

Hardware Security Module

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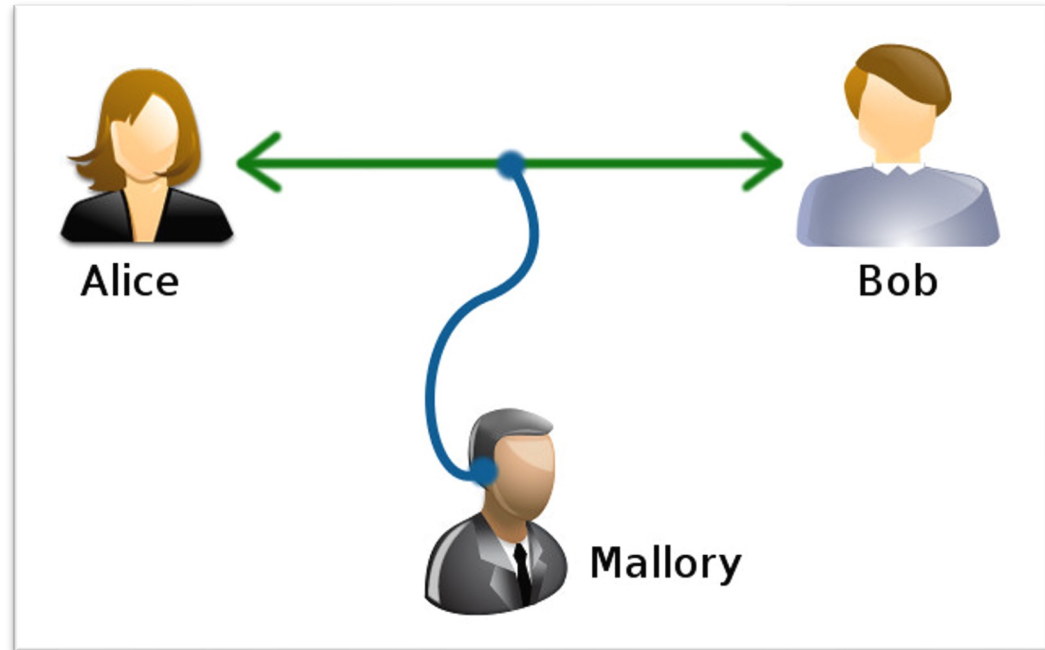


Outline

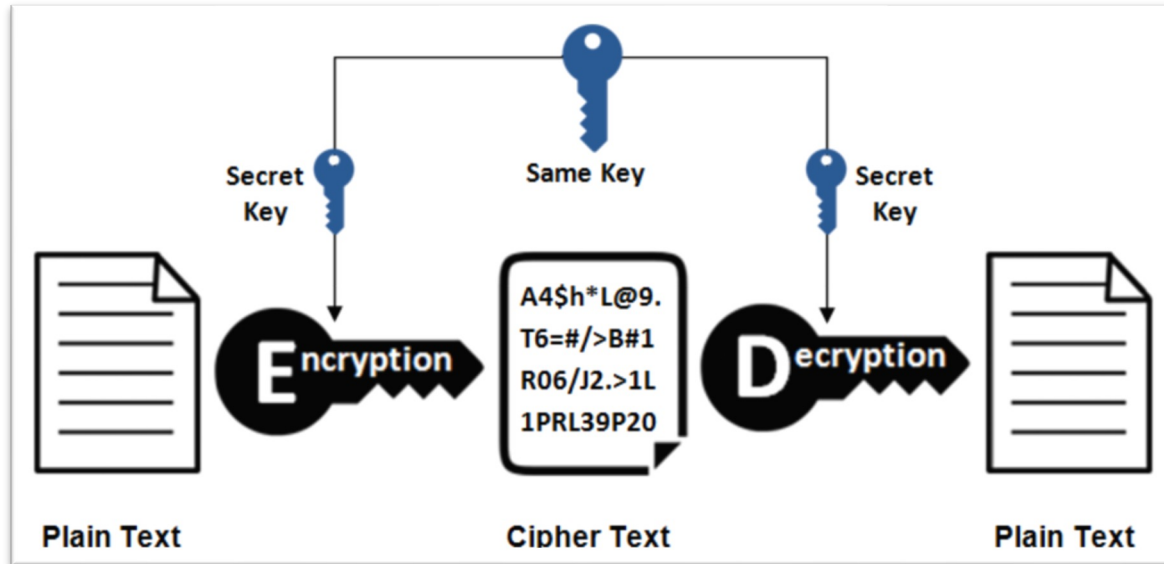
- Crypto is great, but real-world security also needs hardware support
- Design considerations and tradeoffs when designing hardware security modules

Security Property and Crypto Primitives

- Confidentiality
 - Symmetric
 - Asymmetric
- Integrity
- Freshness



Symmetric Cryptography



- One-time-pad (OTP)

Encryption:

$$\text{ciphertext} = \text{key} \oplus \text{plaintext}$$

Decryption:

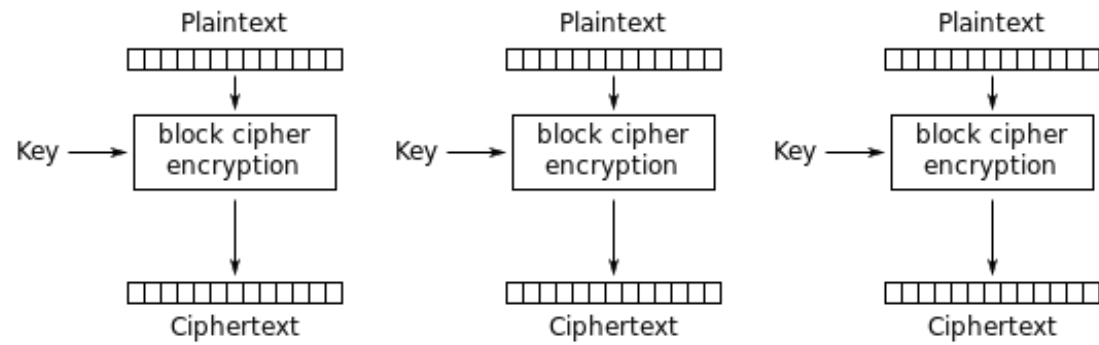
$$\text{plaintext} = \text{key} \oplus \text{ciphertext}$$

How about encrypting arbitrary length message? Any problems?

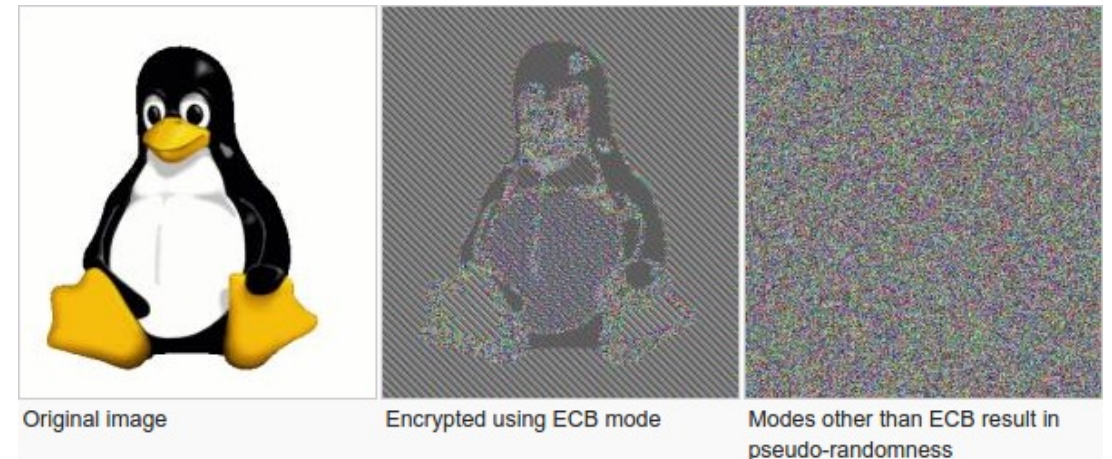
Block ciphers (e.g., DES, AES)

- Divide data in blocks and encrypt/decrypt each block
- AES block size can be 128, 192, 256 bits

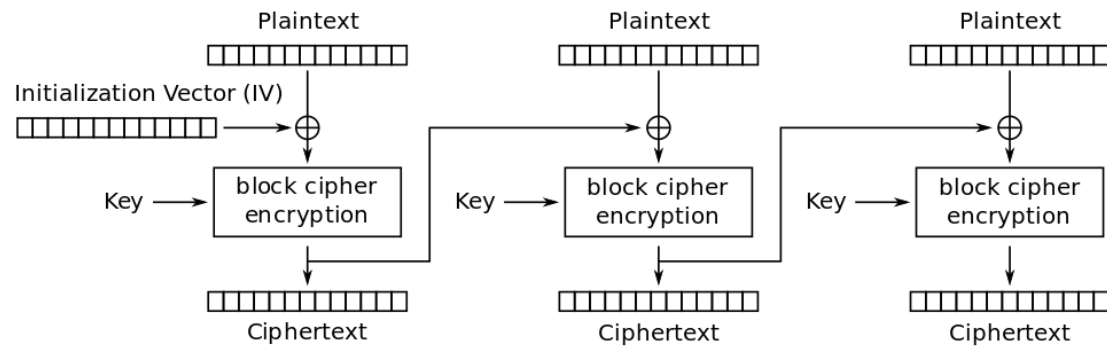
**ECB IS NOT
RECOMMENDED**



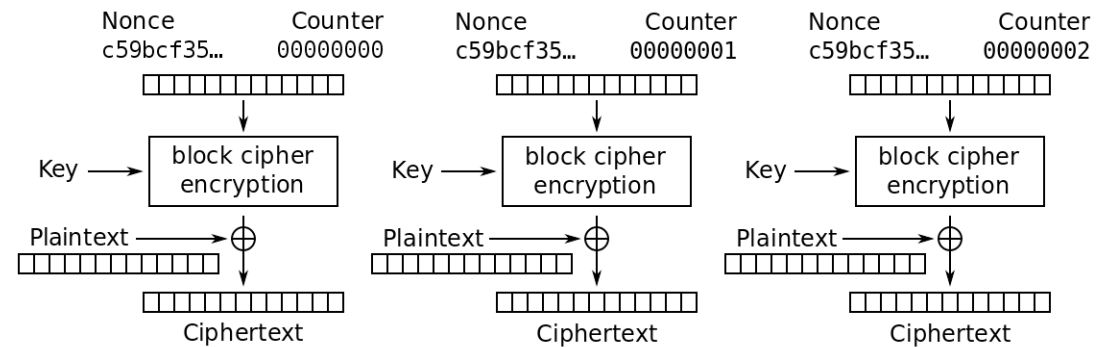
Electronic Codebook (ECB) mode encryption



Block ciphers (e.g., DES, AES)



Cipher Block Chaining (CBC) mode encryption



Counter (CTR) mode encryption

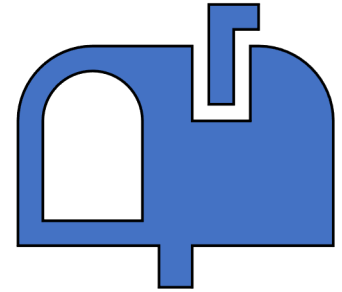
IV can be public, but need to ensure to not reuse IV for the same key.

Real-world application: file/disk encryption and memory encryption.

How to exchange the shared key between two parties?

