An Adaptive Cache Coherence Protocol That Implements Sequential Consistency for DSM Systems With Multi-level Caches

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Xiaowei Shen and Arvind

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545 TECHNOLOGY SQUARE, CAMBRIDGE, MASSACHUSETTS 02139
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1 HCN-opt: An Optimized Protocol

The HCN-opt protocol is an optimized protocol based on HCN-base. While most of HCN-base rules remain unchanged in HCN-opt, several HCN-base rules are slightly modified, and a number of new rules are added. HCN-opt incorporates three new protocol messages, Upgrade-rep, Pushout-rep, and Pushout-req.

Upgrade-rep is Ex-rep without data. It can be used to upgrade a shared cell to an exclusive cell. A common scenario is that when a memory receives an Ex-req from a child memory, the directory shows that the child has a shared copy. In HCN-base, the parent memory first invalidates all shared copies cached by its children, and then sends an Ex-rep message to the requesting child memory. An HCN-opt, the parent memory only invalidates, if any, other shared copies cached by its other child memories, and then sends an Upgrade-rep message to the requesting child memory.

Pushout-rep is a combination of Wb-rep and Inv-rep. In HCN-base, when a memory receives an Ex-req, if currently a child memory has the most up-to-date data and the ownership, the parent first sends a Wb-req message to the child to enforce the data to be written back, and then sends an Inv-req message to the child to have the shared copy left be invalidated. In HCN-opt, the parent simply sends a Pushout-req message.

In HCN-opt, the six reply messages form three pairs, Sh-rep/InvRep, Ex-rep/Pushout-rep and Upgrade-rep/Wb-rep, where the two messages in each pair can cause inverse actions each other. The following table gives the initiator (suspended instruction or message) and the expected reply when a request message is issued.

<table>
<thead>
<tr>
<th>Request Issued</th>
<th>Initiator</th>
<th>Expected Reply</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sh-req</td>
<td>Load Sh-req</td>
<td>Sh-req</td>
</tr>
<tr>
<td>Ex-req</td>
<td>Store Ex-req</td>
<td>Ex-ray / Upgrade-rep</td>
</tr>
<tr>
<td>Wb-req</td>
<td>Sh-req Wb-req</td>
<td>Wb-req</td>
</tr>
<tr>
<td>Pushout-req</td>
<td>Ex-ray Pushout-req</td>
<td>Pushout-req</td>
</tr>
<tr>
<td>Inv-req</td>
<td>Ex-ray Pushout-req Inv-req</td>
<td>Inv-req</td>
</tr>
</tbody>
</table>

1.1 Memory Access Rules

1.1.1 Cache-Hit Rules

Read-Cache-Hit Rule

\[ \text{Sys}(\text{id}, \text{Cell}(a,v,(\text{cs},R(e)))) | m, \text{in}, \text{out}, \text{trecs}) \Rightarrow \text{Proc}(ia, rf, prog) \]

if \[ \text{prog}[ia] = r := \text{Load}(r_1) \text{ and } a = rf[r_1] \]

\[ \Rightarrow \text{Sys}(\text{id}, \text{Cell}(a,v,(\text{cs},R(e)))) | m, \text{in}, \text{out}, \text{trecs}) \Rightarrow \text{Proc}(ia+1, rf[r := v], prog) \]

Write-Cache-Hit Rule

\[ \text{Sys}(\text{id}, \text{Cell}(a,u,(\text{Ex},R(e)))) | m, \text{in}, \text{out}, \text{trecs}) \Rightarrow \text{Proc}(ia, rf, prog) \]

if \[ \text{prog}[ia] = \text{Store}(r_1, r_2) \text{ and } a = rf[r_1] \]

\[ \Rightarrow \text{Sys}(\text{id}, \text{Cell}(a,v,(\text{Ex},R(e)))) | m, \text{in}, \text{out}, \text{trecs}) \Rightarrow \text{Proc}(ia+1, rf, prog) \] where \[ v = rf[r_2] \]
<table>
<thead>
<tr>
<th>SYSTEM</th>
<th>MOD</th>
<th>VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sys(MU, EU)</td>
<td>MU</td>
<td>(id, MEM, INQ, OUTQ, TRECS)</td>
</tr>
<tr>
<td>PROC</td>
<td>SG</td>
<td>System</td>
</tr>
<tr>
<td>Cell(a,v,STATE)</td>
<td>MEM</td>
<td>Memory Unit</td>
</tr>
<tr>
<td>CSTATE</td>
<td>STATE</td>
<td>Execution Unit</td>
</tr>
<tr>
<td>Sh</td>
<td>Ex</td>
<td>System Group</td>
</tr>
<tr>
<td>R(DIR)</td>
<td>W(id)</td>
<td>Cell’s State</td>
</tr>
<tr>
<td>e</td>
<td>id</td>
<td>Cell’s State</td>
</tr>
<tr>
<td>e</td>
<td>MSG ⊕ INQ</td>
<td>Directory</td>
</tr>
<tr>
<td>e</td>
<td>MSG ⊕ OUTQ</td>
<td>Incoming Queue</td>
</tr>
<tr>
<td>e</td>
<td>MSG ⊕ OUTQ</td>
<td>Outgoing Queue</td>
</tr>
<tr>
<td>Msg(idsrc, iddest, CMD, a, v)</td>
<td>MSG</td>
<td>Protocol Message</td>
</tr>
<tr>
<td>REPLY</td>
<td>REQUEST</td>
<td>Message Command</td>
</tr>
<tr>
<td>Sh-rep</td>
<td>Ex-rep</td>
<td>Wb-rep</td>
</tr>
<tr>
<td>Upgrade-rep</td>
<td>Pushout-rep</td>
<td></td>
</tr>
<tr>
<td>Sh-rep</td>
<td>Ex-rep</td>
<td>Wb-rep</td>
</tr>
<tr>
<td>Pushout-rep</td>
<td></td>
<td></td>
</tr>
<tr>
<td>e</td>
<td>Trec(a, INITIATOR)</td>
<td>TRECS</td>
</tr>
<tr>
<td>(ia,Load)</td>
<td>(ia,Store)</td>
<td>(id,REQUEST)</td>
</tr>
</tbody>
</table>

Figure 1: The HCN-opt Model (Initially, all non-outermost memories, all incoming and outgoing message queues, and all transient records are empty; the outermost memory contains a cell in the (Ex,R(e)) state for each address)
1.1.2 Cache-Miss Rules

Read-Cache-Miss Rule
Sys((id, m, in, out, trecs), Proc(id, rf, prog))
if prog[i] = r := Load(r1) and rf[r1] \notin m and rf[r] \notin trecs
→ Sys((id, m, in, out \otimes Msg(id, id_p, Sh-req, a, \bot), Trec(a, (ia,Load)) | trecs), Proc(id, rf, prog))
   where id_p = parent(id) and a = rf[r1]

Write-Cache-Miss Rule
Sys((id, m, in, out, trecs), Proc(id, rf, prog))
if prog[i] = Store(r1, r2) and Cell(rf[r1].c,(Ex,.-)) \notin m and rf[r] \notin trecs
→ Sys((id, m, in, out \otimes Msg(id, id_p, Ex-req, a, \bot), Trec(a, (ia,Store)) | trecs), Proc(id, rf, prog))
   where id_p = parent(id) and a = rf[r1]

1.2 Child-to-Parent Request Processing Rules

1.2.1 Sh-Request Rules

Receive-Sh-Req-And-Send-Sh-Rep Rule
(id, Cell(a,v,(cs,R(dir)))) | m, Msg(id_k, id, Sh-req, a, \bot) \otimes in, out, trecs)
if id_k \notin dir and a \notin trecs
→ (id, Cell(a,v,(cs,R(id_k[dir])))) | m, in, out \otimes Msg(id, id_k, Sh-req, a, v), trecs)

