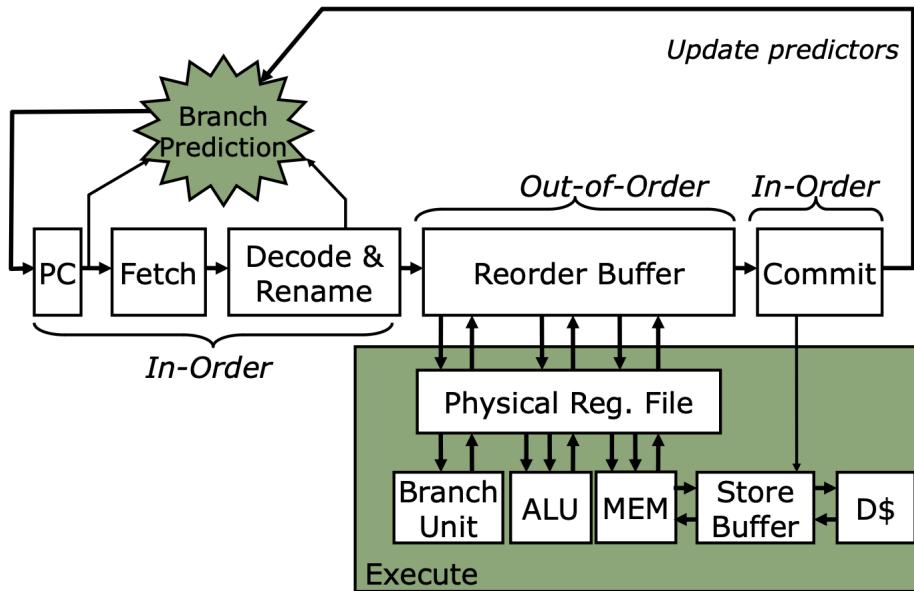


# Quiz 1 Handout - Physical Register File Design

This handout defines a simple RISC-V Out-of-Order processor. All register data is stored directly in a physical register file. The ROB simply contains tags that point to physical registers.

The processor follows the “Physical Register File” design seen in lecture. A diagram representing the processor is shown in Figure 1.

Figure 1: Processor Overview



The processor follows the RISC-V ISA, though we will only focus on integer and control instructions, so you do not need to worry about memory instructions.

The processor has the following key stages:

1. **Fetch:** the instruction at PC is fetched from the instruction cache/memory
2. **Decode & Rename:** the fetched instruction is decoded. If the decoded instruction is a conditional branch, its direction is predicted by the branch predictor. The instruction’s architectural source registers are renamed to the appropriate physical registers. A destination physical register is assigned.
3. **Reorder Buffer:** the instruction waits to be executed in the reorder buffer. When all hazards are cleared, it is able to execute.
4. **Commit:** the instruction is committed.

Figure 2: Current Processor State

Rename Table
Physical Register File
Free List

Arch. Reg.	Physical Reg.
x0	
x1	P7
x2	P8
x3	
x4	P4
x5	P6
x6	P9
x7	

Physical Reg.	Value	Present?
P0	37	
P1	38	
P2	823	
P3	5900	
P4	2816	
P5	0	
P6	1123	
P7	314	
P8	217	
P9	415	
P10		

Physical Reg.
P10

Reorder Buffer

Next to Commit


Next Available


Inum	Use	Ex	Op	PR1 Present?	PR1	PR2 Present?	PR2	Rd	Last PRd	PRd	
...											
I5	1	1	addi		1	P0			x1	P0	P1
I6	1	0	xor			P2		P4	x2	P2	P3
I7	1	0	div			P3		P5	x5	P5	P6
I8	1	0	bne			P1		P9			
I9	1	0	add			P1			x1	P1	P7
I10	1	0	xor			P3		P4	x2	P3	P8

- **Rename Table:** mapping between architectural registers and physical registers
- **Physical Register File:** central physical register file that holds all register data
- **Free List:** available physical registers not currently assigned to any architectural register
- **Reorder Buffer:** holds inflight instructions and the tags of their input and output registers

### **Label List**

- A. Satisfy a dependence on \_\_\_\_\_ by stalling
- B. Satisfy a dependence on \_\_\_\_\_ by bypassing a speculative value
- C. Satisfy a dependence on \_\_\_\_\_ by bypassing a committed value
- D. Satisfy a dependence on \_\_\_\_\_ by speculation using a static prediction
- E. Satisfy a dependence on \_\_\_\_\_ by using a dynamic prediction
- F. Write a speculative value using lazy data management
- G. Write a speculative value using greedy data management
- H. Speculatively update a prediction on \_\_\_\_\_ using lazy value management
- I. Speculatively update a prediction on \_\_\_\_\_ using greedy value management
- J. Non-speculatively update a prediction on \_\_\_\_\_
- K. Check the correctness of a speculation on \_\_\_\_\_ and find a correct speculation
- L. Check the correctness of a speculation on \_\_\_\_\_ and find an incorrect speculation
- M. Abort speculative action and cleanup lazily managed values
- N. Abort speculative action and cleanup greedily managed values
- O. Commit correctly speculated instruction, where there was no value management
- P. Commit correctly speculated instruction, and mark lazily updated values as non-speculative
- Q. Commit correctly speculated instruction, and free log associated with greedily updated values
- R. Illegal or broken actions

### **Blank Choices**

- i. Register Value
- ii. PC value
- iii. Branch direction
- iv. Memory address
- v. Memory value
- vi. Latency of operation
- vii. Functional unit
- viii. Storage space