Asymmetric Cryptography (e.g., RSA)

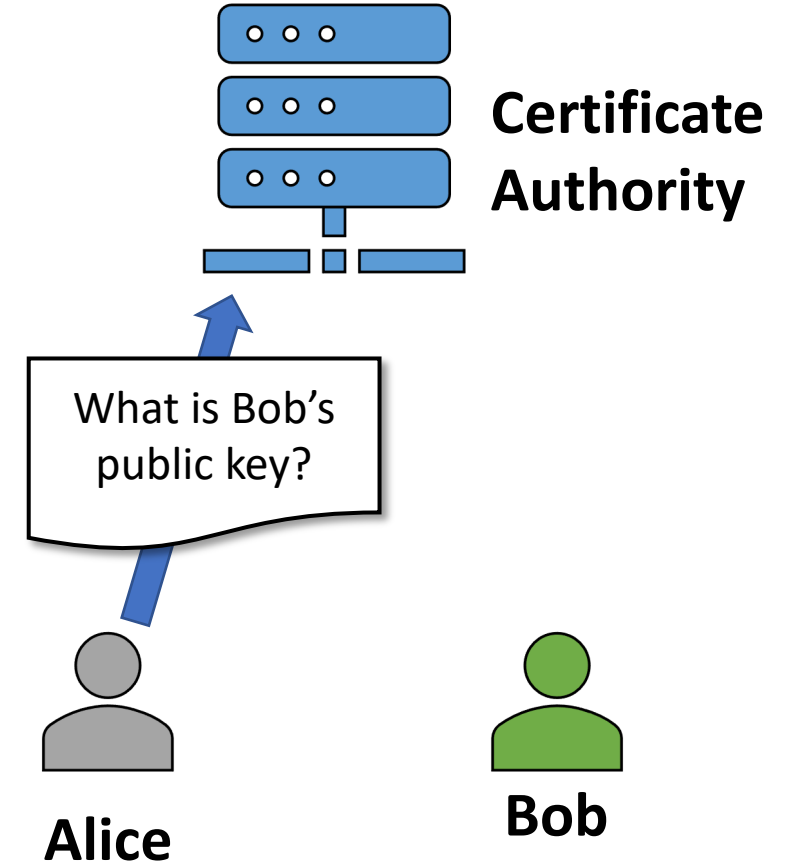
- A pair of keys:
 - Private key (K_{private} – kept as secret)
 - Public key (K_{public} – safe to release publicly)
- Computation:
 - $\text{Encrypt}(\text{plaintext}, K_{\text{public}}) = \text{ciphertext}$
 - $\text{Decrypt}(\text{ciphertext}, K_{\text{private}}) = \text{plaintext}$
- Computationally more expensive, so usually use asymmetric cryptography to negotiate a shared key (e.g., DKE key exchange), then use symmetric cryptography
- How to announce and obtain the public key?



Mail box is public;
Box key is private

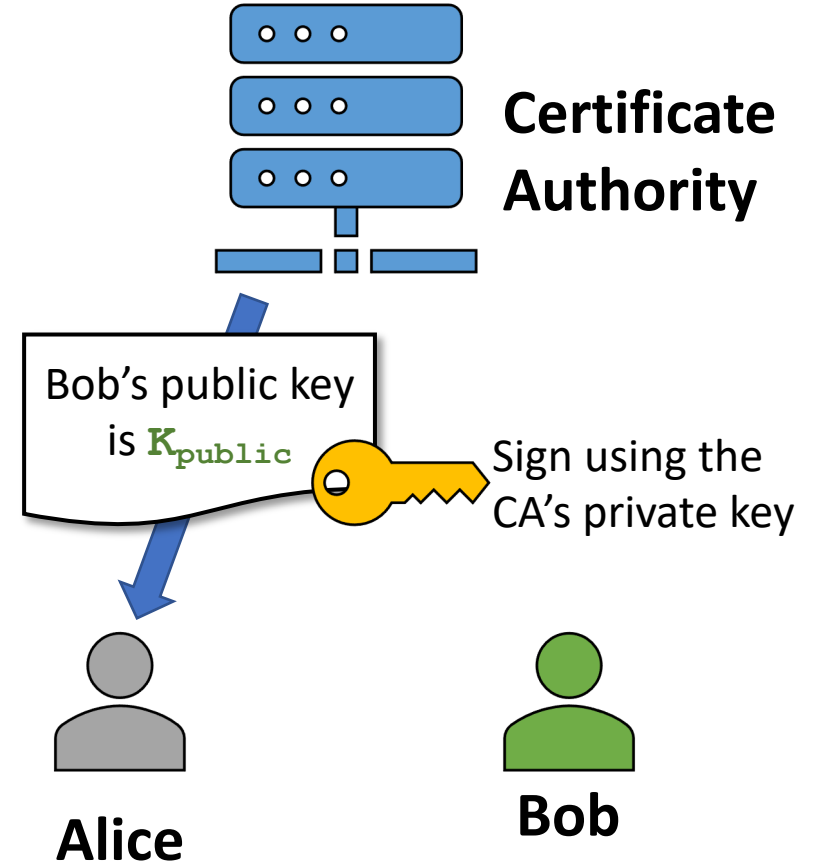
Public Key Infrastructures (PKIs)

- Bob has a private key K_{private} and wants to claim he corresponds to a public key K_{public}
- Analogy: public key is like a government-issued ID, need to be validated by an authority.

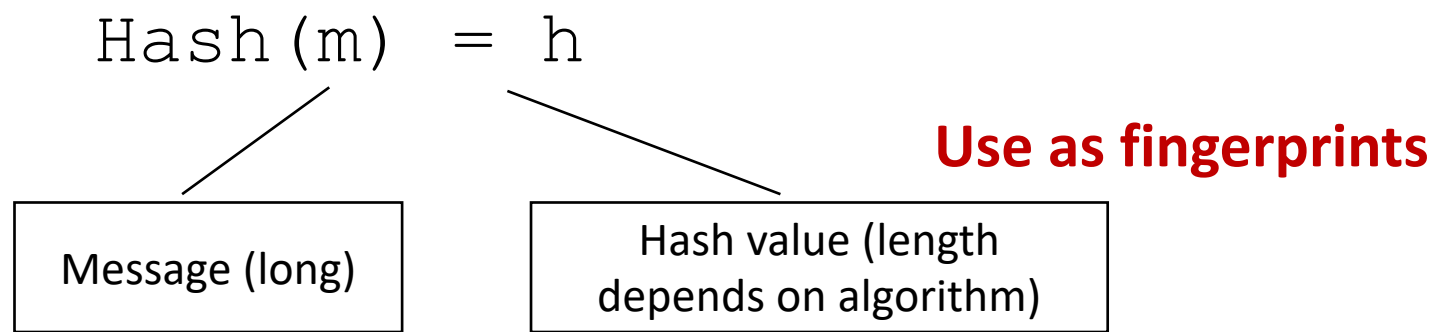


Public Key Infrastructures (PKIs)

- Bob has a private key K_{private} and wants to claim he corresponds to a public key K_{public}
- Analogy: public key is like a government-issued ID, need to be validated by an authority.
- Establish a chain of trust
- **Real-world use cases:** identify website, identify hardware chips/processors



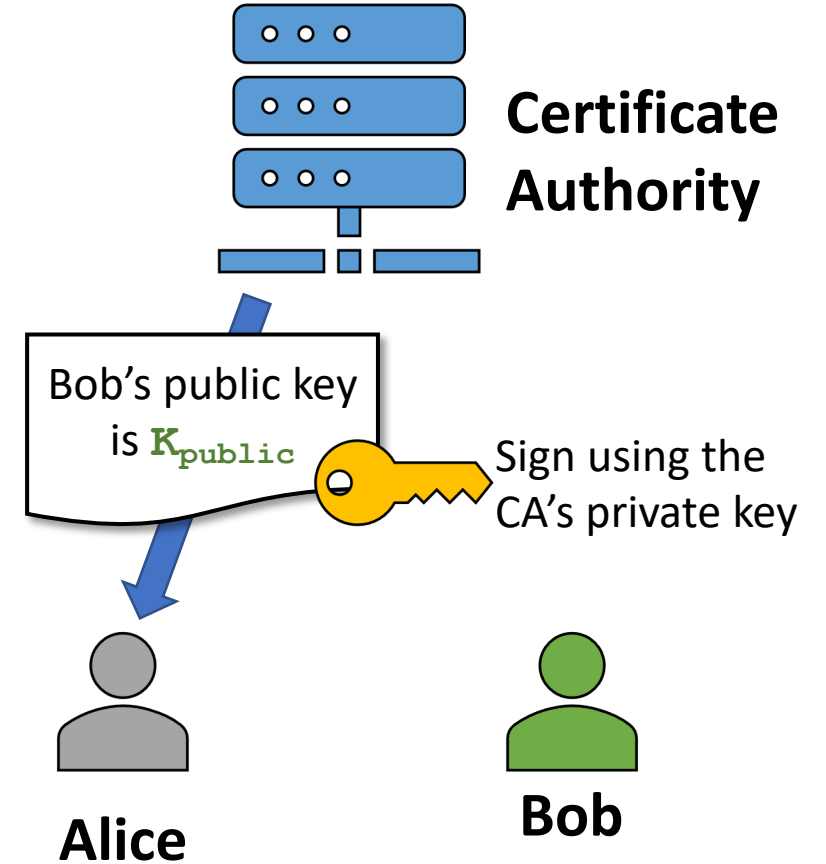
Integrity (MAC/Signature)



- One-way hash
 - Practically infeasible to invert, and difficult to find collision
- Avalanche effect
 - “Bob Smith got an A+ in ELE386 in Spring 2005” → 01eace851b72386c46
 - “Bob Smith got an B+ in ELE386 in Spring 2005” → 936f8991c111f2cefaw
- When message is long
 - Divide message into blocks, and keep extending the hash by adding previous hash

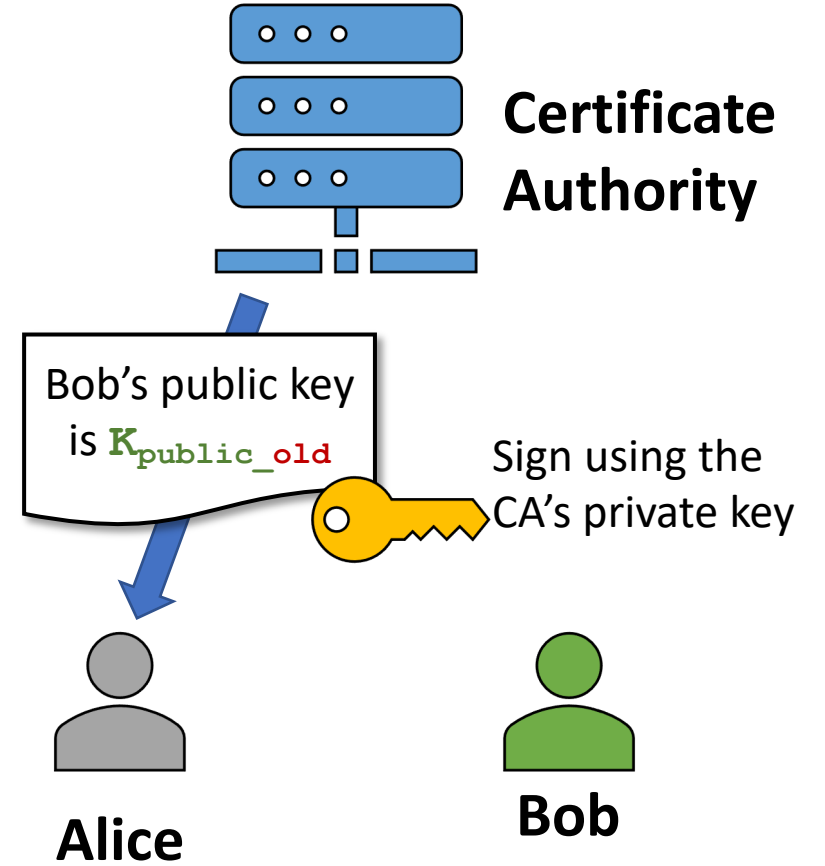
Integrity + Crypto

- Using symmetric crypto:
 - $\text{hash} = \text{SHA}(\text{message})$
 - $\text{HMAC} = \text{enc}(\text{hash}, \text{key})$
- Using asymmetric crypto:
 - Sign: $\text{sig} = \text{enc}(\text{hash}, K_{\text{private}})$
 - Verify:
 - $\text{hash}' = \text{SHA}(\text{message})$
 - $\text{sig} = \text{dec}(\text{sig}, K_{\text{public}})$