Receive-Sh-Req-And-Send-Wb-Req Rule
(id, Cell(a,v,(Ex,W(id)))) | m, Msg(id_k, id, Sh-req, a, \bot) \otimes in, out, trecs)
if id_k \neq idj and a \notin trecs
→ (id, Cell(a,v,(Ex,W(id)))) | m, in, out \otimes Msg(id, id_k, Wb-req, a, \bot), Trec(a, (id_k,Sh-req)) | trecs)

Receive-Sh-Req-And-Send-Sh-Req Rule
(id, m, Msg(id_k, id, Sh-req, a, \bot) \otimes in, out, trecs) if a \notin m and a \notin trecs
→ (id, m, in, out \otimes Msg(id, id_p, Sh-req, a, \bot), Trec(a, (id_k,Sh-req)) | trecs)
   where id_p = parent(id)

1.2.2 Ex-Request Rules

Receive-Ex-Req-And-Send-Ex-Rep Rule
(id, Cell(a,v,(Ex,R(v)))) | m, Msg(id_k, id, Ex-req, a, \bot) \otimes in, out, trecs) if a \notin trecs
→ (id, Cell(a,v,(Ex,W(id)))) | m, in, out \otimes Msg(id, id_k, Ex-req, a, v), trecs)

Receive-Ex-Req-And-Send-Upgrade-Rep Rule
(id, Cell(a,v,(Ex,R(id_k)))) | m, Msg(id_k, id, Ex-req, a, \bot) \otimes in, out, trecs) if a \notin trecs
→ (id, Cell(a,v,(Ex,W(id_k)))) | m, in, out \otimes Msg(id, id_k, Upgrade-req, a, v), trecs)

Receive-Ex-Req-And-Multicast-Inv-Req Rule
(id, Cell(a,v,(Ex,R(dir)))) | m, Msg(id_k, id, Ex-req, a, \bot) \otimes in, out, trecs)
if dir-\{id_k\} \neq \epsilon and a \notin trecs
→ (id, Cell(a,v,(Ex,R(dir)))) | m, in, out \otimes multicast(id, dir-\{id_k\}, Inv-req, a, \bot), Trec(a, (id_k,Ex-req)) | trecs)

Receive-Ex-Req-And-Send-Pushout-Req Rule
(id, Cell(a,v,(Ex,W(id)))) | m, Msg(id_k, id, Ex-req, a, \bot) \otimes in, out, trecs)
if id_k \neq id and a \notin trecs
→ (id, Cell(a,v,(Ex,W(id)))) | m, in, out \otimes Msg(id, id_j, Pushout-req, a, \bot), Trec(a, (id_k,Ex-req)) | trecs)
Receive-Ex-Req-And-Send-Ex-Req Rule
\[ \{ \text{id, m, Msg(\text{id}, \text{id}, \text{Ex-req}, \text{a}, \bot) \odot \text{in, out, trecs)} } \]
\[ \text{if Cell(a,\text{Ex-req})} \notin \text{m and a} \notin \text{trecs} \]
\[ \rightarrow \{ \text{id, m, in, out} \odot \text{Msg(\text{id}, \text{id}, \text{Ex-req}, \text{a}, \bot), Trec(a, (\text{id}, \text{Ex-req})) | trecs} \}
\text{where \text{id} = parent(id)} \]

1.3 Parent-to-Child Request Processing Rules

1.3.1 Wb-Request Rules

Receive-Wb-Req-And-Send-Wb-Req Rule
\[ \{ \text{id, Cell(a,\text{Wb-req}) | m, Msg(\text{id}, \text{id}, \text{Wb-req}, \text{a}, \bot) \odot \text{in, out, trecs)} } \]
\[ \text{if a} \notin \text{trecs} \]
\[ \rightarrow \{ \text{id, Cell(a,\text{Wb-req}) | m, in, out} \odot \text{Msg(\text{id}, \text{id}, \text{Wb-req}, \text{a}, \bot), Trec(a, (\text{id}, \text{Wb-req})) | trecs} \}

1.3.2 Pushout-Request Rules

Receive-Pushout-Req-And-Send-Pushout-Rep Rule
\[ \{ \text{id, Cell(a,\text{Pushout-req}) | m, Msg(\text{id}, \text{id}, \text{Pushout-req}, \text{a}, \bot) \odot \text{in, out, trecs)} } \]
\[ \text{if a} \notin \text{trecs} \]
\[ \rightarrow \{ \text{id, m, in, out} \odot \text{Msg(\text{id}, \text{id}, \text{Pushout-req}, \text{a}, \bot), Trec(a, (\text{id}, \text{Pushout-req})) | trecs} \}

Receive-Pushout-Req-And-Multicast-Inv-Req Rule
\[ \{ \text{id, Cell(a,\text{Pushout-req}) | m, Msg(\text{id}, \text{id}, \text{Pushout-req}, \text{a}, \bot) \odot \text{in, out, trecs)} } \]
\[ \text{if dir} \neq \epsilon \text{ and a} \notin \text{trecs} \]
\[ \rightarrow \{ \text{id, Cell(a,\text{Pushout-req}) | m, in, out} \odot \text{multicast(id, dir, Inv-req, a, \bot), Trec(a, (\text{id}, \text{Pushout-req})) | trecs} \}

Receive-Pushout-Req-And-Send-Pushout-Req Rule
\[ \{ \text{id, Cell(a,\text{Pushout-req}) | m, Msg(\text{id}, \text{id}, \text{Pushout-req}, \text{a}, \bot) \odot \text{in, out, trecs)} } \]
\[ \text{if a} \notin \text{trecs} \]
\[ \rightarrow \{ \text{id, Cell(a,\text{Pushout-req}) | m, in, out} \odot \text{Msg(\text{id}, \text{id}, \text{Pushout-req}, \text{a}, \bot), Trec(a, (\text{id}, \text{Pushout-req})) | trecs} \}

1.3.3 Inv-Request Rules

Receive-Inv-Req-And-Send-Inv-Rep Rule
\[ \{ \text{id, Cell(a,\text{Inv-req}) | m, Msg(\text{id}, \text{id}, \text{Inv-req}, \text{a}, \bot) \odot \text{in, out, trecs)} } \]
\[ \rightarrow \{ \text{id, m, in, out} \odot \text{Msg(\text{id}, \text{id}, \text{Inv-req}, \text{a}, \bot), trecs} \}

Receive-Inv-Req-And-Multicast-Inv-Req Rule
\[ \{ \text{id, Cell(a,\text{Inv-req}) | m, Msg(\text{id}, \text{id}, \text{Inv-req}, \text{a}, \bot) \odot \text{in, out, trecs)} } \]
\[ \text{if dir} \neq \epsilon \]
\[ \rightarrow \{ \text{id, Cell(a,\text{Inv-req}) | m, in, out} \odot \text{multicast(id, dir, Inv-req, a, \bot), Trec(a, (\text{id}, \text{Inv-req})) | trecs} \} \]
1.4 Parent-to-Child Reply Processing Rules

1.4.1 Sh-Reply Rules

Receive-Sh-Rep-And-Execute-Load Rule
\[\text{Sys}(\text{id}, \text{m}, \text{Msg}(\text{id}_p, \text{id}, \text{Sh-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia,Load})) \mid \text{trecs}), \text{Proc(ia, rf, prog)})\]
if \(\text{prog[ia]} = r := \text{Load}(r_1) \text{ and } a = \text{rf}[r_1]\)
\[\rightarrow \text{Sys}(\text{id}, \text{Cell}(a,v,(\text{Sh,R}(r))) \mid \text{m, in, out, trecs}), \text{Proc(ia+1, rf[r := v], prog)})\]

Receive-Sh-Rep-And-Send-Sh-Rep Rule
\[\text{id}, \text{m}, \text{Msg}(\text{id}_p, \text{id}, \text{Sh-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{id}_k, \text{Sh-req})) \mid \text{trecs}\]
\[\rightarrow \text{id}, \text{Cell}(a,v,(\text{Sh,R(id}_k))) \mid \text{m, in, out} \odot \text{Msg}(\text{id}, \text{id}_k, \text{Sh-rep}, a, v), \text{trecs}\]

