, Eq);
```

January 18, 2012

<http://csg.csail.mit.edu/SNU>

L7-12

## Structure type

Example: Complex Addition

```
typedef struct{
    Int#(t) r;
    Int#(t) i;
} Complex#(numeric type t) deriving (Eq,Bits);

function Complex#(t) \+
    (Complex#(t) x, Complex#(t) y);
    Int#(t) real = x.r + y.r;
    Int#(t) imag = x.i + y.i;
    return(Complex{r:real, i:imag});
endfunction
```

January 12, 2012

<http://csg.csail.mit.edu/SNU>

L4-13

## Combinational ALU

```
function Bit#(width) alu(Bit#(width) a,
                        Bit#(width) b, AluFunc op);
    Bit#(width) res = case(op)
        Add    : add(a, b);
        Sub    : subtract(a, b);
        And    : (a & b);
        Or     : (a | b);
        Xor    : (a ^ b);
        Nor    : ~(a | b);
        Slt    : setLessThan(a, b);
        Sltu   : setLessThanUnsigned(a, b);
        LShift : logicalShiftLeft(a, b[4:0]);
        RShift : logicalShiftRight(a, b[4:0]);
        Sra    : signedShiftRight(a, b[4:0]);
    endcase;
    return res;
endfunction
```

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-14

## Comparison operators

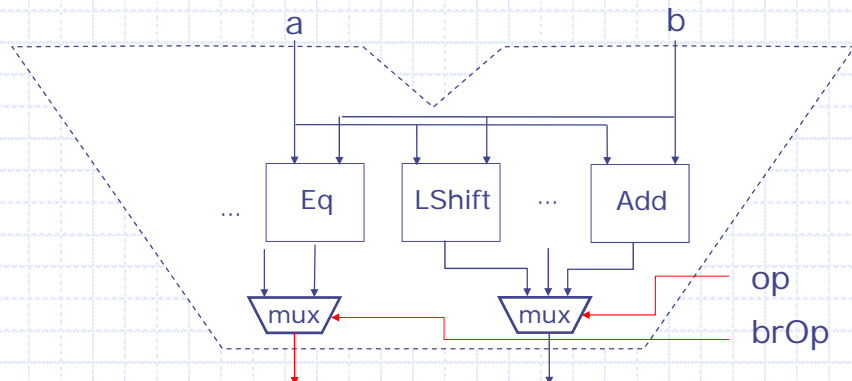
```
function Bool aluBr(Bit#(width) a,  
                  Bit#(width) b, BrType brOp);  
  Bool brTaken = case(brOp)  
    Eq   : (a == b);  
    Neq  : (a != b);  
    Le   : signedLE(a, 0);  
    Lt   : signedLT(a, 0);  
    Ge   : signedGE(a, 0);  
    Gt   : signedGT(a, 0);  
  endcase;  
  return brTaken;  
endfunction
```

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-15

## ALU including Comparison operators



February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-16



# Complex combinational circuits

February 10, 2012

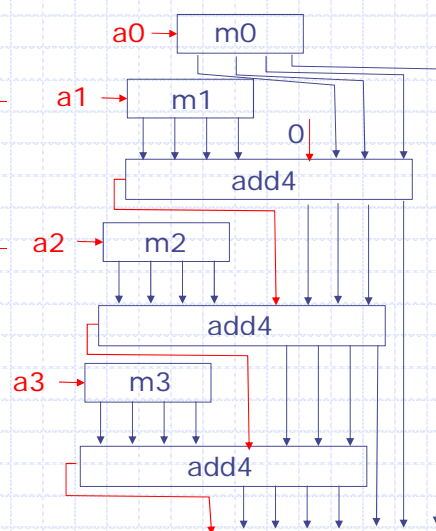
<http://csg.csail.mit.edu/6.S078>

L2-17

# Multiplication by repeated addition

b	Multiplicand	1101	(13)
a	Multiplier	* 1011	(11)
		1101	
		+ 1101	
		+ 0000	
		+ 1101	
		10001111	(143)

$$m_i = (a[i]==0)? 0 : b;$$



February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-18

## Combinational 32-bit multiply

```
function Bit#(64) mul32(Bit#(32) a, Bit#(32) b);
    Bit#(32) prod = 0;
    Bit#(32) tp = 0;
    for(Integer i = 0; i < 32; i = i+1)
    begin
        Bit#(32) m = (a[i]==0)? 0 : b;
        Bit#(33) sum = add32(m,tp,0);
        prod[i] = sum[0];
        tp = truncateLSB(sum);
    end
    return {tp,prod};
endfunction
```

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-19

## Combinational n-bit multiply

```
function Bit#(TAdd#(w,w)) mulN(Bit#(w) a, Bit#(w) b);
    Bit#(w) prod = 0;
    Bit#(w) tp = 0;
    for(Integer i = 0; i < valueOf(w); i = i+1)
    begin
        Bit#(w) m = (a[i]==0)? 0 : b;
        Bit#(TAdd#(w,1)) sum = addN(m,tp,0);
        prod[i] = sum[0];
        tp = truncateLSB(sum);
    end
    return {tp,prod};
endfunction
```

February 10, 2012

<http://csg.csail.mit.edu/6.S078>

L2-20

# Design issues with combinational multiply

- ◆ Lot of hardware
  - 32-bit multiply uses 31 addN circuits
- ◆ Long chains of gates
  - 32-bit ripple carry adder has a 31 long chain of gates
  - 32-bit multiply has 31 ripple carry adders in sequence!

The speed of a combinational circuit is determined by its longest input-to-output path

*next time: Sequential circuits*