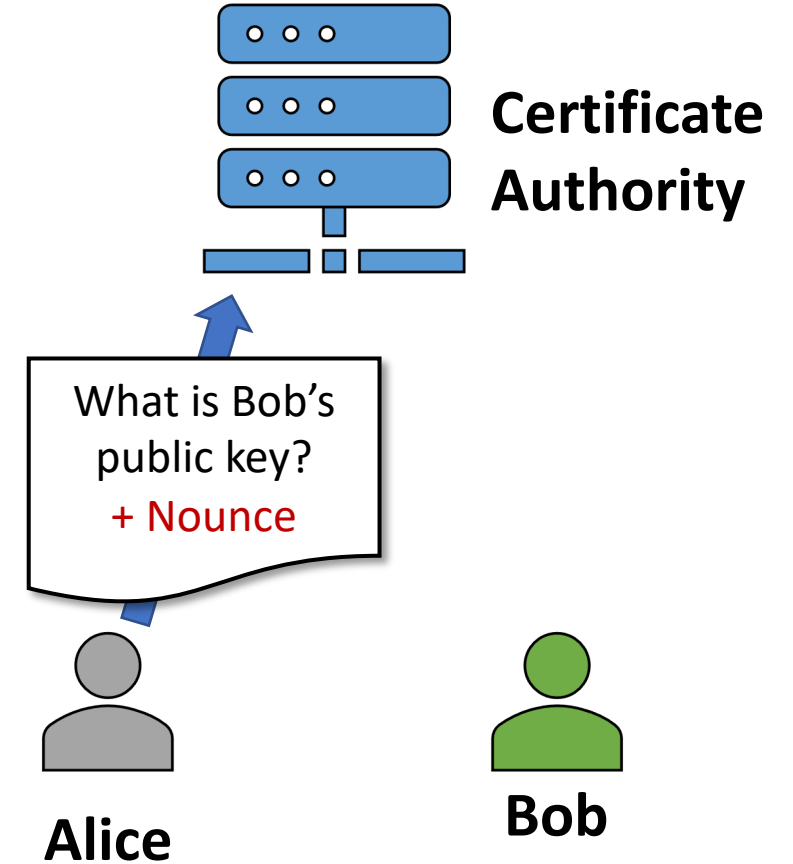
Freshness

- Goal: to block replay attack



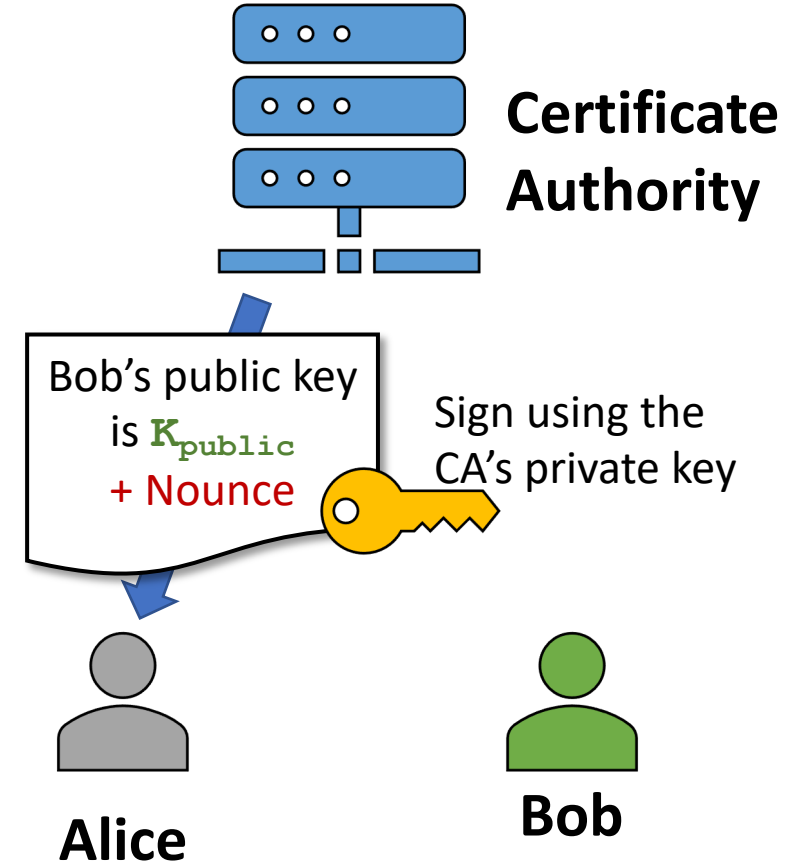
Freshness

- Goal: to block replay attack
- Nounces + Integrity
 - Nonce is a one-time use random number
 - Should not reuse the same nonce

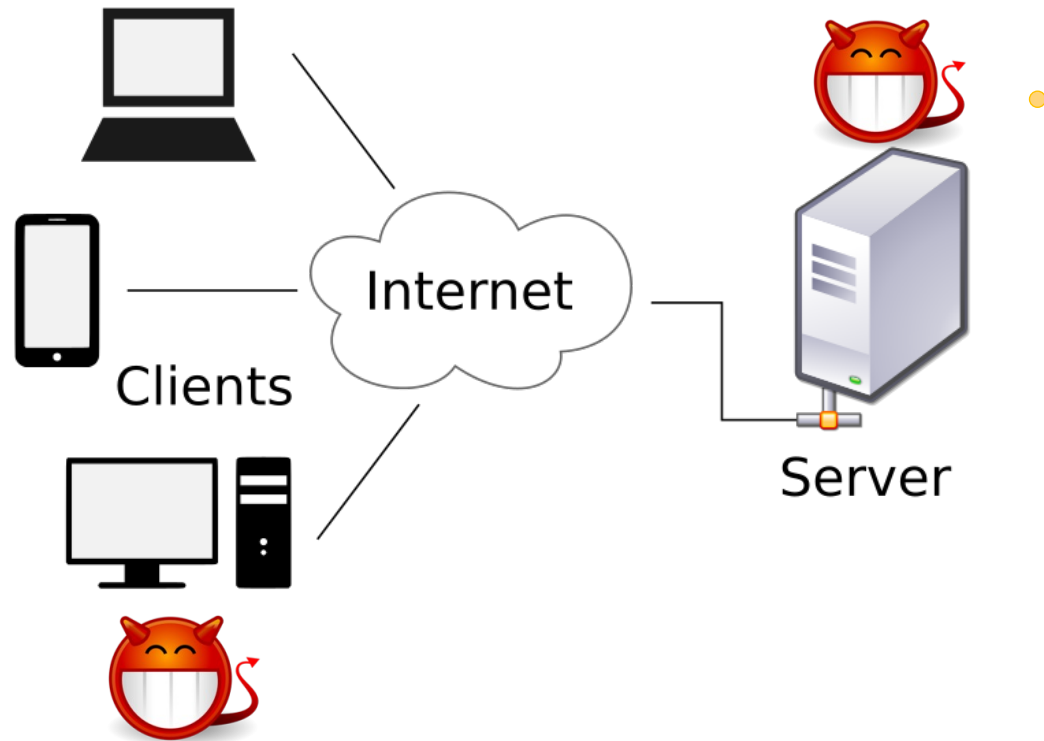


Freshness

- Goal: to block replay attack
- Nonces + Integrity
 - Nonce is a one-time use random number
 - Should not reuse the same nonce
- Challenge-response



Security Contexts #1



Hardware establishes
root of trust.

- a) An end-user wants to trust a remote server, e.g., bank server.
- b) A remote server wants to trust an end-user, e.g., when joining a company's highly-secure network.
- c) Lost device, rootkits?

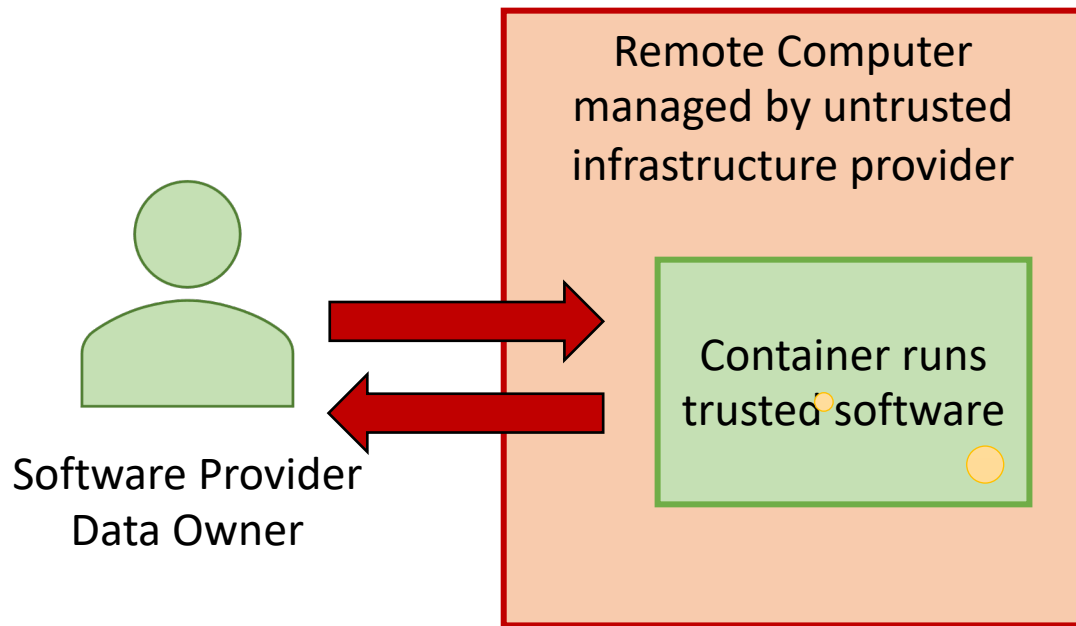
Security Contexts #2



- Software piracy (copying and reselling software to gain benefits).
- Disk lost or removed, leading to confidentiality leakage.
- Data encryption with weak passwords, such as, 6-digit passcode.

Bind data/application
with hardware.

Security Contexts #3



- Remote computation where the host hypervisor and OS is not trusted.

Hardware offers stronger isolation.

What Can Hardware Security Modules Offer?