1.4.2 Ex-Reply Rules

Receive-Ex-Rep-And-Execute-Store Rule
\[\text{Sys}(\text{id}, \text{m}, \text{Msg}(\text{id}_p, \text{id}, \text{Ex-rep}, a, u) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia,Store})) \mid \text{trecs}), \text{Proc(ia, rf, prog)})\]
if \(\text{prog[ia]} = \text{Store}(r_1, r_2) \text{ and } a = \text{rf}[r_1]\)
\[\rightarrow \text{Sys}(\text{id}, \text{Cell}(a,v,(\text{Ex,R}(e))) \mid \text{m, in, out, trecs}), \text{Proc(ia+1, rf, prog)}) \text{ where } v = \text{rf}[r_2]\]

Receive-Ex-Rep-And-Send-Ex-Rep Rule
\[\text{id}, \text{m}, \text{Msg}(\text{id}_p, \text{id}, \text{Ex-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{id}_k, \text{Ex-req})) \mid \text{trecs}\]
\[\rightarrow \text{id}, \text{Cell}(a,v,(\text{Ex,W(id}_k))) \mid \text{m, in, out} \odot \text{Msg}(\text{id}, \text{id}_k, \text{Ex-rep}, a, v), \text{trecs}\]

1.4.3 Upgrade-Reply Rules

Receive-Upgrade-Rep-And-Execute-Store Rule
\[\text{Sys}(\text{id}, \text{Cell}(a,u,(\text{Sh,R}(e))) \mid \text{m, Msg(id}_p, \text{id}, \text{Upgrade-rep}, a, \perp) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia,Store})) \mid \text{trecs}), \text{Proc(ia, rf, prog)})\]
\[\rightarrow \text{Sys}(\text{id}, \text{Cell}(a,v,(\text{Ex,R}(e))) \mid \text{m, in, out, trecs}), \text{Proc(ia+1, rf, prog)}) \text{ where } v = \text{rf}[r_2]\]

Receive-Upgrade-Rep-And-Send-Upgrade-Rep Rule
\[\text{id}, \text{Cell}(a,v,(\text{Sh,R}(e))) \mid \text{m, Msg(id}_p, \text{id}, \text{Upgrade-rep}, a, \perp) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{id}_k, \text{Ex-req})) \mid \text{trecs}\]
\[\rightarrow \text{id}, \text{Cell}(a,v,(\text{Ex,W(id}_k))) \mid \text{m, in, out} \odot \text{Msg}(\text{id}, \text{id}_k, \text{Ex-rep}, a, v), \text{trecs}\]

Receive-Upgrade-Rep-And-Multicast-Inv-Req Rule
\[\text{id}, \text{Cell}(a,v,(\text{Sh,R(dir)}) \mid \text{m, Msg(id}_p, \text{id}, \text{Upgrade-rep}, a, \perp) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{id}_k, \text{Ex-req})) \mid \text{trecs})\]
\[\text{if} \text{ dir-\{id}_k\} \neq \epsilon\]
\[\rightarrow \text{id}, \text{Cell}(a,v,(\text{Ex,R(dir)}) \mid \text{m, in, out} \odot \text{multicast(id, dir-\{id}_k\), Inv-req, a, \perp}), \text{Trec}(a, (\text{id}_k, \text{Ex-req})) \mid \text{trecs})\]

1.5 Child-to-Parent Reply Processing Rules

1.5.1 Wb-Reply Rules

Receive-Wb-Rep-And-Send-Sh-Rep Rule
\[\text{id}, \text{Cell}(a,u,(\text{Ex,W(id}_k))) \mid \text{m, Msg(id}_k, \text{id}, \text{Wb-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{id}_j, \text{Sh-req})) \mid \text{trecs}\]
\[\rightarrow \text{id}, \text{Cell}(a,v,(\text{Ex,R(id}_k,id))) \mid \text{m, in, out} \odot \text{Msg}(\text{id}, \text{id}_j, \text{Sh-rep}, a, v), \text{trecs}\]
Receive-Wb-Rep-And-Send-Wb-Rep Rule
(id, Cell(a,u,(Ex,W(id_k)))) | m, Msg(id_k, id, Wb-rep, a, v) ⊙ in, out, Trec(a, (id_p,Wb-req)) | treCs
→ (id, Cell(a,v,(Sh,R(id_k)))) | m, in, out ⊗ Msg(id, id_p, Wb-rep, a, v), treCs

1.5.2 Pushout-Reply Rules

Receive-Pushout-Rep-And-Send-Ex-Rep Rule
(id, Cell(a,u,(Ex,W(id_k)))) | m, Msg(id_k, id, Pushout-rep, a, v) ⊙ in, out, Trec(a, (id_j,Ex-req)) | treCs
→ (id, Cell(a,v,(Ex,W(id_j)))) | m, in, out ⊗ Msg(id, id_p, Pushout-rep, a, v), treCs

Receive-Pushout-Rep-And-Send-Pushout-Rep Rule
(id, Cell(a,u,(Ex,W(id_k)))) | m, Msg(id_k, id, Pushout-rep, a, v) ⊙ in, out, Trec(a, (id_p,Pushout-req)) | treCs
→ (id, m, in, out ⊗ Msg(id, id_p, Pushout-rep, a, v), treCs)

1.5.3 Inv-Reply Rules

Receive-Inv-Rep-With-Ex-Req-Suspended Rule
(id, Cell(a,v,(Ex,R(id_k,dir)))) | m, Msg(id_k, id, Inv-rep, a, ⊥) ⊙ in, out, Trec(a, (id_j,Ex-req)) | treCs
if dir - {id_j} ≠ ε
→ (id, Cell(a,v,(Ex,R(dir)))) | m, in, out, Trec(a, (id_j,Ex-req)) | treCs

Receive-Inv-Rep-With-Pushout-Req-Suspended Rule
(id, Cell(a,v,(Ex,R(id_k,dir)))) | m, Msg(id_k, id, Inv-rep, a, ⊥) ⊙ in, out, Trec(a, (id_p,Pushout-req)) | treCs
if dir ≠ ε
→ (id, Cell(a,v,(Ex,R(dir)))) | m, in, out, Trec(a, (id_p,Pushout-req)) | treCs

Receive-Inv-Rep-With-Inv-Req-Suspended Rule
(id, Cell(a,v,(Sh,R(id_k,dir)))) | m, Msg(id_k, id, Inv-rep, a, ⊥) ⊙ in, out, Trec(a, (id_p,Inv-req)) | treCs
if dir ≠ ε
→ (id, Cell(a,v,(Sh,R(dir)))) | m, in, out, Trec(a, (id_p,Inv-req)) | treCs

Receive-Inv-Rep-And-Send-Ex-Rep Rule
(id, Cell(a,v,(Ex,R(id_k)))) | m, Msg(id_k, id, Inv-rep, a, ⊥) ⊙ in, out, Trec(a, (id_j,Ex-req)) | treCs
→ (id, Cell(a,v,(Ex,W(id_j)))) | m, in, out ⊗ Msg(id, id_j, Ex-rep, a, v), treCs

Receive-Inv-Rep-And-Send-Upgrade-Rep Rule
(id, Cell(a,v,(Ex,R(id_k,dir)))) | m, Msg(id_k, id, Inv-rep, a, ⊥) ⊙ in, out, Trec(a, (id_j,Ex-req)) | treCs
→ (id, Cell(a,v,(Ex,W(id_j)))) | m, in, out ⊗ Msg(id, id_j, Upgrade-rep, a, ⊥), treCs

Receive-Inv-Rep-And-Send-Pushout-Rep Rule
(id, Cell(a,v,(Ex,R(id_k)))) | m, Msg(id_k, id, Inv-rep, a, ⊥) ⊙ in, out, Trec(a, (id_p,Pushout-req)) | treCs
→ (id, m, in, out ⊗ Msg(id, id_p, Pushout-rep, a, v), treCs)