- Establish root of trust
- Bind data and applications with the hardware device
- Offer stronger isolation
- More efficient

Secure Processors

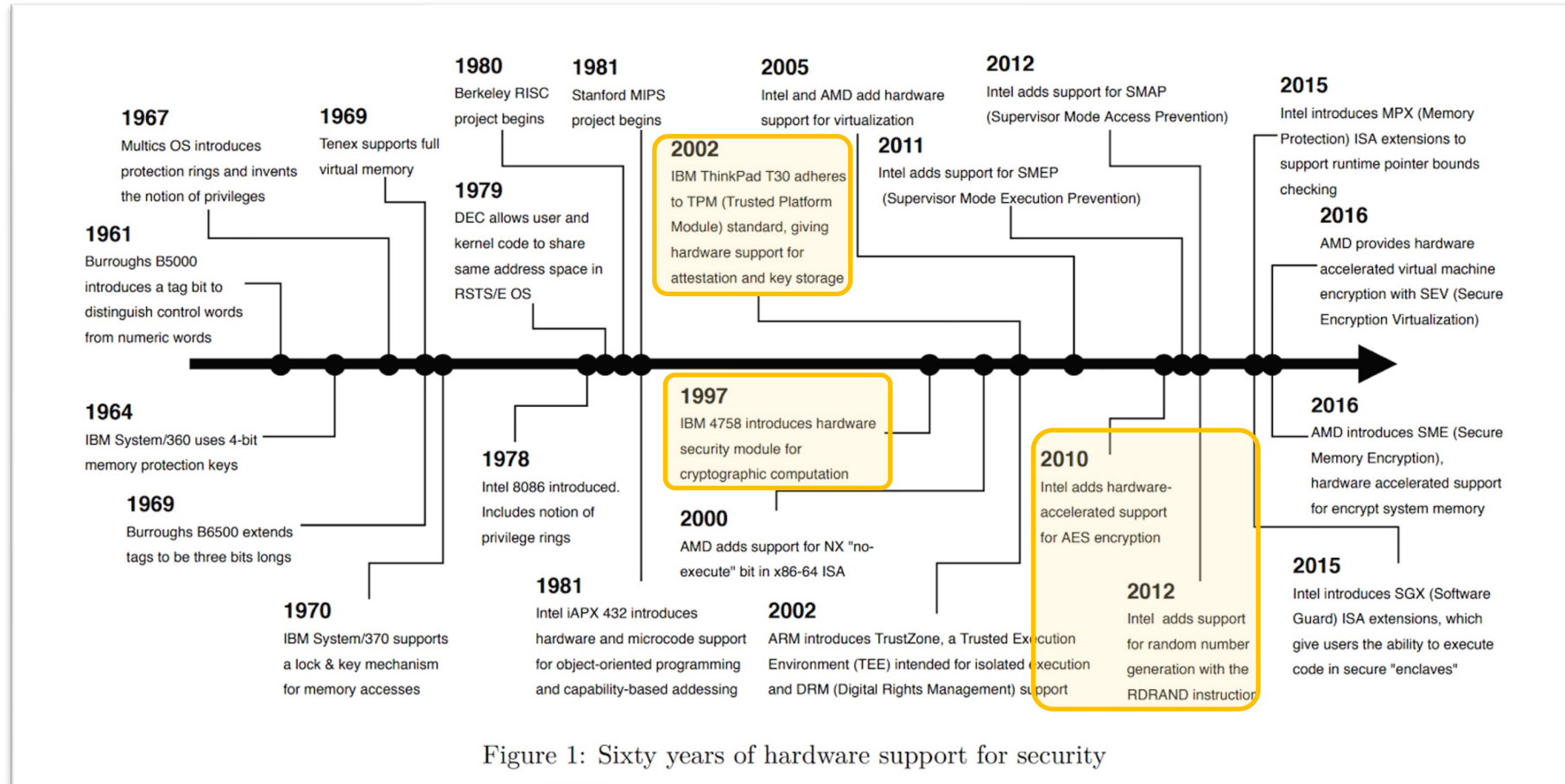
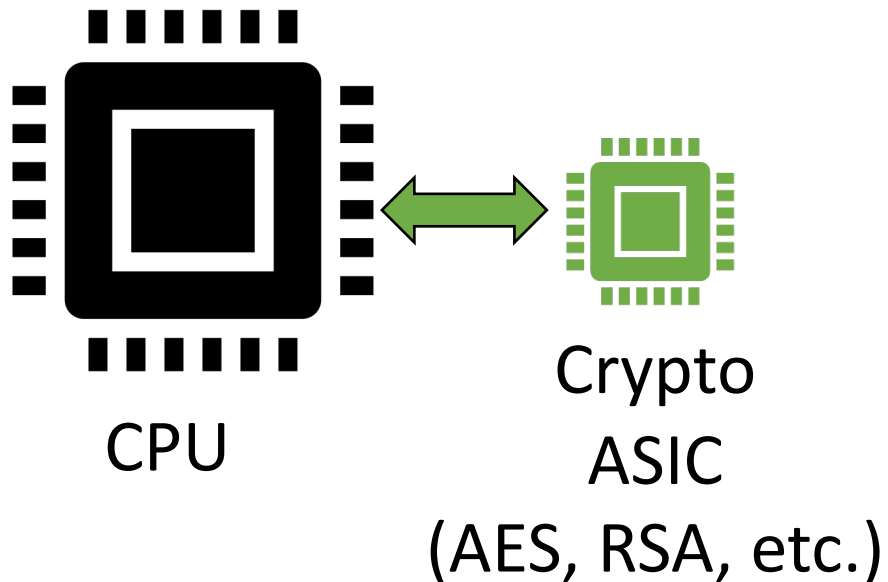


Figure 1: Sixty years of hardware support for security

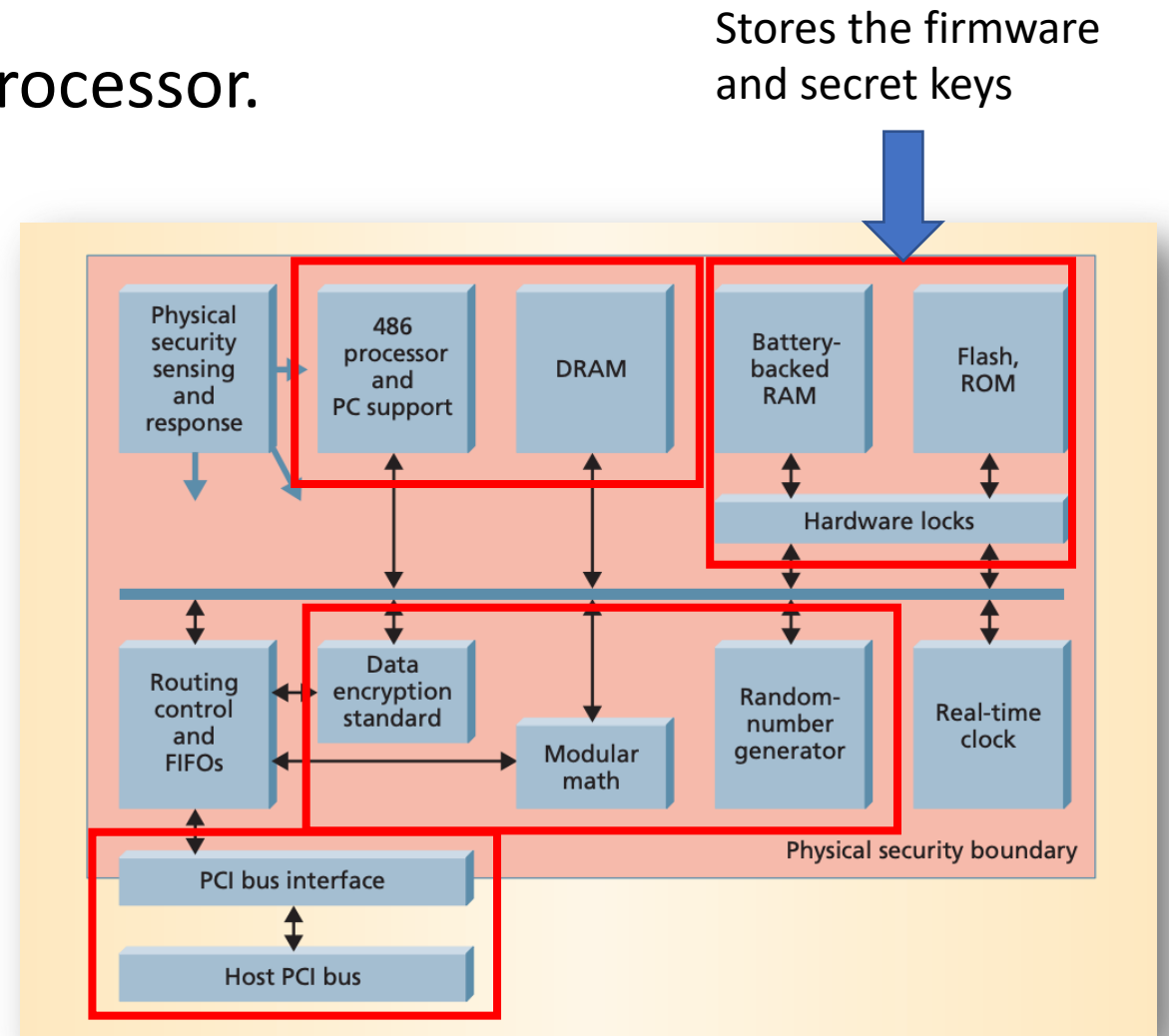
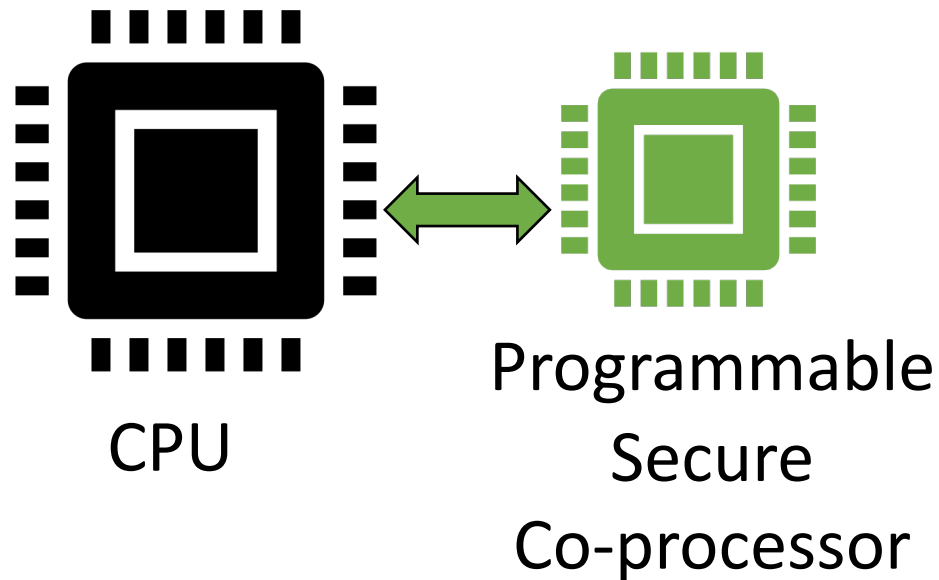
Before IBM 4758 (1999)

- Crypto Accelerators
 - Better performance
 - Simple functionality
 - Narrow interface



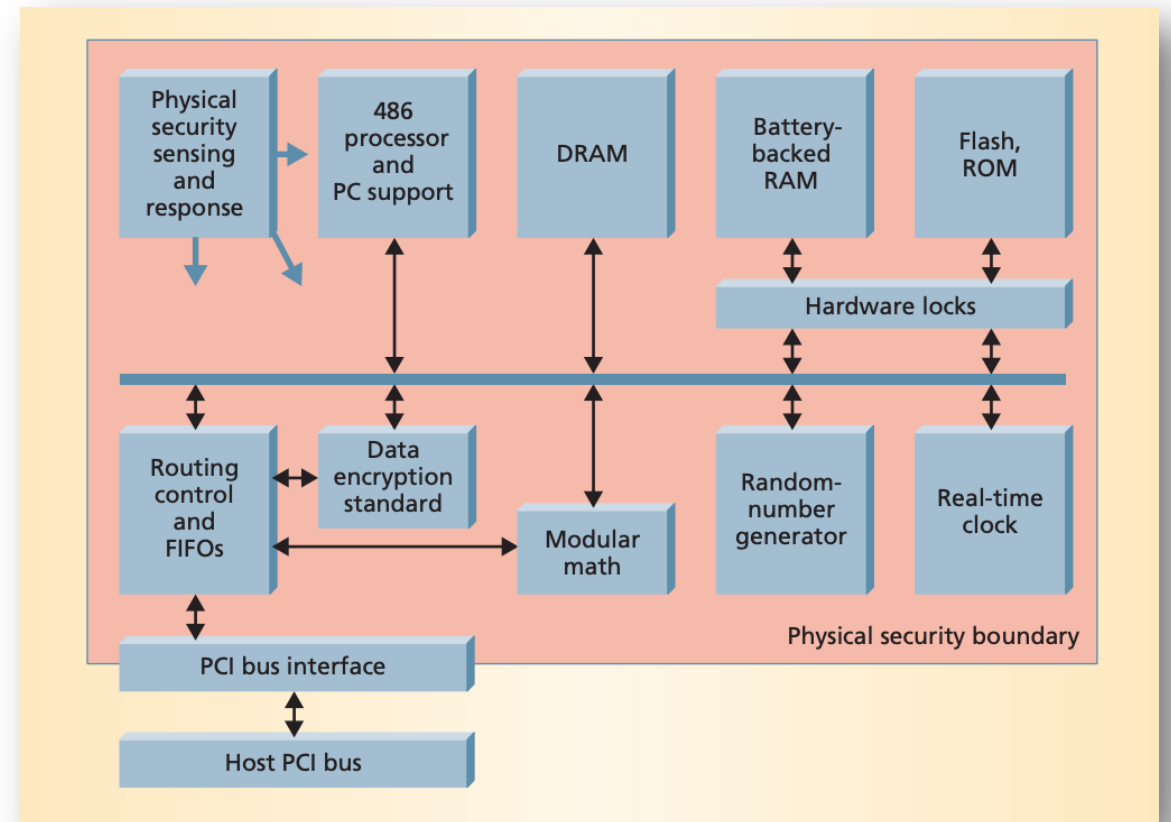
IBM 4758 (1999) -- 4765 (2012)

- Goal: a programmable, secure co-processor.
- High level idea: virtual locker room

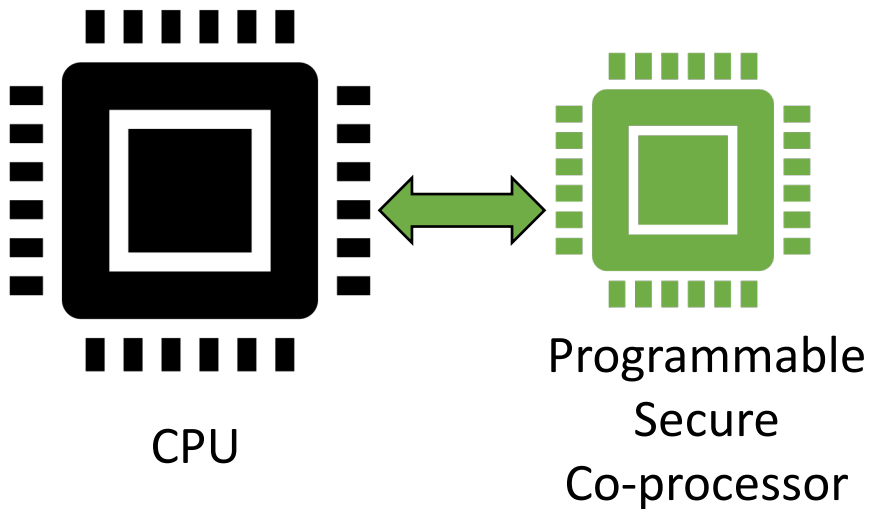


Software Layer Design and Concerns

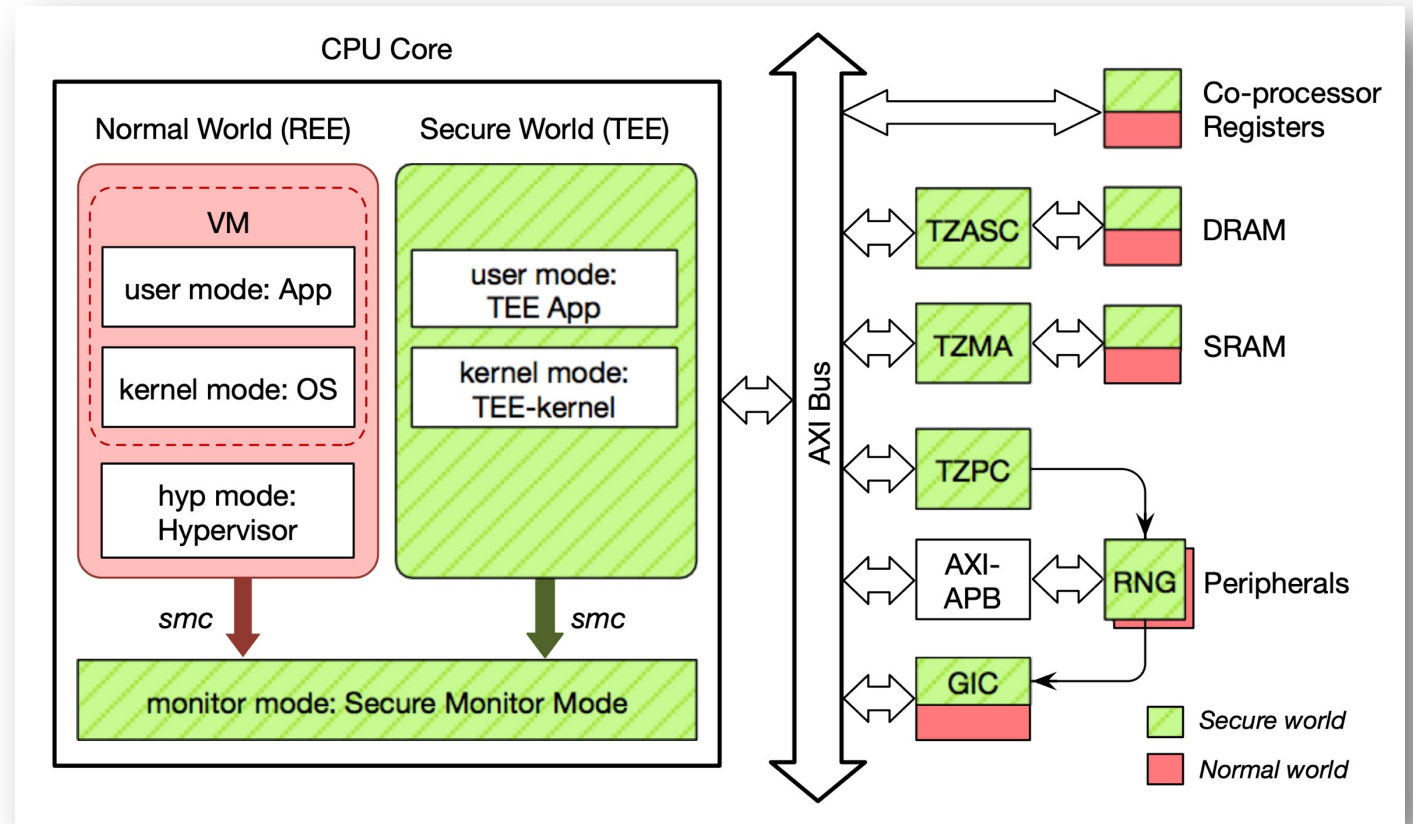
- Software stack:
 - Application
 - OS, kernel (microkernel)
 - Loaders, firmware, etc.
- Use cases:
 - Solve music/software piracy issue
 - Run an SSL server inside to store the agreed symmetric session keys
 - Bank application
- Problems:
 - Update software is tricky
 - Bad programmability due to microkernel



Compare to TrustZone



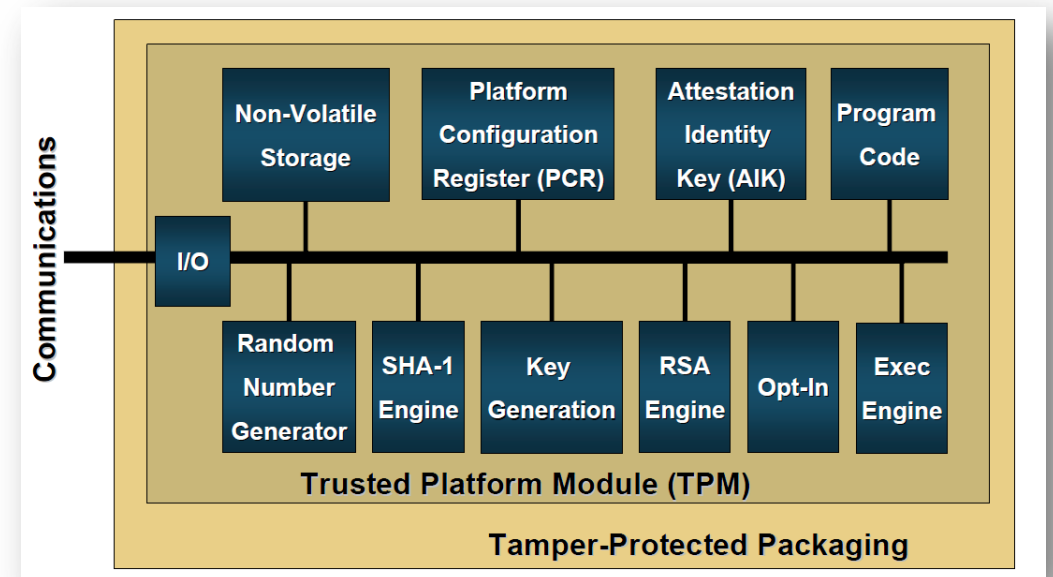
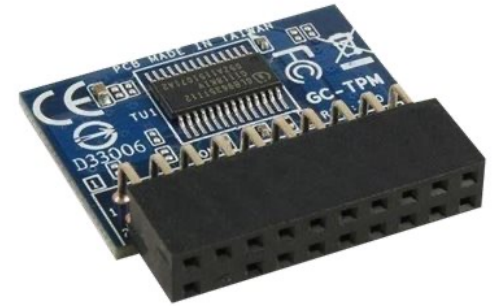
IBM 4758