Receive-Inv-Rep-And-Send-Inv-Rep Rule
(id, Cell(a,v,(Sh,R(id_k)))) | m, Msg(id_k, id, Inv-rep, a, ⊥) ⊙ in, out, Trec(a, (id_p,Inv-req)) | treCs
→ (id, m, in, out ∙ Msg(id, id_p, Inv-rep, a, ⊥), treCs)
1.6 Message Passing Rules

Message-Passing-To-Child Rule

\[ \text{Sys}(\text{id}, \text{m}, \text{in}, \text{Msg}(\text{id}, \text{id}_k, \text{cmd}, \text{a}, \text{v}) \otimes \text{out}, \text{trecs}), \text{Sys}(\text{id}_k, \text{m}_k, \text{in}_k, \text{out}_k, \text{trecs}_k), \text{eu}_k) | \text{sg} \]

\[ \rightarrow \text{Sys}(\text{id}, \text{m}, \text{in}, \text{out}, \text{trecs}), \text{Sys}(\text{id}_k, \text{m}_k, \text{in}_k \odot \text{Msg}(\text{id}, \text{id}_k, \text{cmd}, \text{a}, \text{v}), \text{out}_k, \text{trecs}_k), \text{eu}_k) | \text{sg} \]

Message-Passing-To-Parent Rule

\[ \text{Sys}(\text{id}, \text{m}, \text{in}, \text{out}, \text{trecs}), \text{Sys}(\text{id}_k, \text{m}_k, \text{in}_k, \text{Msg}(\text{id}_k, \text{id}, \text{cmd}, \text{a}, \text{v}) \otimes \text{out}_k, \text{trecs}_k), \text{eu}_k) | \text{sg} \]

\[ \rightarrow \text{Sys}(\text{id}, \text{m}, \text{in} \odot \text{Msg}(\text{id}_k, \text{id}, \text{cmd}, \text{a}, \text{v}), \text{out}, \text{trecs}), \text{Sys}(\text{id}_k, \text{m}_k, \text{in}_k, \text{out}_k, \text{trecs}_k), \text{eu}_k) | \text{sg} \]
2 HCN-adpt: An Adaptive Protocol

The HCN-adpt protocol is an adaptive protocol based on HCN-opt. While all HCN-opt rules remain unchanged, new rules are added to incorporate adaptivity, which allows protocol messages to be issued voluntarily, i.e., without initiator. It is worth noting that voluntary Pushout-rep and Inv-rep effectively model cache line replacements due to capacity or associativity conflict, while voluntary Sh-rep and Ex-rep behave like data prefetch.

The newly added rules fall into four categories: voluntary reply issue rules, voluntary request issue rules, unexpected reply processing rules, and unexpected request processing rules.

2.1 Memory Access Rules

2.1.1 Cache-Hit Rules

*Read-Cache-Hit Rule*

\[
\text{Sys}(id, \text{Cell}(a,v,(cs,R(e)))) \mid m, \text{in}, \text{out}, \text{trecs}), \text{Proc}(ia, rf, prog))
\]

\[
\text{if } \text{prog}[ia] = r := \text{Load}(r_1) \text{ and } a = rf[r_1]
\]

\[
\rightarrow \text{Sys}(id, \text{Cell}(a,v,(cs,R(e)))) \mid m, \text{in}, \text{out}, \text{trecs}), \text{Proc}(ia+1, rf[r := v], prog))
\]

*Write-Cache-Hit Rule*

\[
\text{Sys}(id, \text{Cell}(a,u,(Ex,R(e)))) \mid m, \text{in}, \text{out}, \text{trecs}), \text{Proc}(ia, rf, prog))
\]

\[
\text{if } \text{prog}[ia] = \text{Store}(r_1, r_2) \text{ and } a = rf[r_1]
\]

\[
\rightarrow \text{Sys}(id, \text{Cell}(a,v,(Ex,R(e)))) \mid m, \text{in}, \text{out}, \text{trecs}), \text{Proc}(ia+1, rf, prog)) \text{ where } v = rf[r_2]
\]

2.1.2 Cache-Miss Rules

*Read-Cache-Miss Rule*

\[
\text{Sys}(id, m, \text{in}, \text{out}, \text{trecs}), \text{Proc}(ia, rf, prog))
\]

\[
\text{if } \text{prog}[ia] = r := \text{Load}(r_1) \text{ and } rf[r_1] \notin m \text{ and } rf[r_1] \notin \text{trecs}
\]

\[
\rightarrow \text{Sys}(id, m, \text{in}, \text{out} \otimes \text{Msg}(id, id_p, \text{Sh-rep}, a, \bot), \text{Trec}(a,(ia,\text{Load})) \mid \text{trecs}), \text{Proc}(ia, rf, prog))
\]

\[
\text{where } id_p = \text{parent}(id) \text{ and } a = rf[r_1]
\]

*Write-Cache-Miss Rule*

\[
\text{Sys}(id, m, \text{in}, \text{out}, \text{trecs}), \text{Proc}(ia, rf, prog))
\]

\[
\text{if } \text{prog}[ia] = \text{Store}(r_1, r_2) \text{ and } \text{Cell}(rf[r_1], -(Ex,r)) \notin m \text{ and } rf[r_1] \notin \text{trecs}
\]

\[
\rightarrow \text{Sys}(id, m, \text{in}, \text{out} \otimes \text{Msg}(id, id_p, \text{Ex-rep}, a, \bot), \text{Trec}(a,(ia,\text{Store})) \mid \text{trecs}), \text{Proc}(ia, rf, prog))
\]

\[
\text{where } id_p = \text{parent}(id) \text{ and } a = rf[r_1]
\]

2.2 Voluntary Reply Issue Rules

*Send-Voluntary-Sh-Rep Rule*

\[
(id, \text{Cell}(a,v,(cs,R(dir)))) \mid m, \text{in}, \text{out}, \text{trecs})
\]

\[
\text{if } id_k \in \text{children}(id) \text{ and } id_k \notin \text{dir} \text{ and } a \notin \text{trecs}
\]

\[
\rightarrow (id, \text{Cell}(a,v,(cs,R(id_k|dir)))) \mid m, \text{in}, \text{out} \otimes \text{Msg}(id, id_k, \text{Sh-rep}, a, v), \text{trecs})
\]
Send-Voluntary-Ex-Rep Rule
\[
(id, Cell(a,v,(Ex,R(\epsilon)))) | m, \text{in, out, trecs}) \quad \text{if} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, Cell(a,v,(Ex,W(id_k)))) | m, \text{in, out} \otimes \text{Msg}(id, id_k, Ex-rep, a, v), \text{trecs})
\]
\[
\text{where} \quad id_k \in \text{children}(id)
\]

Send-Voluntary-Wb-Rep Rule
\[
(id, Cell(a,v,(Ex,R(dir)))) | m, \text{in, out, trecs}) \quad \text{if} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, Cell(a,v,(Sh,R(dir)))) | m, \text{in, out} \otimes \text{Msg}(id, id_p, Wb-rep, a, v), \text{trecs})
\]
\[
\text{where} \quad id_p = \text{parent}(id)
\]

Send-Voluntary-Pushout-Rep Rule
\[
(id, Cell(a,v,(Ex,R(\epsilon)))) | m, \text{in, out, trecs}) \quad \text{if} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, m, \text{in, out} \otimes \text{Msg}(id, id_p, Pushout-rep, a, v), \text{trecs}) \quad \text{where} \quad id_p = \text{parent}(id)
\]

Send-Voluntary-Inv-Rep Rule
\[
(id, Cell(a,v,(Sh,R(\epsilon)))) | m, \text{in, out, trecs}) \quad \text{if} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, m, \text{in, out} \otimes \text{Msg}(id, id_p, Inv-rep, a, \perp), \text{trecs}) \quad \text{where} \quad id_p = \text{parent}(id)
\]

2.3 Voluntary Request Issue Rules

Send-Voluntary-Sh-Req Rule
\[
(id, m, \text{in, out, trecs}) \quad \text{if} \quad \alpha \notin m \quad \text{and} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, m, \text{in, out} \otimes \text{Msg}(id, id_p, Sh-req, a, \perp), \text{trecs}) \quad \text{where} \quad id_p = \text{parent}(id)
\]

Send-Voluntary-Ex-Req Rule
\[
(id, m, \text{in, out, trecs}) \quad \text{if} \quad Cell(a,v,(Ex,-)) \notin m \quad \text{and} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, m, \text{in, out} \otimes \text{Msg}(id, id_p, Ex-req, a, \perp), \text{trecs}) \quad \text{where} \quad id_p = \text{parent}(id)
\]

Send-Voluntary-Wb-Req Rule
\[
(id, Cell(a,v,(Ex,W(id_k)))) | m, \text{in, out, trecs}) \quad \text{if} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, Cell(a,v,(Ex,W(id_k)))) | m, \text{in, out} \otimes \text{Msg}(id, id_k, Wb-req, a, \perp), \text{trecs})
\]

Send-Voluntary-Pushout-Req Rule
\[
(id, Cell(a,v,(Ex,W(id_k)))) | m, \text{in, out, trecs}) \quad \text{if} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, Cell(a,v,(Ex,W(id_k)))) | m, \text{in, out} \otimes \text{Msg}(id, id_k, Pushout-req, a, \perp), \text{trecs})
\]

Send-Voluntary-Inv-Req Rule
\[
(id, Cell(a,v,(Sh,R(dir)))) | m, \text{in, out, trecs}) \quad \text{if} \quad \text{dir} \neq \epsilon \quad \text{and} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, Cell(a,v,(Sh,R(dir)))) | m, \text{in, out} \otimes \text{Msg}(id, id_k, Inv-req, a, \perp), \text{trecs}) \quad \text{where} \quad id_k \in \text{dir}
\]

2.4 Child-to-Parent Request Processing Rules

2.4.1 Sh-Request Rules

Receive-Sh-Req-And-Send-Sh-Req Rule
\[
(id, Cell(a,v,(cs,R(dir)))) | m, \text{Msg}(id_k, id, Sh-req, a, \perp) \otimes \text{in, out, trecs})
\]
\[
\text{if} \quad id_k \neq \text{dir} \quad \text{and} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, Cell(a,v,(cs,R(id_k,dir)))) | m, \text{in, out} \otimes \text{Msg}(id, id_k, Sh-rep, a, v), \text{trecs})
\]

Receive-Sh-Req-And-Send-Wb-Req Rule
\[
(id, Cell(a,v,(Ex,W(id_j)))) | m, \text{Msg}(id_k, id, Sh-req, a, \perp) \otimes \text{in, out, trecs})
\]
\[
\text{if} \quad id_k \neq id_j \quad \text{and} \quad \alpha \notin \text{trecs}
\]
\[
\rightarrow (id, Cell(a,v,(Ex,W(id_j)))) | m, \text{in, out} \otimes \text{Msg}(id, id_j, Wb-req, a, \perp), \text{Trec}(a, (id_k,Sh-req)) \otimes \text{trecs})
\]
Receive-Sh-Req-And-Send-Sh-Req Rule

\[(id, m, \text{Msg}(id_k, id, \text{Sh-req}, a, \bot) \odot \text{in, out, trecs}) \quad \text{if} \ a \notin \text{m and} \ a \notin \text{trecs}\]

\[\rightarrow (id, m, \text{in, out} \otimes \text{Msg}(id, id_p, \text{Sh-req}, a, \bot), \text{Trec}(a, (id_k, \text{Sh-req})) | \text{trecs})\]

\[\text{where} \ id_p = \text{parent}(id)\]

Discard-Redundant-Sh-Req-Due-To-Shared-Copy Rule

\[(id, \text{Cell}(a,v,\text{cs,R(id_k,dir)})) | m, \text{Msg}(id_k, id, \text{Sh-req}, a, \bot) \odot \text{in, out, trecs})\]

\[\rightarrow (id, \text{Cell}(a,v,\text{cs,R(id_k,dir)})) | m, \text{in, out, trecs})\]

Discard-Redundant-Sh-Req-Due-To-Exclusive-Copy Rule

\[(id, \text{Cell}(a,v,\text{Ex,W(id_k)})) | m, \text{Msg}(id_k, id, \text{Sh-req}, a, \bot) \odot \text{in, out, trecs})\]

\[\rightarrow (id, \text{Cell}(a,v,\text{Ex,W(id_k)})) | m, \text{in, out, trecs})\]

2.4.2 Ex-Request Rules

Receive-Ex-Req-And-Send-Ex-Rep Rule

\[(id, \text{Cell}(a,v,\text{Ex,R(id_k)})) | m, \text{Msg}(id_k, id, \text{Ex-req}, a, \bot) \odot \text{in, out, trecs}) \quad \text{if} \ a \notin \text{trecs}\]

\[\rightarrow (id, \text{Cell}(a,v,\text{Ex,W(id_k)})) | m, \text{in, out} \otimes \text{Msg}(id, id_k, \text{Ex-req}, a, v), \text{trecs})\]

Receive-Ex-Req-And-Send-Upgrade-Rep Rule

\[(id, \text{Cell}(a,v,\text{Ex,R(id_k)})) | m, \text{Msg}(id_k, id, \text{Ex-req}, a, \bot) \odot \text{in, out, trecs}) \quad \text{if} \ a \notin \text{trecs}\]

\[\rightarrow (id, \text{Cell}(a,v,\text{Ex,W(id_k)})) | m, \text{in, out} \otimes \text{Msg}(id, id_k, \text{Upgrade-req}, a, v), \text{trecs})\]

Receive-Ex-Req-And-Multicast-Inv-Req Rule

\[(id, \text{Cell}(a,v,\text{Ex,R(dir)})) | m, \text{Msg}(id_k, id, \text{Ex-req}, a, \bot) \odot \text{in, out, trecs})\]

\[\text{if} \ dir - \{id_k\} \neq \epsilon \quad \text{and} \ a \notin \text{trecs}\]

\[\rightarrow (id, \text{Cell}(a,v,\text{Ex,R(dir)})) | m, \text{in, out} \otimes \text{multicast}(id, dir - \{id_k\}, \text{inv-req}, a, \bot), \text{Trec}(a, (id_k, \text{Ex-req})) | \text{trecs})\]

Receive-Ex-Req-And-Send-Pushout-Req Rule

\[(id, \text{Cell}(a,v,\text{Ex,W(id_k)})) | m, \text{Msg}(id_k, id, \text{Ex-req}, a, \bot) \odot \text{in, out, trecs})\]

\[\text{if} \ id_k \neq id \quad \text{and} \ a \notin \text{trecs}\]

\[\rightarrow (id, \text{Cell}(a,v,\text{Ex,W(id_k)})) | m, \text{in, out} \otimes \text{Msg}(id, id_j, \text{Pushout-req}, a, \bot), \text{Trec}(a, (id_k, \text{Ex-req})) | \text{trecs})\]

Receive-Ex-Req-And-Send-Ex-Req Rule

\[(id, m, \text{Msg}(id_k, id, \text{Ex-req}, a, \bot) \odot \text{in, out, trecs})\]

\[\text{if Cell}(a,-,\text{Ex,-}) \neq \text{m and} \ a \notin \text{trecs}\]

\[\rightarrow (id, m, \text{in, out} \otimes \text{Msg}(id, id_p, \text{Ex-req}, a, \bot), \text{Trec}(a, (id_k, \text{Ex-req})) | \text{trecs})\]

\[\text{where} \ id_p = \text{parent}(id)\]

Discard-Redundant-Ex-Req Rule

\[(id, \text{Cell}(a,v,\text{Ex,W(id_k)})) | m, \text{Msg}(id_k, id, \text{Ex-req}, a, \bot) \odot \text{in, out, trecs})\]

\[\rightarrow (id, \text{Cell}(a,v,\text{Ex,W(id_k)})) | m, \text{in, out, trecs})\]