ARM TrustZone

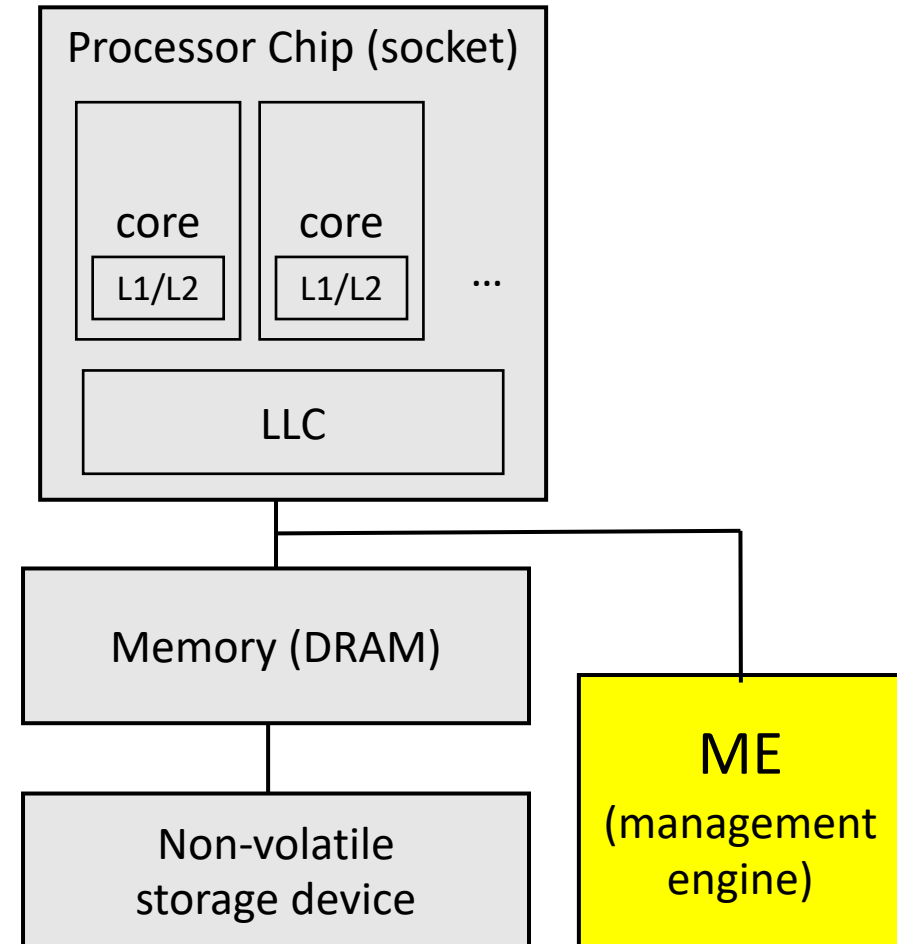
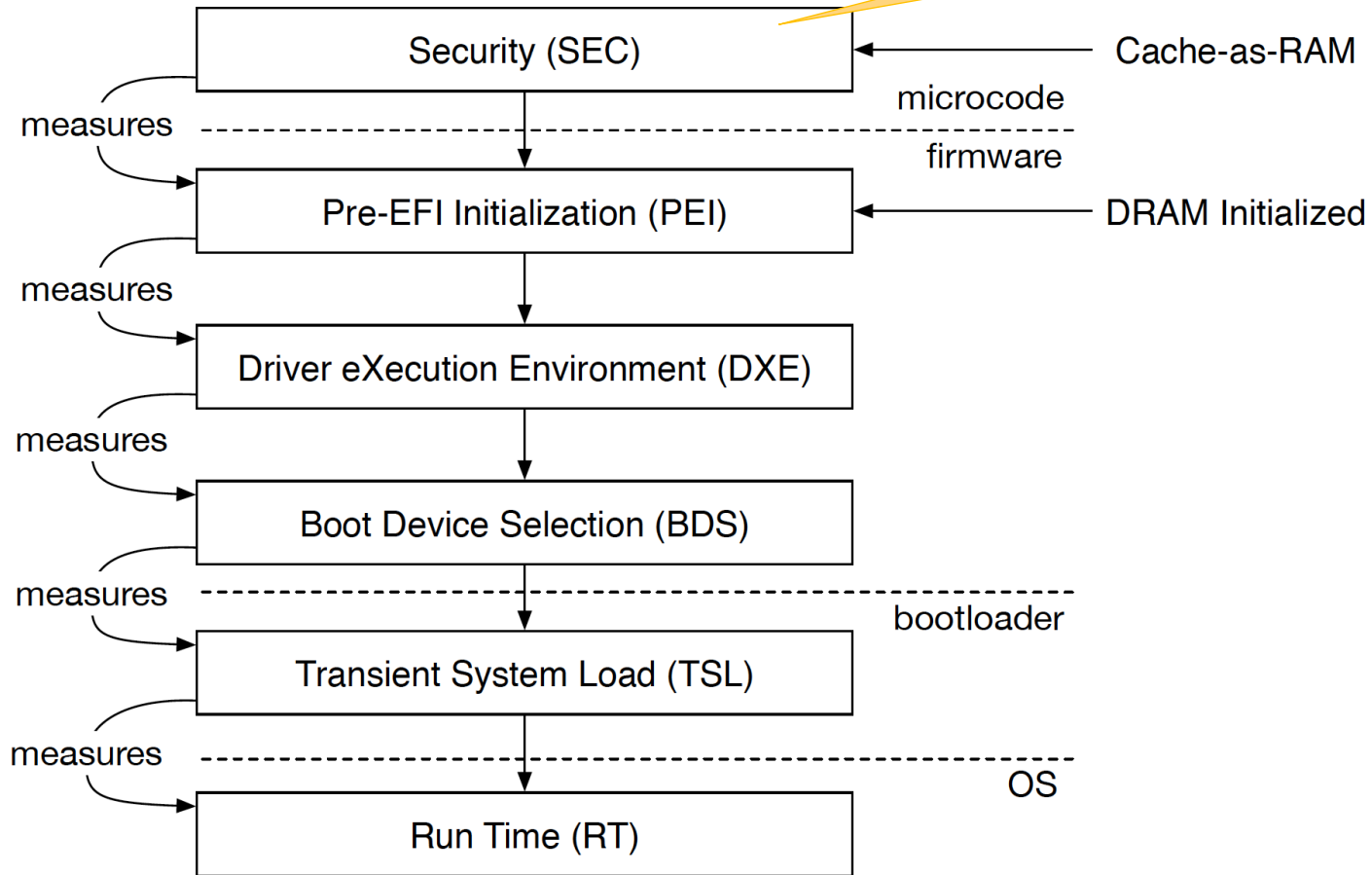
Trusted Platform Module (TPM)

- “Commoditized IBM 4758”
- Standard LPC interface – attaches to commodity motherboards
- Weaker computation capability
- Use cases:
 - Disk encryption and password protection (“seal”)
 - Verify platform integrity (firmware+OS)



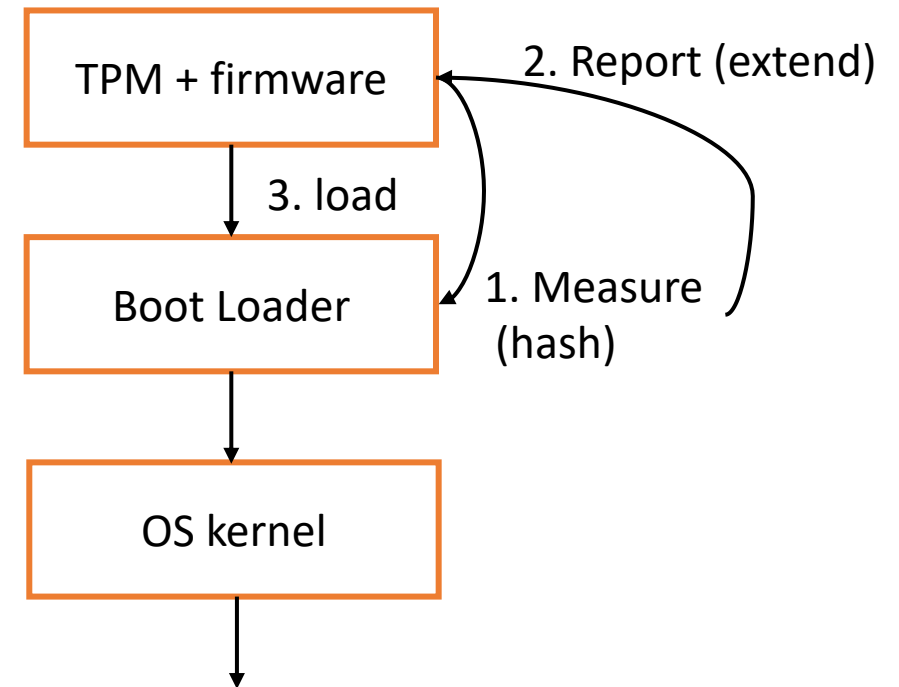
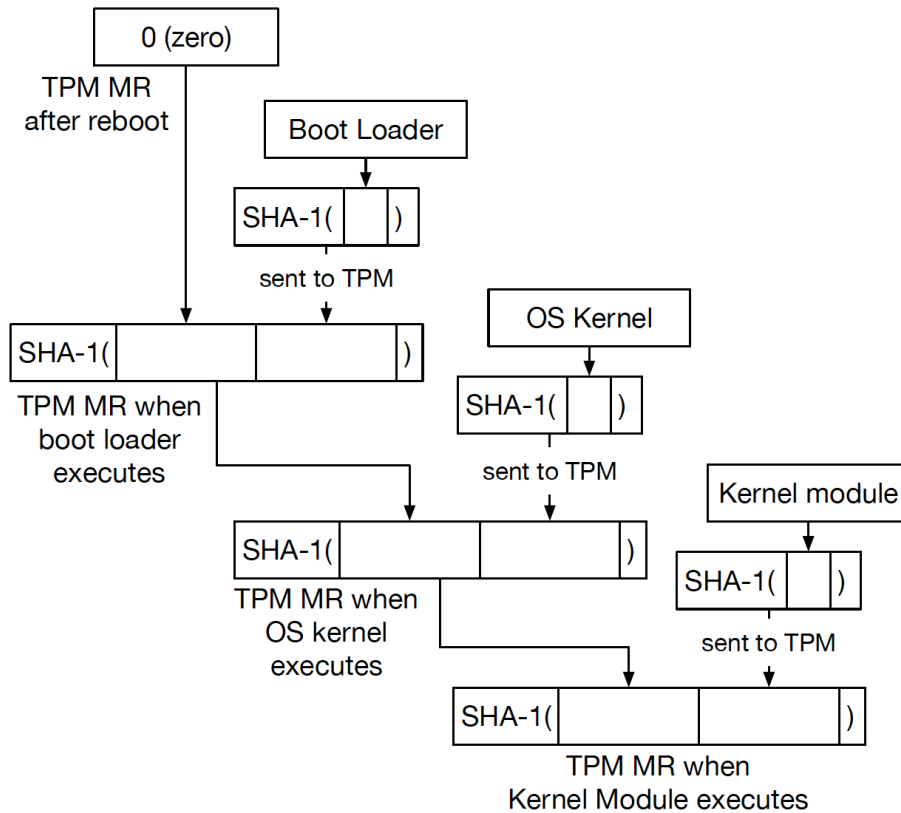
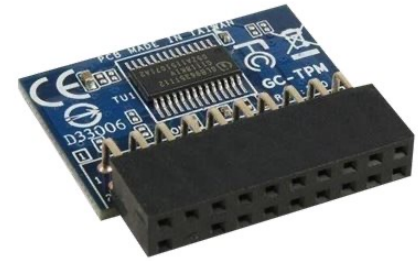
Boot Process (UEFI)

Root of trust



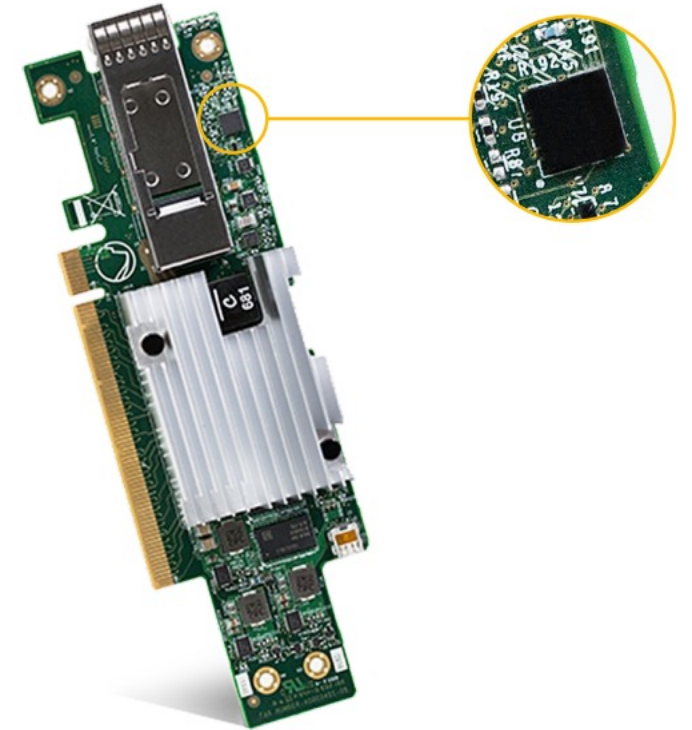
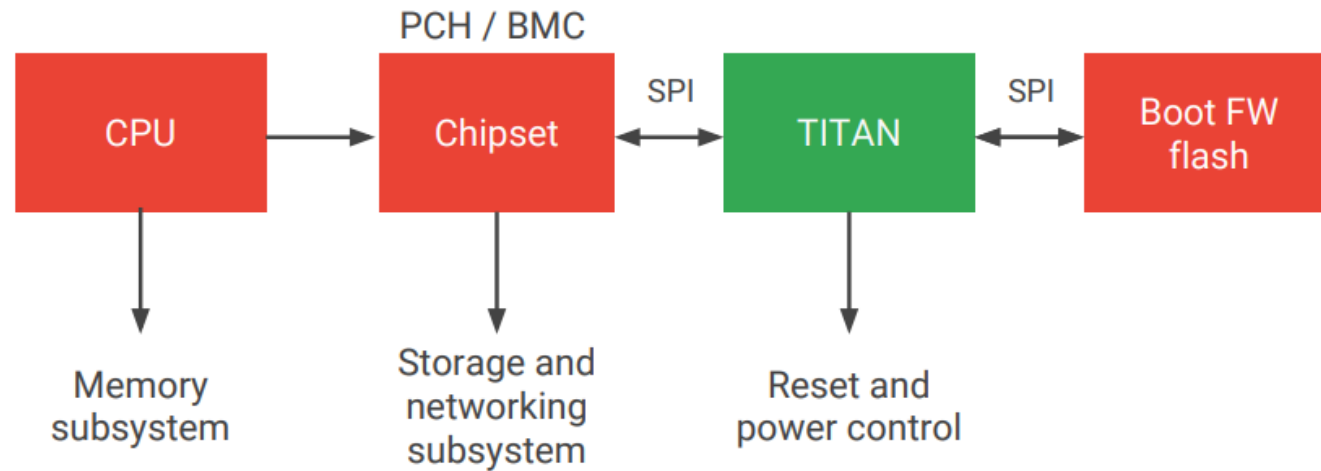
How to perform the measurement?

Secure Boot using TPM



Each step, TPM compares to expected values locally or submitted to a remote attester.

Open-source Choice: Google Titan

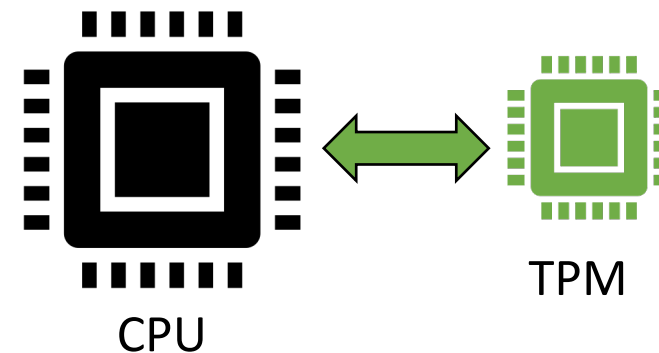
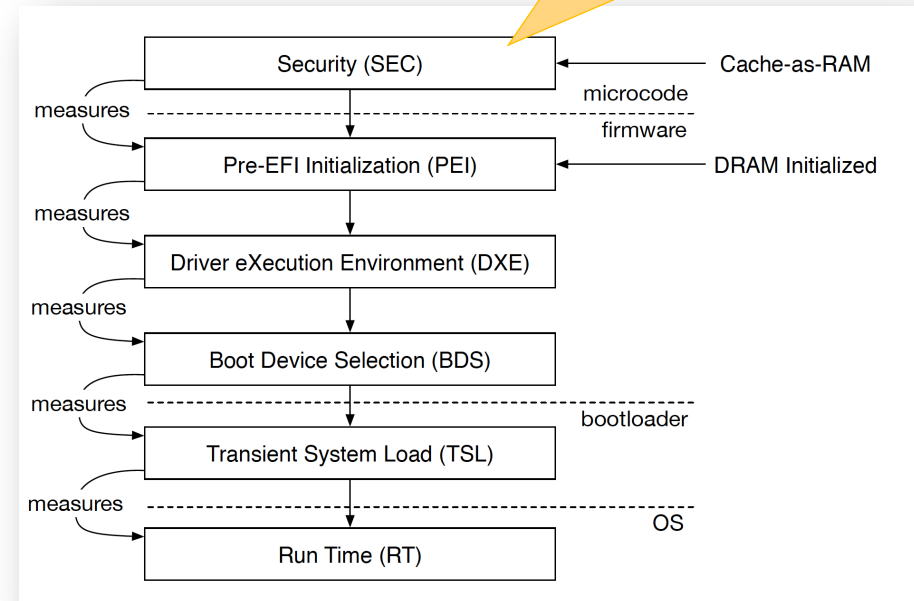


from https://www.hotchips.org/hc30/1conf/1.14_Google_Titan_GoogleFinalTitanHotChips2018.pdf

Security Problems of Using TPM

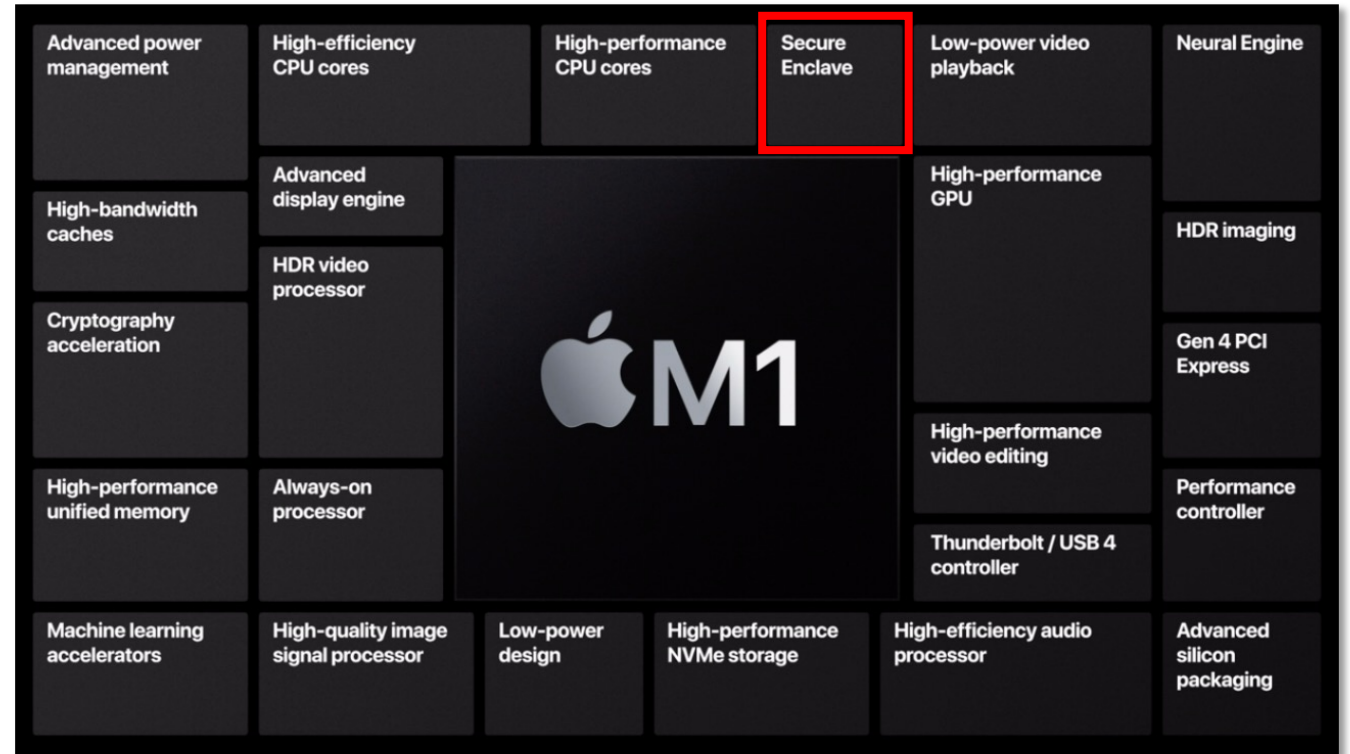
- Assume the first-stage bootloader is securely embedded in motherboard
- Not easy to use with frequent software/kernel update
- Time to check, time to use
- TPM Reset attacks
 - exploiting software vulnerabilities and using software to report false hash values

Root of trust



Apple Secure Enclave

- Additional Goals:
 - Prevent jailbreak
 - Easy to use
- Advantage: one company controls both the hardware and the software



Isolation

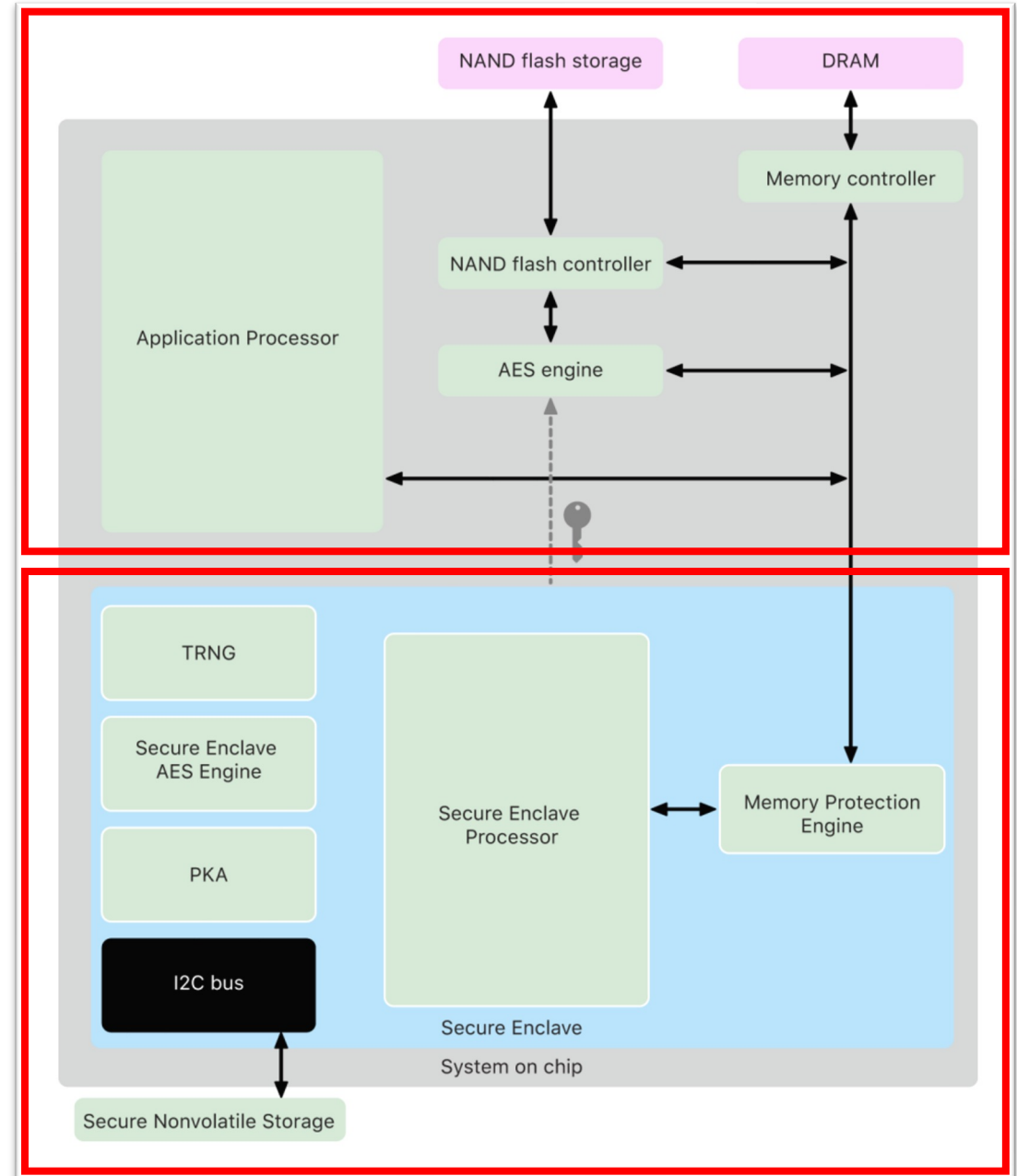
Why separate cores?