2.5 Parent-to-Child Request Processing Rules

2.5.1 Wb-Request Rules

Receive-Wb-Req-And-Send-Wb-Req Rule

\[(id, \text{Cell}(a,v,\text{Ex,R(dir)})) | m, \text{Msg}(id_p, id, \text{Wb-req}, a, \bot) \odot \text{in, out, trecs}) \quad \text{if} \ a \notin \text{trecs}\]

\[\rightarrow (id, \text{Cell}(a,v,\text{Sh,R(dir)})) | m, \text{in, out} \otimes \text{Msg}(id, id_p, \text{Wb-req}, a, v), \text{trecs})\]
Receive-Wb-Req-And-Send-Wb-Req Rule
(id, Cell(a,v,(Ex,W(idk)))) | m, Msg(idp, id, Wb-req, a, ⊥)⊙ in, out, trecs) if a \notin trecs
→ (id, Cell(a,v,(Ex,W(idk)))) | m, in, out⊙ Msg(id, idk, Wb-req, a, ⊥), Trec(a,(idp,Wb-req)) | trecs)

Discard-Redundant-Wb-Req Rule
(id, m, Msg(idp, id, Wb-req, a, ⊥)⊙ in, out, trecs) if Cell(a, -,(Ex,-)) \notin m
→ (id, m, in, out, trecs)

2.5.2 Pushout-Request Rules

Receive-Pushout-Req-And-Send-Pushout-Rep Rule
(id, Cell(a,v,(Ex,R(ε)))) | m, Msg(idp, id, Pushout-req, a, ⊥)⊙ in, out, trecs) if a \notin trecs
→ (id, m, in, out⊙ Msg(id, idp, Pushout-req, a, v), trecs)

Receive-Pushout-Req-And-Multicast-Inv-Req Rule
(id, Cell(a,v,(Ex,R(dir)))) | m, Msg(idp, id, Pushout-req, a, ⊥)⊙ in, out, trecs)
if dir ≠ ε and a \notin trecs
→ (id, Cell(a,v,(Ex,R(dir)))) | m, in, out⊙ multicast(id, dir, Inv-req, a, ⊥), Trec(a,(idp,Pushout-req)) | trecs)

Receive-Pushout-Req-And-Send-Pushout-Req Rule
(id, Cell(a,v,(Ex,W(idk)))) | m, Msg(idp, id, Pushout-req, a, ⊥)⊙ in, out, trecs) if a \notin trecs
→ (id, Cell(a,v,(Ex,W(idk)))) | m, in, out⊙ Msg(id, idk, Pushout-req, a, ⊥), Trec(a,(idp,Pushout-req)) | trecs)

Receive-Unexpected-Pushout-Req-And-Send-Inv-Rep Rule
(id, Cell(a,v,(Sh,R(ε)))) | m, Msg(idp, id, Pushout-req, a, ⊥)⊙ in, out, trecs)
→ (id, m, in, out⊙ Msg(id, idp, Inv-req, a, ⊥), trecs)

Receive-Unexpected-Pushout-Req-And-Multicast-Inv-Req Rule
(id, Cell(a,v,(Sh,R(dir)))) | m, Msg(idp, id, Pushout-req, a, ⊥)⊙ in, out, trecs) if dir ≠ ε
→ (id, Cell(a,v,(Sh,R(dir)))) | m, in, out⊙ multicast(id, dir, Inv-req, a, ⊥), Trec(a,(idp,Inv-req)) | trecs)

Discard-Redundant-Pushout-Req Rule
(id, m, Msg(idp, id, Pushout-req, a, ⊥)⊙ in, out, trecs) if a \notin m
→ (id, m, in, out, trecs)

2.5.3 Inv-Request Rules

Receive-Inv-Req-And-Send-Inv-Req Rule
(id, Cell(a,v,(Sh,R(ε)))) | m, Msg(idp, id, Inv-req, a, ⊥)⊙ in, out, trecs)
→ (id, m, in, out⊙ Msg(id, idp, Inv-req, a, ⊥), trecs)

Receive-Inv-Req-And-Multicast-Inv-Req Rule
(id, Cell(a,v,(Sh,R(dir)))) | m, Msg(idp, id, Inv-req, a, ⊥)⊙ in, out, trecs) if dir ≠ ε
→ (id, Cell(a,v,(Sh,R(dir)))) | m, in, out⊙ multicast(id, dir, Inv-req, a, ⊥), Trec(a,(idp,Inv-req)) | trecs)

Discard-Redundant-Inv-Req Rule
(id, m, Msg(idp, id, Inv-req, a, ⊥)⊙ in, out, trecs) if Cell(a, -,(Sh,-)) \notin m
→ (id, m, in, out, trecs)
2.6 Parent-to-Child Reply Processing Rules

2.6.1 Sh-Reply Rules

*Receive-Sh-Rep-And-Execute-Load Rule*

\[ \text{Sys}(\text{id}, \text{msg}(\text{id}_p, \text{id}, \text{Sh-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia}, \text{Load})) | \text{trecs}), \quad \text{Proc}(\text{ia}, \text{rf}, \text{prog})] \\
\quad \text{if} \quad \text{prog}[\text{ia}] = \text{r} := \text{Load}(r_1) \quad \text{and} \quad \text{a} = \text{rf}[r_1] \\
\rightarrow \quad \text{Sys}(\text{id}, \text{Cell}(a, v, (\text{Sh}, R(e))) | \text{m}, \text{in}, \text{out}, \text{trecs}), \quad \text{Proc}(\text{ia}+1, \text{rf}[r := v], \text{prog})) \\

*Receive-Sh-Rep-And-Send-Sh-Rep Rule*

\[ \text{id}, \text{msg}(\text{id}_p, \text{id}, \text{Sh-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia}_k, \text{Sh-req})) | \text{trecs} \]
\[ \rightarrow \quad \text{id}, \text{Cell}(a, v, (\text{Sh}, R(id_k))) | \text{m}, \text{in}, \text{out} \otimes \text{Msg}(\text{id}, \text{id}_k, \text{Sh-rep}, a, v), \text{trecs} \]

*Receive-Unexpected-Sh-Rep-With-Suspended-Store Rule*

\[ \text{id}, \text{msg}(\text{id}_p, \text{id}, \text{Sh-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia}, \text{Store})) | \text{trecs} \]
\[ \rightarrow \quad \text{id}, \text{Cell}(a, v, (\text{Sh}, R(e))) | \text{m}, \text{in}, \text{out}, \text{Trec}(a, (\text{ia}, \text{Store})) | \text{trecs} \]

*Receive-Unexpected-Sh-Rep-With-Suspended-Ex-Req Rule*

\[ \text{id}, \text{msg}(\text{id}_p, \text{id}, \text{Sh-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia}_k, \text{Ex-req})) | \text{trecs} \]
\[ \rightarrow \quad \text{id}, \text{Cell}(a, v, (\text{Sh}, R(e))) | \text{m}, \text{in}, \text{out}, \text{Trec}(a, (\text{ia}_k, \text{Ex-req})) | \text{trecs} \]

*Receive-Unexpected-Sh-Rep-Without-Initiator Rule*

\[ \text{id}, \text{msg}(\text{id}_p, \text{id}, \text{Sh-rep}, a, v) \odot \text{in}, \text{out}, \text{trecs} \] \quad \text{if} \quad a \notin \text{trecs} \\
\rightarrow \quad \text{id}, \text{Cell}(a, v, (\text{Sh}, R(e))) | \text{m}, \text{in}, \text{out}, \text{trecs} \]

2.6.2 Ex-Reply Rules

*Receive-Ex-Rep-And-Execute-Store Rule*

\[ \text{Sys}(\text{id}, \text{msg}(\text{id}_p, \text{id}, \text{Ex-rep}, a, u) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia}, \text{Store})) | \text{trecs}), \quad \text{Proc}(\text{ia}, \text{rf}, \text{prog})] \\
\quad \text{if} \quad \text{prog}[\text{ia}] = \text{Store}(r_1, r_2) \quad \text{and} \quad \text{a} = \text{rf}[r_1] \\
\rightarrow \quad \text{Sys}(\text{id}, \text{Cell}(a, v, (\text{Ex}, R(e))) | \text{m}, \text{in}, \text{out}, \text{trecs}), \quad \text{Proc}(\text{ia}+1, \text{rf}, \text{prog}) \quad \text{where} \quad v = \text{rf}[r_2] \\