Similar to IBM 4758

- Strong isolation
- Block vulnerabilities due to software bugs (running L4 microkernel) and side channels

Different from IBM 4758

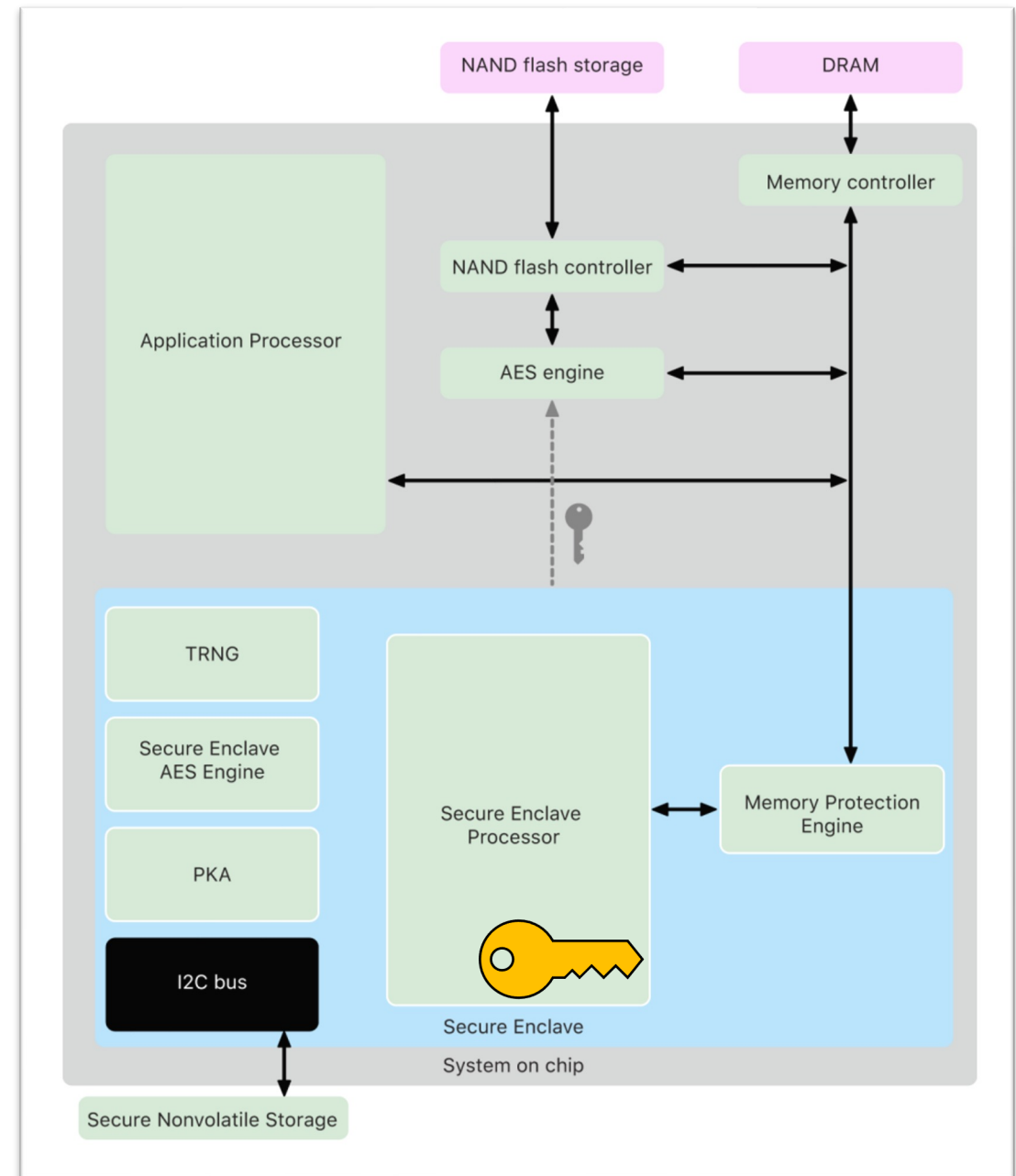
- Not general-purpose, only run secure enclave functionality



Crypto Keys

The Secure Enclave includes a unique ID (**UID**) root cryptographic key.

- Unique to each device
- Randomly generated
- Fused into the SoC at manufacturing time
- Not visible outside the device



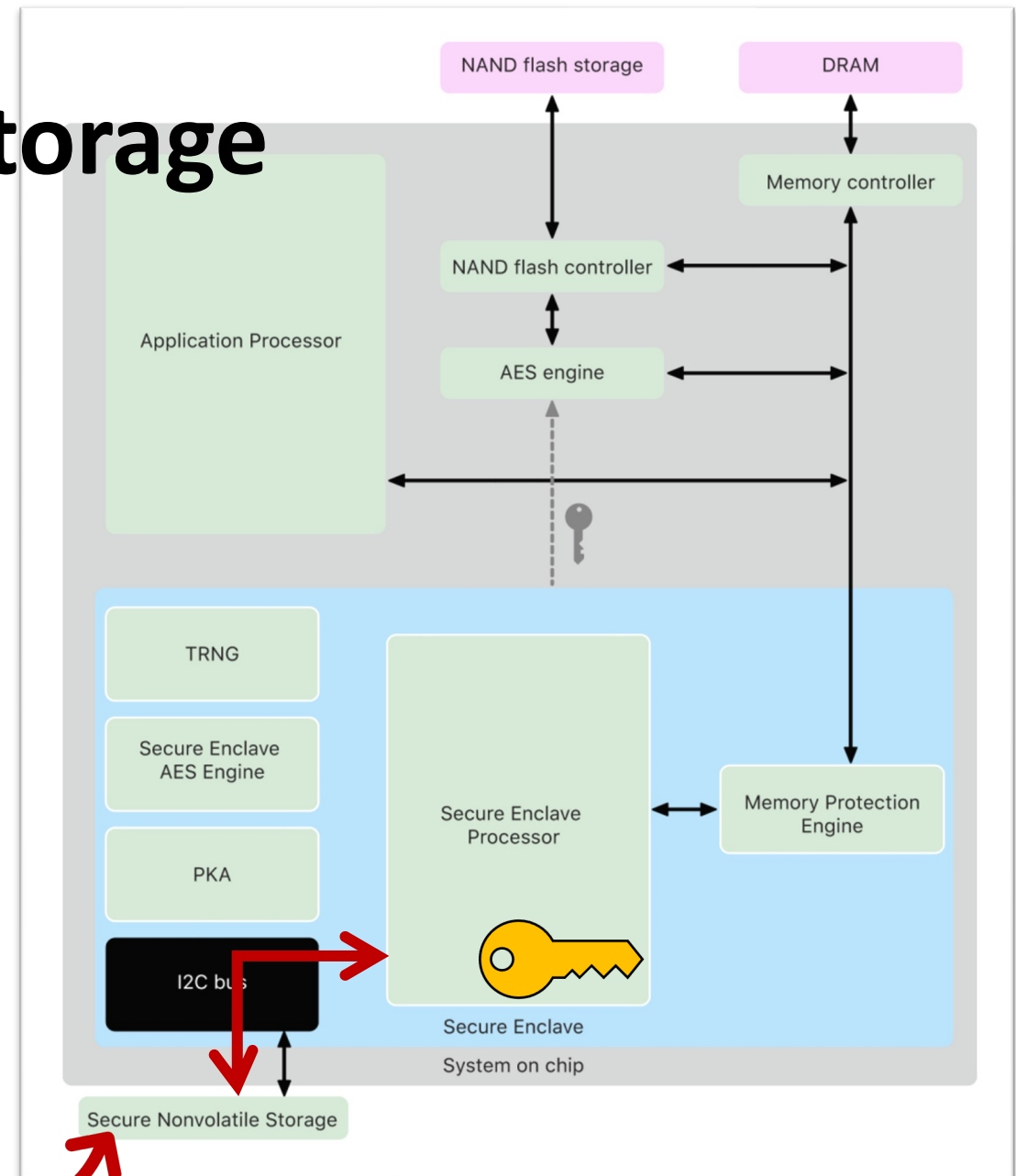
Secure Non-volatile Storage

For easy to use: short passcode. But weaker security?

Passcode + UID -> passcode entropy

Brute-force has to be performed on the device under attack

- Escalating time delays
- Erase data when exceeding attempt count

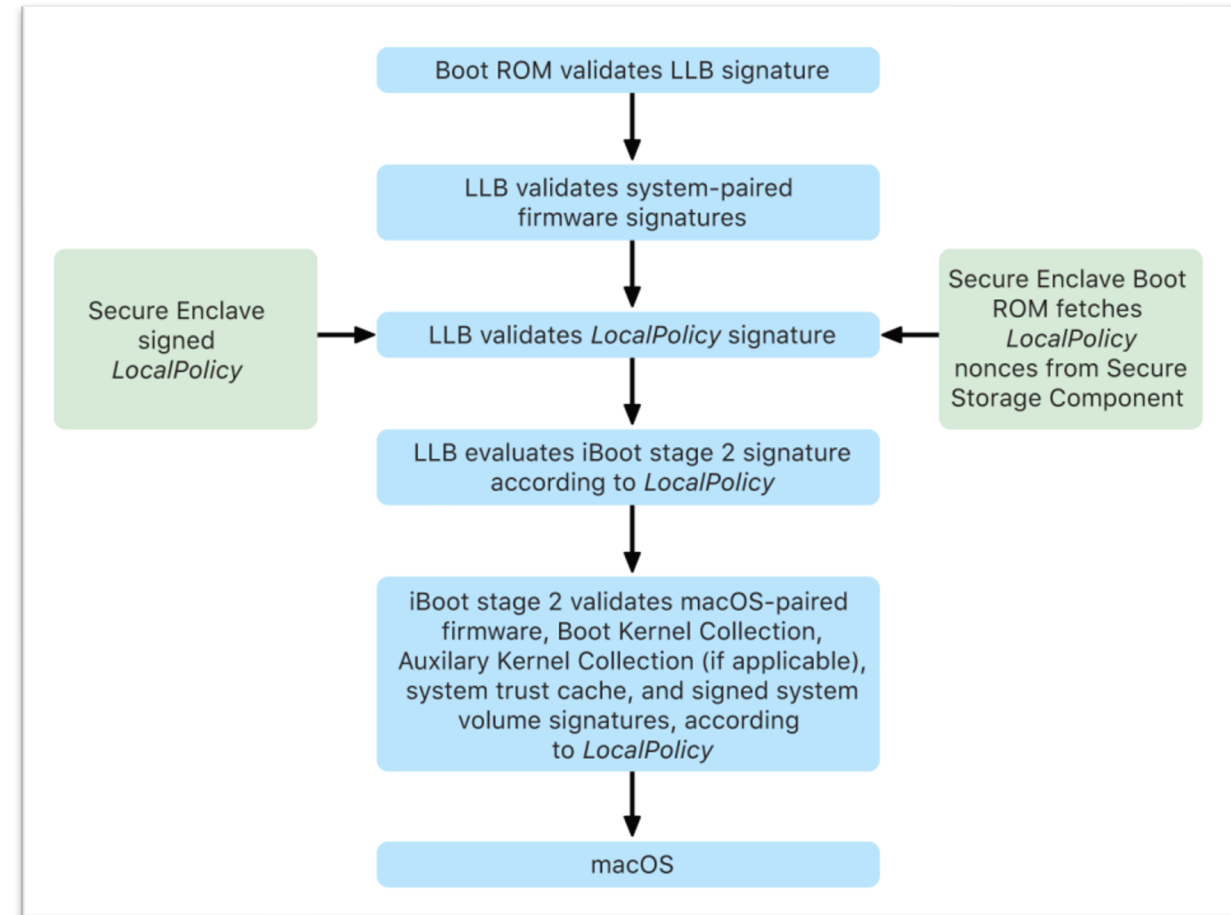


All user data encryption keys

Secure Boot

Similar to TPM but with more constraints

- Each step is signed by Apple to prevent loading non-Apple systems
 - Using Apple Root Certificate authority public key
- Verify more components, including operating system, kernel extensions, etc.
- Keep track of version number to prevent rolling back to older/vulnerable versions



Summary

What Can Hardware Security Modules Offer?

- Establish root of trust
- Bind data and applications with the hardware device
- Offer stronger isolation
- More efficient

Next: Physical Attacks