*Receive-Ex-Rep-And-Send-Ex-Rep Rule*

\[ \text{id}, \text{msg}(\text{id}_p, \text{id}, \text{Ex-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia}_k, \text{Ex-req})) | \text{trecs} \]
\[ \rightarrow \quad \text{id}, \text{Cell}(a, v, (\text{Ex}, W(id_k))) | \text{m}, \text{in}, \text{out} \otimes \text{Msg}(\text{id}, \text{id}_k, \text{Ex-rep}, a, v), \text{trecs} \]

*Receive-Unexpected-Ex-Rep-With-Suspended-Load Rule*

\[ \text{Sys}(\text{id}, \text{msg}(\text{id}_p, \text{id}, \text{Ex-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia}, \text{Load})) | \text{trecs}), \quad \text{Proc}(\text{ia}, \text{rf}, \text{prog})] \\
\quad \text{if} \quad \text{prog}[\text{ia}] = \text{r} := \text{Load}(r_1) \quad \text{and} \quad \text{a} = \text{rf}[r_1] \\
\rightarrow \quad \text{Sys}(\text{id}, \text{Cell}(a, v, (\text{Ex}, R(e))) | \text{m}, \text{in}, \text{out}, \text{trecs}), \quad \text{Proc}(\text{ia}+1, \text{rf}[r := v], \text{prog})) \\

*Receive-Unexpected-Ex-Rep-With-Suspended-Sh-Req Rule*

\[ \text{id}, \text{msg}(\text{id}_p, \text{id}, \text{Ex-rep}, a, v) \odot \text{in}, \text{out}, \text{Trec}(a, (\text{ia}_k, \text{Sh-req})) | \text{trecs} \]
\[ \rightarrow \quad \text{id}, \text{Cell}(a, v, (\text{Ex}, R(id_k))) | \text{m}, \text{in}, \text{out} \otimes \text{Msg}(\text{id}, \text{id}_k, \text{Sh-rep}, a, v), \text{trecs} \]

*Receive-Unexpected-Ex-Rep-Without-Initiator Rule*

\[ \text{id}, \text{msg}(\text{id}_p, \text{id}, \text{Ex-rep}, a, v) \odot \text{in}, \text{out}, \text{trecs} \] \quad \text{if} \quad a \notin \text{trecs} \\
\rightarrow \quad \text{id}, \text{Cell}(a, v, (\text{Ex}, R(e))) | \text{m}, \text{in}, \text{out}, \text{trecs} \]
2.6.3 Upgrade-Reply Rules

Receive-Upgrade-Rep-And-Execute-Store Rule
Sys((id, Cell(a,u,(Sh,R(e)))) | m, Msg(idp, id, Upgrade-rep, a, ⊥) ⊗ in, out, Trec(a, (ia,Store)) | trecs),
Proc(ia, rf, prog) if prog[ia] = Store(r1, r2) and a = rf[r1]
→ Sys((id, Cell(a,v,(Ex,R(e)))) | m, in, out, trecs), Proc(ia+1, rf, prog) where v = rf[r2]

Receive-Upgrade-Rep-And-Send-Ex-Req Rule
(id, Cell(a,v,(Sh,R(id)))) | m, Msg(idp, id, Upgrade-rep, a, ⊥) ⊗ in, out, Trec(a, (idk,Ex-req)) | trecs)
→ (id, Cell(a,v,(Ex,W(idk)))) | m, in, out ⊗ Msg(id, idk, Ex-rep, a, v), trecs)

Receive-Upgrade-Rep-And-Send-Upgrade-Rep Rule
(id, Cell(a,v,(Sh,R(idk)))) | m, Msg(idp, id, Upgrade-rep, a, ⊥) ⊗ in, out, Trec(a, (idk,Ex-req)) | trecs)
→ (id, Cell(a,v,(Ex,W(idk)))) | m, in, out ⊗ Msg(id, idk, Upgrade-rep, a, ⊥), trecs)

Receive-Upgrade-Rep-And-Multicast-Inv-Req Rule
(id, Cell(a,v,(Sh,R(dir)))) | m, Msg(idp, id, Upgrade-rep, a, ⊥) ⊗ in, out, Trec(a, (idk,Ex-req)) | trecs)
if dir-{idk} ̸= ε
→ (id, Cell(a,v,(Ex,R(dir)))) | m, in, out ⊗ multicast(dir-{idk}, Inv-req, a, ⊥), Trec(a, (idk,Ex-req)) | trecs)

2.7 Child-to-Parent Reply Processing Rules

2.7.1 Wb-Reply Rules

Receive-Wb-Rep-And-Send-Sh-Rep Rule
(id, Cell(a,u,(Ex,W(idk)))) | m, Msg(idk, id, Wb-rep, a, v) ⊗ in, out, Trec(a, (idj,Sh-req)) | trecs)
→ (id, Cell(a,v,(Ex,R(idk)))) | m, in, out ⊗ Msg(id, idj, Sh-rep, a, v), trecs)

Receive-Wb-Rep-And-Send-Wb-Rep Rule
(id, Cell(a,u,(Ex,W(idk)))) | m, Msg(idk, id, Wb-rep, a, v) ⊗ in, out, Trec(a, (idp,Wb-req)) | trecs)
→ (id, Cell(a,v,(Sh,R(idk)))) | m, in, out ⊗ Msg(id, idp, Wb-rep, a, v), trecs)

Receive-Unexpected-Wb-Rep-With-Suspended-Ex-Req Rule
(id, Cell(a,u,(Ex,W(idk)))) | m, Msg(idk, id, Wb-rep, a, v) ⊗ in, out, Trec(a, (idj,Ex-req)) | trecs)
→ (id, Cell(a,v,(Ex,R(idk)))) | m, in, out, Trec(a, (idj,Ex-req)) | trecs)

Receive-Unexpected-Wb-Rep-With-Suspended-Pushout-Req Rule
(id, Cell(a,u,(Ex,W(idk)))) | m, Msg(idk, id, Wb-rep, a, v) ⊗ in, out, Trec(a, (idj,Pushout-req)) | trecs)
→ (id, Cell(a,v,(Ex,R(idk)))) | m, in, out, Trec(a, (idj,Pushout-req)) | trecs)

Receive-Unexpected-Wb-Rep-Without-Initiator Rule
(id, Cell(a,u,(Ex,W(idk)))) | m, Msg(idk, id, Wb-rep, a, v) ⊗ in, out, trecs) if a ̸∈ trecs
→ (id, Cell(a,v,(Ex,R(idk)))) | m, in, out, trecs)

2.7.2 Pushout-Reply Rules

Receive-Pushout-Rep-And-Send-Ex-Rep Rule
(id, Cell(a,u,(Ex,W(idk)))) | m, Msg(idk, id, Pushout-rep, a, v) ⊗ in, out, Trec(a, (idj,Ex-req)) | trecs)
→ (id, Cell(a,v,(Ex,W(idj)))) | m, in, out ⊗ Msg(id, idj, Ex-rep, a, v), trecs)

Receive-Pushout-Rep-And-Send-Pushout-Rep Rule
(id, Cell(a,u,(Ex,W(idk)))) | m, Msg(idk, id, Pushout-rep, a, v) ⊗ in, out, Trec(a, (idp,Pushout-req)) | trecs)
→ (id, m, in, out ⊗ Msg(id, idp, Pushout-rep, a, v), trecs)
Receive-Unexpected-Pushout-Rep-With-Suspended-Sh-Req Rule
\( \{id, Cell(a,u,(Ex,W(id_k))) \} m, Msg(id_k, id, Pushout-rep, a, v) \odot in, out, Trec(a, (id_j, Sh-req)) | trecs \)  
\( \rightarrow \{id, Cell(a,v,(Ex,R(id_j))) \} m, in, out \otimes Msg(id, id_j, Sh-rep, a, v), trecs \)  

Receive-Unexpected-Pushout-Rep-With-Suspended-Wb-Req Rule
\( \{id, Cell(a,u,(Ex,W(id_k))) \} m, Msg(id_k, id, Pushout-rep, a, v) \odot in, out, Trec(a, (id_p, Wb-req)) | trecs \)  
\( \rightarrow \{id, Cell(a,v,(Sh,R(e))) \} m, in, out \otimes Msg(id, id_p, Wb-rep, a, v), trecs \)  

Receive-Unexpected-Pushout-Rep-Without-Initiator Rule
\( \{id, Cell(a,u,(Ex,W(id_k))) \} m, Msg(id_k, id, Pushout-rep, a, v) \odot in, out, trecs \) \( a \notin trecs \)  
\( \rightarrow \{id, Cell(a,v,(Ex,R(e))) \} m, in, out, trecs \)  

2.7.3 Inv-Reply Rules

Receive-Inv-Rep-With-Ex-Req-Suspended Rule
\( \{id, Cell(a,v,(Ex,R(id_k),dir)) \} m, Msg(id_k, id, Inv-rep, a, \bot) \odot in, out, Trec(a, (id_j, Ex-req)) | trecs \)  
\( \text{if dir} - \{id_j\} \neq e \)  
\( \rightarrow \{id, Cell(a,v,(Ex,R(dir))) \} m, in, out, Trec(a, (id_j, Ex-req)) | trecs \)  

Receive-Inv-Rep-With-Inv-Req-Suspended Rule
\( \{id, Cell(a,v,(Sh,R(id_k),dir)) \} m, Msg(id_k, id, Inv-rep, a, \bot) \odot in, out, Trec(a, (id_p, Inv-req)) | trecs \)  
\( \text{if dir} \neq e \)  
\( \rightarrow \{id, Cell(a,v,(Sh,R(dir))) \} m, in, out, Trec(a, (id_j, Inv-req)) | trecs \)  

Receive-Inv-Rep-With-Pushout-Req-Suspended Rule
\( \{id, Cell(a,v,(Ex,R(id_k),dir)) \} m, Msg(id_k, id, Inv-rep, a, \bot) \odot in, out, Trec(a, (id_p, Pushout-req)) | trecs \)  
\( \text{if dir} \neq e \)  
\( \rightarrow \{id, Cell(a,v,(Ex,R(dir))) \} m, in, out, Trec(a, (id_p, Pushout-req)) | trecs \)  

Receive-Inv-Rep-And-Send-Ex-Rep Rule
\( \{id, Cell(a,v,(Ex,R(id_k))) \} m, Msg(id_k, id, Inv-rep, a, \bot) \odot in, out, Trec(a, (id_j, Ex-req)) | trecs \)  
\( \rightarrow \{id, Cell(a,v,(Ex,W(id_k))) \} m, in, out \otimes Msg(id, id_j, Ex-rep, a, v), trecs \)  

Receive-Inv-Rep-And-Send-Upgrade-Rep Rule
\( \{id, Cell(a,v,(Ex,R(id_k))) \} m, Msg(id_k, id, Inv-rep, a, \bot) \odot in, out, Trec(a, (id_j, Ex-req)) | trecs \)  
\( \rightarrow \{id, Cell(a,v,(Ex,W(id_k))) \} m, in, out \otimes Msg(id, id_j, Upgrade-rep, a, \bot), trecs \)  

Receive-Inv-Rep-And-Send-Inv-Rep Rule
\( \{id, Cell(a,v,(Sh,R(id_k))) \} m, Msg(id_k, id, Inv-rep, a, \bot) \odot in, out, Trec(a, (id_p, Inv-req)) | trecs \)  
\( \rightarrow \{id, m, in, out \otimes Msg(id, id_p, Inv-rep, a, \bot), trecs \} \)  

Receive-Inv-Rep-And-Send-Pushout-Rep Rule
\( \{id, Cell(a,v,(Ex,R(id_k))) \} m, Msg(id_k, id, Inv-rep, a, \bot) \odot in, out, Trec(a, (id_p, Pushout-req)) | trecs \)  
\( \rightarrow \{id, m, in, out \otimes Msg(id, id_p, Pushout-rep, a, v), trecs \} \)  

Receive-Unexpected-Inv-Rep-Without-Initiator Rule
\( \{id, Cell(a,v,(cs,R(id_k),dir)) \} m, Msg(id_k, id, Inv-rep, a, \bot) \odot in, out, trecs \) \( a \notin trecs \)  
\( \rightarrow \{id, Cell(a,v,(cs,R(dir))) \} m, in, out, trecs \)  

2.8 Message Passing Rules

Message-Passing-To-Child Rule
\( Sys(id, m, in, Msg(id, id_k, cmd, a, v) \otimes out, trecs), Sys(id_k, m_k, in_k, out_k, trecs_k), eu_k | sg \)  
\( \rightarrow \)  
\( Sys(id, m, in, out, trecs), Sys(id_k, m_k, in_k \otimes Msg(id, id_k, cmd, a, v), out_k, trecs_k), eu_k | sg \)
Message-Passing-To-Parent Rule

\[ \text{Sys}(\{\text{id}_k, \text{m}_k, \text{in}_k, \text{Msg}(\text{id}_k, \text{id}, \text{cmd}, a, \nu) \otimes \text{out}_k, \text{trecs}_k\}, \text{eu}_k) \mid \text{sg}) \]

\[ \rightarrow \text{Sys}(\{\text{id}_k, \text{m}_k, \text{in}_k \circ \text{Msg}(\text{id}_k, \text{id}, \text{cmd}, a, \nu), \text{out}, \text{trecs}\}, \text{Sys}(\{\text{id}_k, \text{m}_k, \text{in}_k, \text{out}_k, \text{trecs}_k\}, \text{eu}_k) \mid \text{sg}) \]

2.9 Buffer Management

2.10 Discussion

Voluntary Request Issue Rules: In theory, the same voluntary request message can be issued repeatedly. In practice, when a voluntary request message is issued, we can encode this information in Trecs to prevent the same request messages being issued multiple times. Moreover, we can distinguish between voluntary request messages and normal request messages by encoding a bit information in the message format. An advantage of this is that voluntary request messages can be discarded without affecting the correctness of the protocol. We may want to do this to control flow of message passing, or the voluntary request message is deemed inappropriate upon arrival. The protocol liveness liveness is guaranteed regardless of how voluntary messages are issued.

2.11 FIFO Message Passing

In HCN-adpt, there are two scenarios that reply messages can be reordered due to non-FIFO message passing: \text{inv-rep} overtakes \text{Wb-rep}, and \text{Upgrade-rep} overtakes \text{ShRep}. When this happens, the reply message that arrives first (i.e., \text{inv-rep} and \text{Upgrade-rep}) cannot be processed before the other reply message (i.e., \text{Wb-rep} and \text{Sh-rep}) is processed (unless some other rules are added, which will inevitably increase the number of states). In HCN-adpt, we deal with this issue by imposing circular buffer management for reply messages.

With FIFO message passing, this reordering cannot happen. Thus, the circular buffer management is no longer needed.

\[ \text{msg}_1 \circ \text{msg}_2 \equiv \text{msg}_2 \circ \text{msg}_1 \]

if \( \text{addr}(\text{msg}_1) \neq \text{addr}(\text{msg}_2) \) or \( \text{dest}(\text{msg}_1) \neq \text{dest}(\text{msg}_1) \)

\[ \text{msg}_1 \circ \text{msg}_2 \equiv \text{msg}_2 \circ \text{msg}_1 \]

if \( \text{addr}(\text{msg}_1) \neq \text{addr}(\text{msg}_2) \) or \( \text{src}(\text{msg}_1) \neq \text{src}(\text{msg}_1) \)

**FIFO Message Passing:** We assume message passing is FIFO in the sense that messages that have the same source and destination are received in the same order as they are issued if they are regarding the same memory location.